

# 1 Primes

Primes would seem to be the ultimate in precision. A number 317 is either prime or it isn't (this one is!), there is no approximation to its primality. Nonetheless, Asymptopia is the proper place to examine primes in the aggregate.

**Definition 1** For  $n \geq 2$ ,  $\pi(n)$  denotes the number of primes  $p$  with  $2 \leq p \leq n$ .

Our goal in this chapter is to show one of the great theorems of mathematics.

## Theorem 1.1 (The Prime Number Theorem)

$$\pi(n) \sim \frac{n}{\ln n} \quad (1)$$

This result was first conjectured in the early nineteenth century. (While the conjecture is sometimes attributed to Gauss the history is murky.) It was a central problem for that century, finally being proven independently by Hadamard and Vallée-Poussin in 1898. Their proofs involved complex variables and a long search continued for an elementary proof. This was finally obtained in 1949 by Selberg and Erdős. Still, a full proof of Theorem 1 is beyond the limits of this work. We shall come close to it with the following results:

**Theorem 1.2** *There exists a positive constant  $c_1$  such that*

$$(c_1 + o(1)) \frac{n}{\ln n} \leq \pi(n) \quad (2)$$

*That is,  $\pi(n) = \Omega(n/\ln n)$ . Further, our argument gives  $c_1 = \ln 2$ .*

**Theorem 1.3** *There exists a positive constant  $c_2$  such that*

$$\pi(n) \leq (c_2 + o(1)) \frac{n}{\ln n} \quad (3)$$

*That is,  $\pi(n) = O(n/\ln n)$ . Further, our argument gives  $c_2 = 2 \ln 2$ .*

Together, Theorem 1.2, 1.3 yield:

$$\pi(n) = \Theta\left(\frac{n}{\ln n}\right) \quad (4)$$

With more effort we shall show

**Theorem 1.4** *If there exists a positive constant  $c$  such that*

$$\pi(n) \sim c \frac{n}{\ln n} \quad (5)$$

then  $c = 1$ .

## 1.1 Fun with Primes

*A Break! No asymptotics in this section!*

How many factors of the prime 7 are there in 100!? The numbers 7, 14, ..., 98 all have a factor of 7 so that gives  $\frac{98}{7} = 14$  factors. And, 49 and 98 have a second factor of 7 which gives an additional  $\frac{98}{49} = 2$  factors. In total there are  $16 = 14 + 2$  factors of 7.

**Definition 2** *For  $n \geq 1$  and  $p$  prime,  $v_p(n)$  denotes the number of factors  $p$  in  $n$ . Equivalently,  $v_p(n)$  is that nonnegative integer  $a$  such that  $p^a$  divides  $n$  but  $p^{a+1}$  does not divide  $n$ .*

**Theorem 1.5** *For any  $n \geq 1$  and  $p$  prime*

$$v_p(n!) = \sum_{i=1}^{\infty} \lfloor \frac{n}{p^i} \rfloor \quad (6)$$

*Equivalently*

$$v_p(n!) = \sum_{i=1}^s \lfloor \frac{n}{p^i} \rfloor \text{ with } s = \lfloor \log_p n \rfloor \quad (7)$$

When  $i > \lfloor \log_p n \rfloor$ ,  $p < n^i$  so the addend in (6), explaining the equivalence. The argument with  $p = 7, n = 100$  easily generalizes. For any  $i \leq s$  there are  $\lfloor np^{-i} \rfloor$  numbers  $1 \leq j \leq n$  that have (at least)  $i$  factors of  $p$ . We count each such  $i$  and  $j$  once, as then an  $i$  with precisely  $u$  factors of  $p$  will be counted precisely  $u$  times.

We apply Theorem 1.5 to study binomial coefficients. Let  $n = a + b$  and set  $C = \binom{n}{a} = \frac{n!}{a!b!}$ . Applying (7)

$$v_p(C) = v_p(n!) - v_p(a!) - v_p(b!) = \sum_{i=1}^s \lfloor \frac{n}{p^i} \rfloor - \lfloor \frac{a}{p^i} \rfloor - \lfloor \frac{b}{p^i} \rfloor \quad (8)$$

with  $s = \lfloor \log_p n \rfloor$  as in (7).

**Theorem 1.6** *With  $n = a + b$ ,  $p$  prime, and  $C = \binom{n}{a}$ ,*

$$0 \leq v_p(C) \leq \lfloor \log_p n \rfloor \quad (9)$$

**Proof:** Set  $\alpha = ap^{-i}$ ,  $\beta = bp^{-i}$ . Then the addend in (8) is

$$\lfloor \alpha + \beta \rfloor - \lfloor \alpha \rfloor - \lfloor \beta \rfloor \tag{10}$$

This term is zero if the fractional parts of  $\alpha, \beta$  sum to less than one and one if they sum to one or more. The sum (8) consists of  $s = \lfloor \log_p n \rfloor$  terms, each one or zero, and so lies between 0 and  $s$ .

**Remark:** With  $n = a + b$  there are two arguments why  $a!b!$  divides  $n!$ . One: the proof of Theorem 8 gives that, for all primes  $p$ ,  $v_p(n!) \geq v_p(a!) + v_p(b!) = v_p(a!b!)$  and thus  $a!b!$  divides  $n!$ . Two: The quotient  $\frac{n!}{a!b!} = \binom{n}{a}$  counts the  $a$ -subsets of an  $n$ -sets and hence must be a nonnegative integer. Which proof one prefers is an esthetic question <sup>1</sup> but it is frequently useful to know more than one proof of a theorem.

There is an amusing way of calculating  $v_p(C)$  with  $C = \binom{n}{a}$  and  $a+b = n$ . Write  $a, b$  in base  $p$ . Add them (in base  $p$ ) so that you will get  $n$  in base  $p$ .

**Theorem 1.7**  $v_p(C)$  is the number of carries when you add  $a, b$  getting  $n$ , all in base  $p$ .

For example, let  $a = 33$ ,  $b = 25$  so  $n = 58$  (written in decimal), and set  $p = 7$ . In base 7,  $a = 45$ ,  $b = 34$ . When we add them <sup>2</sup>

$$\begin{array}{r} 45 \\ + 34 \\ ---- \\ 112 \end{array}$$

There we two carries and  $v_7\left(\binom{45}{34}\right) = 2$ .

We indicate the argument. For each  $1 \leq i$  we get a carry from the  $i-1$ -st place (counting from the right, starting at 0) to the  $i$ -th place if and only if the fractional parts of  $ap^{-i}$  and  $bp^{-i}$  add to at least one and that occurs if and only if term (10) is one.

## 1.2 PMT - Lpper Bound

Let  $n$  be even ( $n$  odd will be similar). The upper *and* lower bounds come from examining the prime factorization of binomial coefficients. Set  $r = \pi(n)$

<sup>1</sup>This author prefers the “counts” argument.

<sup>2</sup>To paraphrase the wonderful songwriter Tom Lehrer, base seven is just like base ten – if you are missing three fingers!

and let  $p_1, \dots, p_r$  denote the primes up to  $n$  and write

$$\binom{n}{n/2} = p_1^{\alpha_1} p_2^{\alpha_2} \cdots p_r^{\alpha_r} \quad (11)$$

(There might not be a factor of  $p_i$ . In that case we simply write  $\alpha_i = 0$ .) We rewrite the upper bound of Theorem 1.6 as:

$$p_i^{\alpha_i} \leq n \quad (12)$$

Thus

$$\binom{n}{n/2} \leq n^r \quad (13)$$

Stirling's Formula gives an asymptotic formula for  $\binom{n}{n/2}$  but here we use only the weaker  $\binom{n}{n/2} = 2^{n(1+o(1))}$ . Taking  $\ln$  of both sides of (13) and dividing gives

$$\pi(n) = r \geq \frac{\ln \binom{n}{n/2}}{\ln n} = \frac{n}{\ln n} (\ln 2)(1 + o(1)) \quad (14)$$

What if  $n$  is odd? In Asymptopia we simply apply (14) to the even  $n - 1$ . Thus

$$\pi(n) \geq \pi(n - 1) \geq \frac{\ln \binom{n-1}{(n-1)/2}}{\ln(n - 1)} \quad (15)$$

which is again  $\frac{n}{\ln n} (\ln 2)(1 + o(1))$ .

### 1.3 PMT-Upper Bound

Again assume  $n$  is even. There are  $\pi(n) - \pi(n/2)$  primes  $p$  with  $\frac{n}{2} < p < n$ . Each of them appears in  $\binom{n}{n/2}$  to the first power. (They appear once in the numerator as a factor of  $p$  and never in the denominator.) Thus, with the product over these primes,

$$\prod p \leq \binom{n}{n/2} \quad (16)$$

We again do not need a more precise estimate and here simply bound  $\binom{n}{n/2} \leq 2^n$ . Each factor  $p$  is a factor of at least  $\frac{n}{2}$ . Thus

$$\left(\frac{n}{2}\right)^{\pi(n) - \pi(n/2)} \leq 2^n \quad (17)$$

Taking  $\ln$  of both sides gives

$$\pi(n) - \pi\left(\frac{n}{2}\right) \leq \frac{n}{\ln(n/2)}(\ln 2) \quad (18)$$

For  $n = 2k + 1$  odd we apply the same argument to  $\binom{n}{k}$  getting an upper bound on  $\pi(n) - \pi(k + 1)$ . We combine the even and odd cases by writing

$$\pi(n) - \pi\left(\lceil \frac{n}{2} \rceil\right) \leq \frac{n}{\ln(n/2)}(\ln 2) \quad (19)$$

Turning (19) into an upper bound on  $\pi(n)$  is a typical problem in Asymptopia. Set  $x_0 = n$  and  $x_{i+1} = \lceil \frac{x_i}{2} \rceil$ . This sequence decreases until finally reaching  $x_s = 1$ . Applying (19) to  $n = x_0, \dots, x_{s-1}$  and adding we get

$$\pi(n) \leq \sum_{i=0}^{s-1} \frac{x_i}{\ln(x_i/2)}(\ln 2) \quad (20)$$

In the exact world this would be a daunting sum. In Asymptopia we will split the sum into the main terms and the small terms. Where to make the split is part of the *art* of Asymptopia which we discuss further below. For now, let  $u$  be the first index with  $x_u \leq n \ln^{-2} n$ . Applying (19) only down to  $x_{u-1}$  and adding we get

$$\pi(n) - \pi(x_u) \leq \sum_{i=0}^{u-1} \frac{x_i}{\ln(x_i/2)}(\ln 2) \quad (21)$$

Now we use the trivial bound  $\pi(x_u) \leq x_u \leq n \ln^{-2} n$ . While this is a “bad” bound for  $\pi(x_u)$  it is a negligible value for us and

$$\pi(n) \leq o\left(\frac{n}{\ln n}\right) + \sum_{i=0}^{u-1} \frac{x_i}{\ln(x_i/2)}(\ln 2) \quad (22)$$

As  $x_i$  is decreasing so is the denominator  $\ln(x_i/2)$  which pushes the sum (22) up. However, all terms in the sum have  $x_i/2 > n \ln^{-2} n/2$ . The  $\ln$  function is going down, but not too far down. Each denominator

$$\ln(x_i/2) \geq \ln(n \ln^{-2} n/2) = \ln n - 2 \ln \ln n - \ln 2 = (1 - o(1)) \ln n \quad (23)$$

Thus

$$\sum_{i=0}^{u-1} \frac{x_i}{\ln(x_i/2)}(\ln 2) \leq \frac{1 + o(1)}{(\ln n)(\ln 2)} \sum_{i=0}^{u-1} x_i \quad (24)$$

Now  $x_0 = n$  and  $x_i \sim n2^{-i}$  (indeed, to be totally formal,  $x_i \leq n2^{-i} + 1$ ) so that

$$\sum_{i=0}^{u-1} x_i \leq 2n(1 + o(1)) \quad (25)$$

and (22) gives

$$\pi(n) \leq \frac{n}{\ln n} \frac{2}{\ln 2} (1 + o(1)) \quad (26)$$

**Selecting the Split:** When we chose  $u$  above there was a lot of room but still, care had to be taken. Knowing the answer in advance helps. Suppose we let  $u$  be the first index with  $x_u < S$  and consider which values of  $S$  might work. It helps (as is frequently the case) to know <sup>3</sup> that  $\pi(n) = \Theta(n/\ln n)$ . In the argument we will be adding  $S$  and so we want  $S = o(n/(\ln n))$ . But also the densities are going down in  $i$  when we look at  $\pi(x_i) - \pi(x_{i+1})$  and we want them all to be  $(1 + o(1))/(\ln n)$ . As the last one will be  $\sim 1/\ln(S)$  we will want  $\ln(S) \sim \ln(n)$  which in turn requires  $S = n^{1-o(1)}$ . Indeed, any  $S = n^{1-o(1)}$  with  $S \ll (n/(\ln n))$  could have been used. Looking ahead at the argument we will be adding  $S$ . This leads us to require that  $S = o(n/\ln n)$ . Having finished the argument it is instructive to look back. The main intervals are roughly  $[n, n/2), [n/2, n/4), \dots$ . In the first interval the upper bound for the density of primes from (19) is roughly  $2/(\ln n)(\ln 2)$ . This upper bound continues down to  $S$ , as  $\ln(S) \sim \ln(n)$ . Thus the upper bound on the total number of primes is at most  $S$  (which we choose to be negligible) plus what the number of primes would be if each interval had prime density  $\frac{2}{\ln 2} \frac{1}{\ln n}$ . The intervals total at most  $n$  values (actually a bit less since we cut it off at  $S$ ) and so the main contribution to the prime count is  $\sim \frac{2}{\ln n} \frac{n}{\ln n}$ .

## 1.4 PMT with Constant

**Note:** This section gets quite technical and should be considered optional.

Here we show Theorem 1.4. That is, we *assume* that there is a constant  $c$  such that  $\pi(n) \sim c(n/(\ln n))$  and then show that  $c$  must be 1. It is a big *if*. *A priori*, from Theorems 1.2,1.3 the ratio of  $\pi(n)$  to  $n/(\ln n)$  could oscillate between two positive constants, never approaching a limit.

We consider the factorization (11) more carefully. Our goal will be to show that if  $c \neq 1$  then the left and right hand sides cannot match. We split the primes from 1 to  $n$  into intervals. We shall let  $K$  be a large but fixed

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<sup>3</sup>Actually, a good hunch is useful. If the hunch turns out to be wrong the calculations will not come out as you wanted.

constant. (More about just how large later.) For  $1 \leq i < K$  let  $P_i$  denote the set of primes  $p$  with

$$\frac{n}{i+1} < p \leq \frac{n}{i} \quad (27)$$

and let  $SP$  (small primes) denote the set of primes  $p$  with  $p < \frac{n}{K}$ . Let  $V_i$ ,  $1 \leq i < K$  denote the contribution of the  $p \in P_i$  to the factorization (11). That is,  $V_i$  is the product of  $p_j^{\alpha_j}$  in (11), where  $p_j$  is restricted to  $P_i$ . Similarly let  $V_{SP}$  denote the contribution of the  $p \in SP$  to the factorization (11). That is,  $V_i$  is the product of  $p_j^{\alpha_j}$  in (11), where  $p_j$  is restricted to  $SP$ .

We first show that  $SP$  makes a relatively small contribution to (11). There are  $\leq \pi(n/K)$  primes  $p \in SP$  and each (12) contributes at most a factor of  $n$  so that  $V_{SP} \leq n^{\pi(n/K)}$ . From Theorem 1.3 gives  $\pi(n/K) < (2 \ln 2 + o(1))(n/K)/\ln(n/K)$ . With  $K$  fixed,  $\ln(n/K) \sim \ln(n)$  so that  $\pi(n/K) < (\ln 2 + o(1))(n/K)/\ln(n)$ . Thus (27),

$$V_{SP} < n^{(2 \ln 2 + o(1))(n/K)/\ln(n)} = 2^{(2n/K)(1+o(1))} \quad (28)$$

so that

$$\ln(V_{SP}) < \frac{2n \ln 2}{K}(1 + o(1)) \quad (29)$$

While this is not a small number in absolute terms it will be relatively small compared to the total contribution which is  $2^{n(1+o(1))}$ .

For  $1 \leq i < K$  we now look at  $V_i$ . As all primes considered have  $p > \frac{n}{K}$  and  $K$  is fixed they have  $p > \sqrt{n}$ . Thus the sum of Theorem 1.5 has only one term. Theorem 1.6 with  $a = n/2$  is then simply

$$v_p\left(\binom{n}{n/2}\right) = \lfloor n/p \rfloor - 2\lfloor n/2p \rfloor \quad (30)$$

This is either zero or one and is one precisely when  $\lfloor n/p \rfloor$  is odd. We have *designed*  $P_i$  so that  $\lfloor n/p \rfloor = i$  for  $p \in P_i$ . When  $i$  is even no primes  $p \in P_i$  appear in the factorization (11) (or, the same thing, they appear with exponent zero) and so  $V_i = 1$ . (For example, with  $\frac{n}{7} < p \leq \frac{n}{6}$ ,  $n!$  has six factors of  $p$  and  $(n/2)!$  has twice three factors of  $p$  and they all cancel.)

Now suppose  $1 \leq i < K$  is odd. Then  $V_i$  is simply the product of all primes  $p \in P_i$ . Each such prime  $p$  lies between  $\frac{n}{K}$  and  $n$  and so can be considered  $p = n^{1+o(1)}$ . The number of such primes is  $\pi(n/i) - \pi(n/(i+1))$ . In this range  $\ln(n/i) \sim \ln n$ . Our assumption for Theorem 1.5 then gives that  $\pi(n/i) \sim c \frac{n}{i \ln n}$  and that  $\pi(n/(i+1)) \sim c \frac{n}{(i+1) \ln n}$ . We deduce that the number of primes is  $\sim c \frac{n}{\ln n} \left(\frac{1}{i} - \frac{1}{i+1}\right)$ . (**Caution:** Subtraction in Asymptopia is dangerous! It is critical here that  $i \leq K$  and that  $K$  is a

fixed constant, so  $\frac{1}{i}$  and  $\frac{1}{i+1}$  is a positive constant. Were, say,  $K = \ln \ln n$  we could not do the subtraction. With  $i \sim (\ln \ln n)/2$ , for example, the asymptotics of  $\pi(n/i)$  and  $\pi(n/(i+1))$  would be the same and so one could *not* deduce the asymptotics of their difference!) Thus

$$V_i = n^{c(1+o(1))(n/(\ln n))(\frac{1}{i}-\frac{1}{i+1})} \quad (31)$$

and

$$\ln(V_i) \sim cn\left(\frac{1}{i} - \frac{1}{i+1}\right) \quad (32)$$

From the factorization (11) Then

$$\ln\left(\binom{n}{n/2}\right) = \ln V_{SP} + \sum \ln(V_i) \quad (33)$$

For convenience, assume  $K = 2T$  is even so we can write the odd  $i < K$  as  $2j - 1$ ,  $1 \leq j \leq T$ . From Chapter xxx, the left hand side is  $\sim n \ln 2$ . Thus

$$(1 + o(1))n \ln 2 = cn(1 + o(1)) \sum j = 1^T \left(\frac{1}{2j-1} - \frac{1}{2j}\right) + \ln V_{SP} \quad (34)$$

Dividing by  $n$

$$(1 + o(1))(\ln 2) = c(1 + o(1)) \sum_{k=1}^{2T-1} \frac{(-1)^{k+1}}{k} + \frac{1}{n} \ln V_{SP} \quad (35)$$

We need <sup>4</sup> the fact that

$$\ln 2 = \sum_{k=1}^{\infty} \frac{(-1)^{k+1}}{k} = 1 - \frac{1}{2} + \frac{1}{3} - \frac{1}{4} + \dots \quad (36)$$

We can now see the idea. The  $\ln(V_{SP})$  will be negligible and (35) becomes  $\ln 2 = c(\ln 2)$ . The actual argument consists of eliminating all  $c \neq 1$ .

Suppose  $c > 1$ . Select  $K = 2T$  so that  $c \sum_{k=1}^{2T-1} \frac{(-1)^{k+1}}{k} > \ln 2$ . As  $\ln V_{SP} \geq 0$  the right hand side of (35) would be bigger than the left hand side.

Suppose  $c < 1$ . Applying the upper bound (29), the right hand side of (35) would be at most  $c \sum_{k=1}^{2T-1} \frac{(-1)^{k+1}}{k} + \frac{2 \ln 2}{K}$ . As  $K \rightarrow \infty$ , this sum approaches  $c \ln 2$  which is less than  $\ln 2$ . Thus we may select  $K$  <sup>5</sup> so that

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<sup>4</sup>Again, from Calculus!

<sup>5</sup>A subtle wrinkle here, while we examine behavior as  $K \rightarrow \infty$  we select  $K$  a constant, dependent only on  $c$ .



this sum is less than  $\ln 2$ . But now the right hand side of (35) would be smaller than the left hand side.

Both assumptions led to a contradiction and since we *assumed* that  $c$  existed, it must be that  $c = 1$ .

## 1.5 Telescoping

Suppose we have a reasonable function  $f(x)$  and we wish to asymptotically evaluate  $\sum_{p \leq n} f(p)$ . We assume the Prime Number Theorem 1, giving the asymptotics of  $\pi(s)$  as  $s \rightarrow \infty$ . On an intuitive level we think of  $1 \leq s \leq n$  as being prime with “probability”  $\pi(s)/s \sim 1/(\ln s)$ . Then  $s$ ,  $1 \leq s \leq n$  would contribute  $f(s)/(\ln s)$  to the sum and  $\sum_{p \leq n} f(p)$  would be roughly  $\sum_{s \leq n} f(s)/(\ln s)$ . This is not a proof, integers are either prime or they aren’t, yet surprisingly we can often get this intuitive result. The key is called telescoping. We write

$$\sum_{p \leq n} f(p) = \sum_{s=2}^n f(s)(\pi(s) - \pi(s-1)) \quad (37)$$

Reversing sums (and noting  $\pi(1) = 0$ )

$$\sum_{s=2}^n f(s)(\pi(s) - \pi(s-1)) = f(n)\pi(n) + \sum_{s=2}^{n-1} \pi(s)(f(s) - f(s+1)) \quad (38)$$

While (38) its effectiveness depends on our ability to asymptotically calculate the sum. An important success is when  $f(s) = \frac{1}{s}$ , we ask for the asymptotics of

$$F(n) = \sum_{p \leq n} \frac{1}{p} \quad (39)$$

The first term of (38) is then  $\sim \frac{1}{n} \frac{n}{\ln n} = o(1)$ . The sum is asymptotically  $\sum \frac{s}{\ln s} \frac{1}{s(s+1)} \sim \sum \frac{1}{s \ln s}$ , the sum from  $s = 1$  to  $n - 1$ . From Chapter xxx,

$$\sum_{s=2}^{n-1} \frac{1}{s \ln s} \sim \int_1^n \frac{dx}{x \ln x} = \ln \ln n \quad (40)$$

That is,  $F(n) \sim \ln \ln n$ . For another example, take  $f(s) = s$  so that  $F(n) = \sum_{p \leq n} p$ . Then

$$F(n) = n\pi(n) - \sum_{s=2}^{n-1} \pi(s) \sim \frac{n^2}{\ln n} - \int_2^{n-1} \frac{s}{\ln s} ds \quad (41)$$

While the integrand cannot be precisely integrated we can handle it in Asymptopia. Our notion is that  $\ln s \sim \ln n$  for “most”  $2 \leq s \leq n - 1$ . We split the integral at some  $n^{1-o(1)}$ , let us take  $u(n) = n \ln^{-10} n$  for definiteness. For  $u(n) \leq s$ ,  $\ln(s) \geq \ln n - 10 \ln \ln n \sim \ln n$  so that

$$\int_{u(n)}^{n-1} \frac{s}{\ln s} ds \sim \int_{u(n)}^{n-1} \frac{s}{\ln n} ds \sim \frac{n^2}{2 \ln n} \quad (42)$$

For  $s \leq u(n)$  we bound  $\frac{s}{\ln s} \leq s$  so that

$$\int_2^{u(n)} \frac{s}{\ln s} ds \leq \int_0^{u(n)} s ds \sim \frac{n^2}{2 \ln^{20} n} \quad (43)$$

As the upper bound (43) is  $o(n^2/\ln n)$  it has a negligible effect and the total integral

$$\int_2^{n-1} \frac{s}{\ln s} ds \sim \frac{n^2}{2 \ln n} \quad (44)$$

Subtracting, (41) gives

$$\sum_{p \leq n} p \sim \frac{n^2}{\ln n} - \frac{n^2}{2 \ln n} \sim \frac{n^2}{2 \ln n} \quad (45)$$