

# HW10 Solutions

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## 1. File Systems

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### 1.1

The data pointed to by the direct block pointers can be as large as  $1024 * 8$  bytes. There is a single indirect block pointer, and it can contain as many as  $1024 / 4$  pointers, each of which points to 1024 bytes of data. Thus, the total is:

$$1024 * 8 + 1024 * (1024 / 4) = (2^{13} + 2^{18}) = 270,336 \text{ B.}$$

### 1.2

Assume directories are implemented as files. Also, each directory entry is 16 bytes (14 bytes for the file name and 2 bytes for the inode number, since there are  $2^{16}$  inodes). Thus, we take our answer to the previous question and divide by 16 bytes:

$$2^{13} + 2^{18} / 2^4 = 2^9 + 2^{14} = 16896.$$

Since there are  $2^{16}$  inodes on the disk, our answer is  $\min\{2^{16}, 16896\} = 16896$ .

## 2. Synchronous Operations

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Synchronous operations cannot return until both locations are updated and persistent. If these two inode arrays are stored on the first and second half of the same disk, it requires a seek for all operations. Reads, however, could be slightly faster: the system could seek to the closer version of the inode (but since there is only one copy of the data itself, the total gain is likely to be limited).

## 3. Virtual Memory Issue

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Your friend forgot that contiguous virtual addresses may not be contiguous in physical memory. For example, if we have mappings:

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va: 0x1000 --> pa: 0x8000
va: 0x2000 --> pa: 0x5000
```

then the scheme described would do the wrong thing for disk writes larger than a page size: the disk would write pa 0x8000 through 0x9FFF.