

# HW9 Solutions

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## 1. Memory Mapping

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### 1.1

3 entries. The page size is 4096, so map1 consumes 2 entries, and map2 consumes one entry.

### 1.2

2 pages. The 3 virtual pages map to two physical pages in the OS buffer cache.

## 2. Thread Switching

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```
void swtch(struct thread *t1, struct thread *t2) {
    // On entry this function is run by thread t1.

    push_register(RBP);
    push_register(R0);
    push_register(R1);
    push_register(R2);
    push_register(R3);

    t1->stack = read_register(RSP);
    write_register(RSP, t2->stack);

    pop_register(R3);
    pop_register(R2);
    pop_register(R1);
    pop_register(R0);
    pop_register(RBP);

    return; // The function should return to the
           // point where thread t2 called swtch.
}
```

## 3. Interrupts

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### 3.1

$200 \text{ words/min} \times 1 \text{ min/60 seconds} \times 7 \text{ letters/word} \times 1 \text{ interrupt/character} = 23 \text{ interrupts/second.}$

### 3.2

It should use interrupts. Your interrupts cost the computer 23 microseconds out of every second; this is a trivial fraction (23 parts in one million). If the computer used polling, you (the human) would notice the lag and get annoyed.

## 4. Disk Calculations

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### 4.1

20 GB/disk.  $10 \text{ platters/disk} \times 4096 \text{ tracks/platter} \times 1024 \text{ sectors/track} \times 512 \text{ bytes/sector} = 10 \times 4 \times 1024 \times 1024 \times 512 = 40 \times 1 \text{ MB} \times 512 = 40 \times .5 \text{ GB} = 20 \text{ GB}.$

### 4.2

100MB/second. First note that 12,000 RPM = 200 rotations per second (or one rotation per 5 ms). In a single rotation, we can read an entire track. A track consists of  $512 \text{ bytes/sector} \times 1024 \text{ sectors} = 0.5 \text{ MB}.$  So the sequential transfer bandwidth is  $200 \text{ rotations/second} \times 0.5 \text{ MB/rotation} = 100 \text{ MB/second}.$  Because the track-to-track seek time and the I/O bus overhead are both modeled as negligible, 100 MB/second is our answer.

### 4.3

Roughly 30 KB/second. First note that in one read, we get 512 bytes. What is the time to issue this read? The disk incurs seek delay and rotational delay. The average seek latency is 15 ms. After the disk head reaches the desired track, the disk has to wait until the desired sector rotates under the disk head. Since the sector could be anywhere on the track, ranging from right under the head to the most pessimal position, the average rotational delay is 2.5 ms (half of the 5 ms per rotation). So the total delay on average is 17.5 ms. Our total effective bandwidth, then, is  $512 \text{ bytes} / 17.5 \text{ ms} \approx 525 / 17.5 \text{ bytes/ms} = 30,000 \text{ bytes/second}.$

## 5. Disk Scheduling

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### Shortest-seek-time-first

20 - 22 - 10 - 6 - 2 - 38 - 40

So the time is  $(2 + 12 + 4 + 4 + 36 + 2) \times 6 \text{ ms} = 360 \text{ ms}$

### LOOK

If travelling up first:

20 - 22 - 38 - 40 - 10 - 6 - 2

So the seek time is  $(2 + 16 + 2 + 30 + 4 + 4) \times 6 \text{ ms} = 348 \text{ ms}$

Or if it travels down first:

20 - 10 - 6 - 2 - 22 - 38 - 40

$(10 + 4 + 4 + 20 + 16 + 2) \times 6 \text{ ms} = 336 \text{ ms}$