New York University CSCI-UA.0202-003: Operating Systems (Undergrad): Fall 2025

Quiz 3

•	Write	your	full	name	on	both:
---	-------	------	------	------	----	-------

- the bubble sheet in the "Name" field
- the quiz booklet
- Write your NYU NetID on the quiz booklet and the bubble sheet in the "ID" field
- Use a #2 pencil to fill in your answers on the bubble sheet (preferred, but you can also use a pen)
- This quiz contains 6 questions only. Each question has choices from A to D
- Fill the bubbles completely by darkening the entire circle, as shown in the example
- Only mark answers for questions 1-6. Do not mark any bubbles beyond question 6
- Choose only one answer per question
- Submit your bubble sheet together with your quiz booklet

Name:			
NetId:			

- 1. According to what we discussed in lectures, why must a thread that calls <code>cond_wait()</code> use a while loop to re-check the condition upon waking?
 - (A) To ensure the lock is held when calling <code>cond_wait()</code> , which is a requirement.
 - (B) To handle spurious wakeups, where a thread might be awakened even if the condition it's waiting for is not yet true.
 - (C) To prevent other threads from signaling the same condition variable simultaneously.
 - (D) To guarantee that the thread with the highest priority is always awakened first.
- 2. In the context of a Monitor, what is the relationship between its mutex and its condition variables?
 - (A) The mutex is only used to protect the condition variables, not the shared data.
 - (B) A monitor can have multiple mutexes but only one condition variable.
 - (C) The mutex that protects the monitor's methods is the same mutex that must be held when operating on its condition variables.
 - (D) Condition variables operate independently and do not require the monitor's mutex to be held.
- 3. In the mutex implementation we introduced in the lecture, a thread attempting to acquire a held lock performs these steps (among others): adds itself to the waiters queue, calls sched_mark_blocked(), releases the internal spinlock, and finally calls sched_swtch(). Why is it critical to release the internal spinlock *before* calling sched_swtch() to go to sleep?
 - (A) To allow the sched_swtch() function to reuse the spinlock for its own internal synchronization.
 - (B) To ensure the thread still holds the spinlock when it wakes up, so it can resume safely.
 - (C) To prevent the scheduler from immediately waking the thread back up again, which is known as a spurious wakeup.
 - (D) To avoid no other thread could ever acquire it to modify the mutex's state (e.g., to release the mutex).

- 4. The entire reader-writer lock mechanism uses several components working together: integer counters (AR, AW, WR, WW), two condition variables (okToRead and okToWrite), and a single Mutex. The condition variables are responsible for putting threads to sleep and waking them up. Given that, what is the essential role of the single Mutex in this mechanism?
 - (A) It serves as the primary lock on the database itself; no thread can read or write without holding this <code>mutex</code> .
 - (B) It is a "master switch" that determines whether the lock should currently favor readers or writers.
 - (C) It protects the integer counters (AR , AW , etc.) from race conditions, ensuring that when one thread is changing their values, no other thread can read or write them at the same time.
 - (D) It is only used to prevent two threads from calling wait() on the same condition variable at the exact same time.
- 5. According to "An Investigation of the Therac-25 Accidents", what was the name of the company that manufactured the Therac-25?
 - (A) General Electric (GE)
 - (B) Atomic Energy of Canada Limited (AECL)
 - (C) Siemens AG
 - (D) Varian Medical Systems
- 6. "An Investigation of the Therac-25 Accidents" emphasizes that the Therac-25 incidents were "system accidents." Which of the following best describes what this means?
 - (A) The accidents were caused exclusively by a single, catastrophic software bug that was impossible to detect.
 - (B) The disasters resulted from a complex interplay of factors, including flawed software design, the removal of hardware safeguards, and poor organizational procedures.
 - (C) The operators were the primary cause of the accidents due to their repeated failure to follow the correct procedures.
 - (D) The machine's hardware was fundamentally unreliable and would have failed regardless of the software installed.