

1. **Improving the Sudan list-decoding algorithm:** In class we saw how to list decode RS codes from agreements as small as  $2\sqrt{kn}$ . Improve this to  $\sqrt{2kn}$ . The best-known algorithm, due to Guruswami and Sudan, needs only  $\sqrt{kn}$  agreement and is based on multiplicities. See Sudan's lecture notes or Guruswami's thesis.
2. **Yekhanin's approach to locally decodable codes:** A binary code  $C : \{0, 1\}^n \rightarrow \{0, 1\}^N$  is said to be  $(q, \delta, \varepsilon)$ -locally decodable if there exists a randomized decoding algorithm  $\mathcal{A}$  that on input  $y \in \{0, 1\}^N$  and  $i \in [n]$  where  $y$  satisfies  $\delta(y, C(x)) \leq \delta$  for some  $x \in \{0, 1\}^n$  outputs  $x_i$  with probability at least  $1 - \varepsilon$  and makes at most  $q$  queries to  $y$ . In class we saw a simple  $(2, \frac{1}{4} - \varepsilon, \frac{1}{2} - 2\varepsilon)$ -local decoder for Hadamard codes.

**Definition 1** Let  $N, R, n \geq 1$  be some integers. For each  $i \in [n], r \in [R]$ , let  $T_i$  and  $Q_{ir}$  be subsets of  $[N]$ . We say that  $T_i$  and  $Q_{ir}$  form a  $(q, n, N, R, s)$  regular intersecting family if the following holds:

- (a)  $q$  is odd;
- (b) For all  $i \in [n], |T_i| = s$ ;
- (c) For all  $i \in [n]$  and  $r \in [R], Q_{ir} \subseteq T_i$  and  $|Q_{ir}| = q$ ;
- (d) For all  $i \in [n]$  and  $w \in T_i, |\{r \in [R] \mid w \in Q_{ir}\}| = Rq/s$  (i.e.,  $T_i$  is uniformly covered by the  $Q_{ir}$ );
- (e) For all  $i, j \in [n], i \neq j$  and  $r \in [R], |Q_{ir} \cap T_j| \equiv 0 \pmod 2$ .

Show how to use a  $(q, n, N, R, s)$  regular intersecting family to construct binary linear code encoding  $n$  bits into  $N$  bits that is  $(q, \delta, \delta Nq/s)$  locally decodable for all  $\delta > 0$ .

3. **One small step in Yekhanin's construction of a regular intersecting family:** Let  $p$  be an odd prime number and  $m \geq p - 1$  be an integer. Define  $n = \binom{m}{p-1}$ . Show that there exist two families of vectors  $\{u_1, \dots, u_n\}$  and  $\{v_1, \dots, v_n\}$  in  $\mathbb{F}_p^m$  such that
  - for all  $i \in [n], \langle u_i, v_i \rangle = 0$ ;
  - for all  $i, j \in [n], i \neq j$ , we have  $\langle u_i, v_j \rangle \neq 0$ .