Sound Verification by Abstract Interpretation

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Motivation

Formal methods

Reasonings on programs are

- Reasonings on properties of their semantics (*i.e.* execution behaviors)
- Always involve some form of abstraction

Abstract interpretation

A theory establishing a correspondance between

- Concrete semantic properties

  ↑ what you want to prove on the semantics

- Abstract properties

  ↑ how to prove it in the abstract

Objective: formalize

- formal methods
- algorithms for reasoning on programs

Fundamental motivations

Scientific research

in Mathematics/Physics:

trend towards unification and synthesis through universal principles

in Computer science:

trend towards dispersion and parcelization through a collection of local techniques for specific applications

An exponential process, will stop!

Example: reasoning on computational structures

WCET
Security protocole verification
Systems biology analysis
Operational semantics
Axiomatic semantics
Dataflow analysis
Model checking
Database query
Abstraction refinement
Confidentiality analysis
Partial evaluation
Obfuscation
Dependence analysis
Type inference
Program synthesis
Effect systems
Denotational semantics
Separation logic
Grammar analysis
CEGAR
Theories combination
Program transformation
Termination proof
Statistical model-checking
Trace semantics
Code contracts
Interpolants
Abstract model checking
Shape analysis
Invariance proof
Symbolic execution
Integrity analysis
Malware detection
Probabilistic verification
Quantum entanglement detection
Bisimulation
SMT solvers
Code refactoring
Parsing
Type theory
Steganography
Tautology testers

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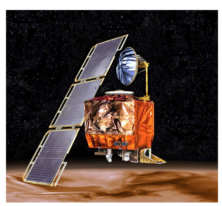
Practical motivations

All computer scientists have experienced bugs


Ariane 5.01 failure
(overflow)


Patriot failure
(float rounding)


Mars orbiter loss
(unit error)


Heartbleed
(buffer overrun)

Checking the presence of bugs by debugging is great

Proving their absence by static analysis is even better!

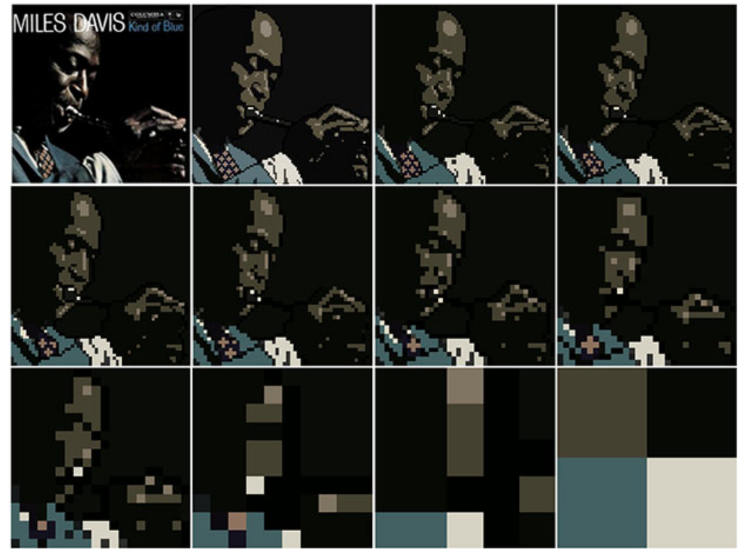
Undecidability and complexity is the challenge for automation

Informal examples of abstraction

Abstractions of Dora Maar by Picasso



Pixelation

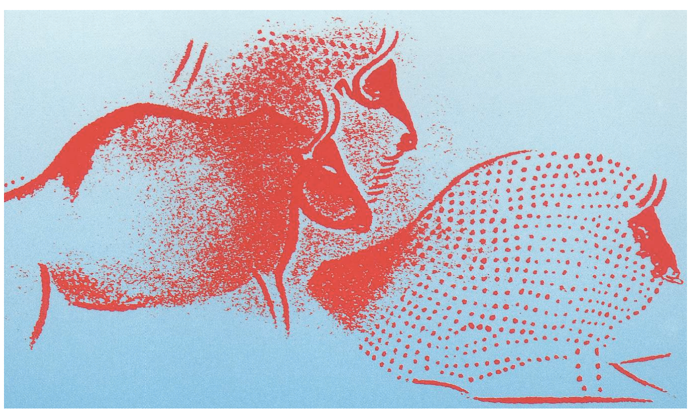


/www.petapixel.com/2011/06/23/how-much-pixelation-is-needed-before-a-photo-becomes-transformed/

Image credit: *Photograph by Jay Maisel*

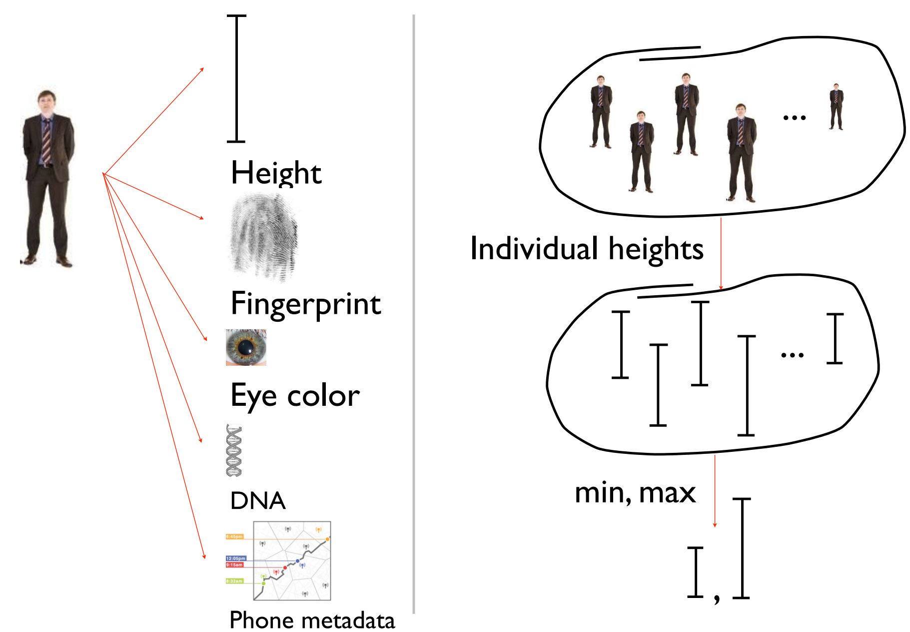
An old idea...

20 000 years old picture in a spanish cave:

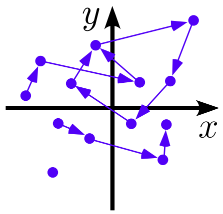


(the concrete is unknown)

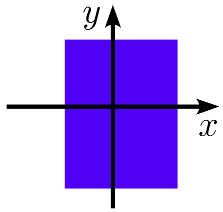
Abstractions of a man / crowd



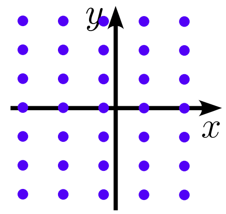
Numerical abstractions in Astrée



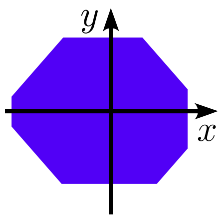
Collecting semantics:[1,5]
partial traces



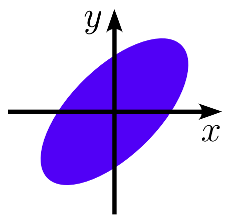
Intervals:[20]
$x \in [a, b]$



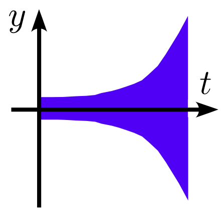
Simple congruences:[24]
$x \equiv a[b]$



Octagons:[25]
$\pm x \pm y \leqslant a$



Ellipses:[26]
$x^2 + by^2 - axy \leqslant d$



Exponentials:[27]
$-a^{bt} \leqslant y(t) \leqslant a^{bt}$

Difficulties

Making it easy…

No induction:

- Model-checking finite systems
- Decidable cases

No soundness: the last trend to fall in the easy, e.g.

- Analyze Linux the easy way (ignoring aliases, overflows, recursion, etc.) → 700 potential bugs
- Ask PhD students to analyze manually the potential bug (3mn per bug maximum)
- Claim 50 true bugs → best paper award

Abstract Interpretation

Abstract interpretation is all about:

Soundness

Induction

A very short introduction to abstract interpretation

Patrick Cousot & Radhia Cousot. Vérification statique de la cohérence dynamique des programmes. In *Rapport du contrat IRIA SESORI No 75-035*, Laboratoire IMAG, University of Grenoble, France. 125 pages. 23 September 1975.

Patrick Cousot & Radhia Cousot. Static Determination of Dynamic Properties of Programs. In B. Robinet, editor, *Proceedings of the second international symposium on Programming*, Paris, France, pages 106—130, April 13-15 1976, Dunod, Paris.

Patrick Cousot, Radhia Cousot: Abstract Interpretation: A Unified Lattice Model for Static Analysis of Programs by Construction or Approximation of Fixpoints. POPL 1977: 238-252

Patrick Cousot, Radhia Cousot: Systematic Design of Program Analysis Frameworks. POPL 1979: 269-282

Patrick Cousot. Méthodes itératives de construction et d'approximation de points fixes d'opérateurs monotones sur un treillis, analyse sémantique des programmes. *Thèse És Sciences Mathématiques*, Université Joseph Fourier, Grenoble, France, 21 March 1978

Patrick Cousot. Semantic foundations of program analysis. In S.S. Muchnick & N.D. Jones, editors, *Program Flow Analysis: Theory and Applications,* Ch. 10, pages 303—342, Prentice-Hall, Inc., Englewood Cliffs, New Jersey, U.S.A., 1981.

Properties and their Abstractions

Concrete properties

A concrete property is represented by the set of elements which have that property:

- universe (set of elements) $\mathscr{D}$ (e.g. a semantic domain)
- properties of these elements: $P \in \wp(\mathscr{D})$
- "x has property P" is $x \in P$

$\langle \wp(\mathscr{D}), \subseteq, \cup, \cap, ... \rangle$ is a *complete lattice* for inclusion $\subseteq$ (i.e. *logical implication*)

Abstract properties

Abstract properties: $Q \in \mathscr{A}$

Abstract domain $\mathscr{A}$: encodes a subset of the concrete properties (e.g. a program logic, type terms, linear algebra, *etc*)

Poset: $\langle \mathscr{A}, \sqsubseteq, \sqcup, \sqcap, ... \rangle$

Partial order: $\sqsubseteq$ is *abstract implication*

Concretization

Concretization $\gamma \in \mathscr{A} \longrightarrow \wp(\mathscr{D})$

$\gamma(Q)$ is the semantics (concrete meaning) of Q

γ is *increasing* (so $\sqsubseteq$ abstracts $\subseteq$)

The concrete properties in $\gamma(\mathscr{A})$ are exactly representable in the abstract $\mathscr{A}$, all others in $\wp(\mathscr{D}) \setminus \gamma(\mathscr{A})$ can only be approximated in $\mathscr{A}$

Best abstraction

A concrete property $P \in \wp(\mathscr{D})$ has a best abstraction $Q \in \mathscr{A}$ iff

- it is sound (over-approximation):

$$P \subseteq \gamma(Q)$$

- and more precise than any sound abstraction:

$$P \subseteq \gamma(Q') \implies Q \sqsubseteq Q' \implies \gamma(Q) \subseteq \gamma(Q')$$

The best abstraction is unique (by antisymmetry)

Under-approximation is order-dual

Galois connection

Any $P \in \wp(\mathscr{D})$ has a (unique) best abstraction $\alpha(P)$ in $\mathscr{A}$ if and only if

$$\forall P \in \wp(\mathscr{D}) : \forall Q \in \mathscr{A} : \ \alpha(P) \sqsubseteq Q \iff P \subseteq \gamma(Q)$$

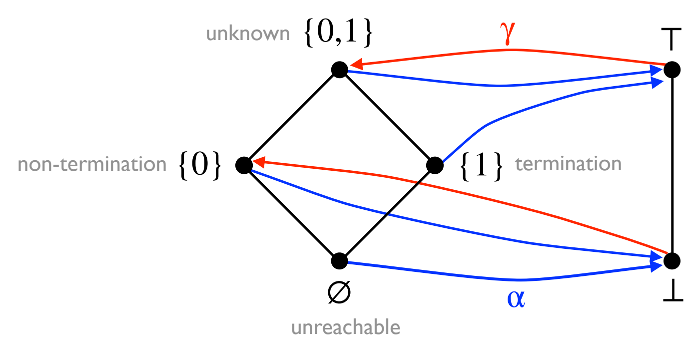
$\Rightarrow$: *over-approximation*
$\Leftarrow$: *best abstraction*

written

$$\langle \wp(\mathscr{D}), \subseteq \rangle \xleftarrow[\alpha]{\gamma} \langle \mathscr{A}, \sqsubseteq \rangle$$

Examples

Needness/strictness analysis (80's)



Similar abstraction ($\gamma(\top) \triangleq \{$true, false$\}$) for scalable harware symbolic trajectory evaluation STE (90)

Alan Mycroft: The Theory and Practice of Transforming Call-by-need into Call-by-value. Symposium on Programming 1980: 269-281

Carl-Johan H. Seger, Randal E. Bryant: Formal Verification by Symbolic Evaluation of Partially-Ordered Trajectories. Formal Methods in System Design 6(2): 147-189 (1995)

Example: Homomorphic abstraction $\wp(\mathscr{D}) \longrightarrow \wp(\mathscr{A})$

$\hbar \in \mathscr{D} \longrightarrow \mathscr{A}$

$\alpha \triangleq \lambda X \cdot \{\hbar(x) \mid x \in X\}$

$\gamma \triangleq \lambda Y \cdot \{x \in \mathscr{D} \mid \hbar(x) \in Y\}$

$\Longrightarrow \langle \wp(\mathscr{D}), \subseteq \rangle \underset{\alpha}{\overset{\gamma}{\leftrightarrows}} \langle \wp(\mathscr{A}), \subseteq \rangle$ ($\longrightarrow\!\!\!\!\rightarrow$ iff $\hbar$ onto)

Example [(*)]: rule of signs: $A = \mathbb{Z}$, $B = \{-1, 0, 1\}$, $\hbar(z) = z/|z|$

Counter-example [(**)]: intervals (octagons, polyhedra, etc)

(*) Patrick Cousot, Radhia Cousot: Systematic Design of Program Analysis Frameworks. POPL 1979: 269-282

(**) Patrick Cousot, Radhia Cousot: Abstract Interpretation: A Unified Lattice Model for Static Analysis of Programs by Construction or Approximation of Fixpoints. POPL 1977: 238-252

Properties of Galois connections

α preserves existing lubs (by order-duality, γ preserves existing glbs)

One adjoint uniquely determine the other

α is surjective (iff γ injective iff $\alpha \circ \gamma = 1$), written

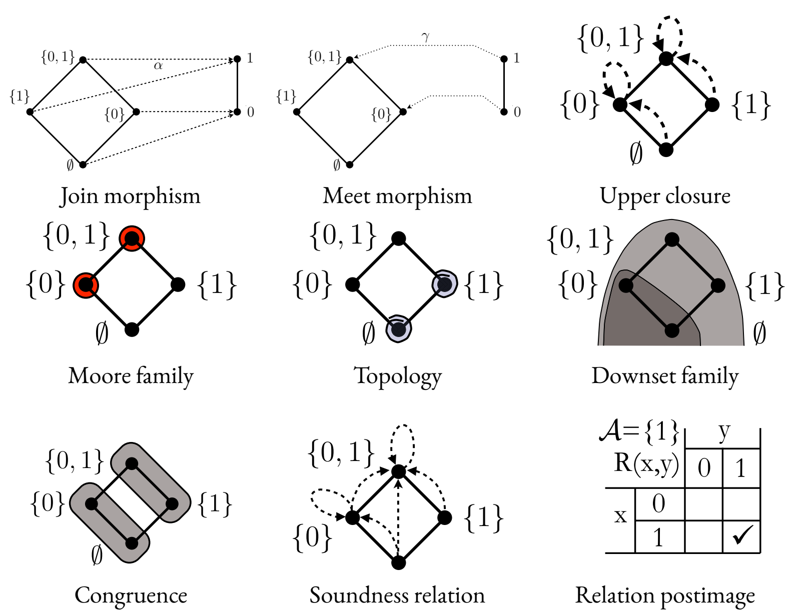
$$\langle P, \preccurlyeq \rangle \underset{\alpha}{\overset{\gamma}{\leftarrows\!\!\!\!\twoheadrightarrow}} \langle Q, \sqsubseteq \rangle$$

The composition of Galois connections is a Galois connection

$\alpha(x)$ is the best over-approximation of $x \in P$:

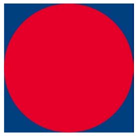
- $x \preccurlyeq \gamma(\alpha(x))$ over-approximation
- $x \preccurlyeq \gamma(y) \Longrightarrow \alpha(x) \sqsubseteq y$ more precise than any other over-approximation

Equivalent mathematical structures

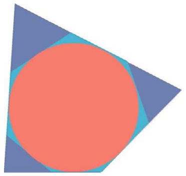


In absence of best abstraction?

Best abstraction of a disk by a rectangular parallelogram (intervals)



No best abstraction of a disk by a polyhedron (Euclid)



use only abstraction or concretization or widening [(*)]

(*) Patrick Cousot, Radhia Cousot: Abstract Interpretation Frameworks. J. Log. Comput. 2(4): 511-547 (1992)

Sound semantics abstraction

program $\mathtt{P} \in \mathbb{L}$ programming language

standard semantics $S[\![\mathtt{P}]\!] \in \mathscr{D}$ semantic domain

collecting semantics $\{S[\![\mathtt{P}]\!]\} \in \wp(\mathscr{D})$ semantic property

abstract semantics $\overline{S}[\![\mathtt{P}]\!] \in \mathscr{A}$ abstract domain

concretization $\gamma \in \mathscr{A} \longrightarrow \wp(\mathscr{D})$

soundness $\{S[\![\mathtt{P}]\!]\} \subseteq \gamma(\overline{S}[\![\mathtt{P}]\!])$

i.e. $S[\![\mathtt{P}]\!] \in \gamma(\overline{S}[\![\mathtt{P}]\!])$, $\mathtt{P}$ *has abstract property* $\overline{S}[\![\mathtt{P}]\!]$

Best abstract semantics

If $\langle \wp(\mathscr{D}), \subseteq \rangle \xrightleftharpoons[\alpha]{\gamma} \langle \mathscr{A}, \sqsubseteq \rangle$ then the best abstract semantics is the abstraction of the collecting semantics

$$\overline{S}[\![P]\!] \triangleq \alpha(\{S[\![P]\!]\})$$

Proof:

- It is *sound*: $\overline{S}[\![P]\!] \triangleq \alpha(\{S[\![P]\!]\}) \sqsubseteq \overline{S}[\![P]\!] \Longrightarrow \{S[\![P]\!]\} \subseteq \gamma(\overline{S}[\![P]\!]) \Longrightarrow S[\![P]\!] \in \gamma(\overline{S}[\![P]\!])$
- It is the *most precise*: $S[\![P]\!] \in \gamma(\overline{\overline{S}}[\![P]\!]) \Longrightarrow \{S[\![P]\!]\} \subseteq \gamma(\overline{\overline{S}}[\![P]\!]) \Longrightarrow \overline{S}[\![P]\!] \triangleq \alpha(\{S[\![P]\!]\}) \sqsubseteq \overline{\overline{S}}[\![P]\!]$ ■

Calculational design of the abstract semantics

The (standard hence collecting) semantics are defined by composition of mathematical structures (such as set unions, products, functions, fixpoints, *etc*)

If you know best abstractions of properties, you also know best abstractions of these mathematical structures

So, by composition, you also know the best abstraction of the collecting semantics ⇝ calculational design of the abstract semantics

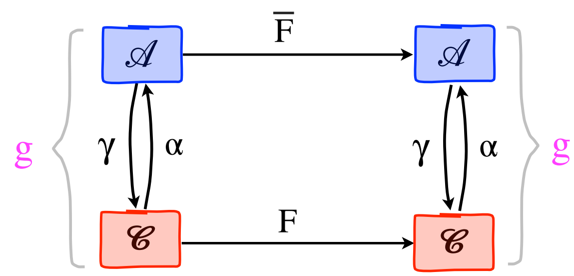
Orthogonally, there are many styles of

- *semantics* (traces, relations, transformers,…)
- *induction* (transitional, structural, segmentation [POPL 2012])
- *presentations* (fixpoints, equations, constraints, rules [CAV 1995])

Example: functional connector

If $g = \langle \mathscr{C}, \subseteq \rangle \xrightleftharpoons[\alpha]{\gamma} \langle \mathscr{A}, \sqsubseteq \rangle$ then

$$g \Longmapsto g = \langle \mathscr{C} \overset{\nearrow}{\longrightarrow} \mathscr{C}, \subseteq \rangle \xrightleftharpoons[\lambda F.\alpha \circ F \circ \gamma]{\lambda \overline{F}.\gamma \circ \overline{F} \circ \alpha} \langle \mathscr{A} \overset{\nearrow}{\longrightarrow} \mathscr{A}, \sqsubseteq \rangle$$



($\Longmapsto$ is a called a *Galois connector*)

Fixpoint abstraction

Best abstraction (completeness case)

if $\alpha \circ F = \overline{F} \circ \alpha$ then $\overline{F} = \alpha \circ F \circ \gamma$ and $\alpha(\text{lfp } F) = \text{lfp } \overline{F}$

e.g. semantics, proof methods, static analysis of finite state systems

Best approximation (incompleteness case)

if $\overline{F} = \alpha \circ F \circ \gamma$ but $\alpha \circ F \sqsubseteq \overline{F} \circ \alpha$ then $\alpha(\text{lfp } F) \sqsubseteq \text{lfp } \overline{F}$

e.g. static analysis of infinite state systems

idem for equations, constraints, rule-based deductive systems, *etc*

Fixpoint abstraction

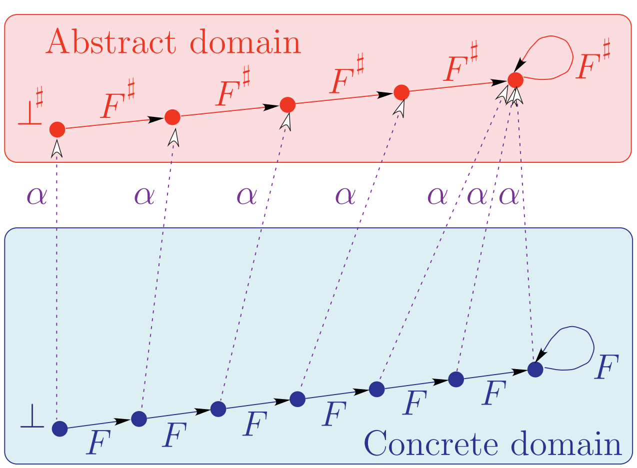
Theorem 1 If $\langle C, \sqsubseteq\rangle \underset{\alpha}{\overset{\gamma}{\leftrightarrows}} \langle A, \preceq\rangle$ in cpos for infinite/transfinite chains, $F \in C \mapsto C$ and $G \in A \mapsto A$ are continuous/increasing then

$$\begin{array}{rcll} \alpha(\mathbf{lfp}^{\sqsubseteq} F) & = & \mathbf{lfp}^{\preceq} G & \Longleftarrow \quad \alpha \circ F = G \circ \alpha \quad \text{(commutation condition)} \\ G & = & \alpha \circ F \circ \gamma & \\ \alpha(\mathbf{lfp}^{\sqsubseteq} F) & \preceq & \mathbf{lfp}^{\preceq} G & \Longleftarrow \quad \alpha \circ F \mathrel{\dot{\preceq}} G \circ \alpha \quad \text{(semi-commutation condition)} \end{array}$$

[Cousot and Cousot, 1979b, theorem 7.1.0.4(2–3)], see also [de Bakker et al., 1984, lemma 4.3], [Apt and Plotkin, 1986, fact 2.3], [Backhouse, 2000, theorem 95], etc.

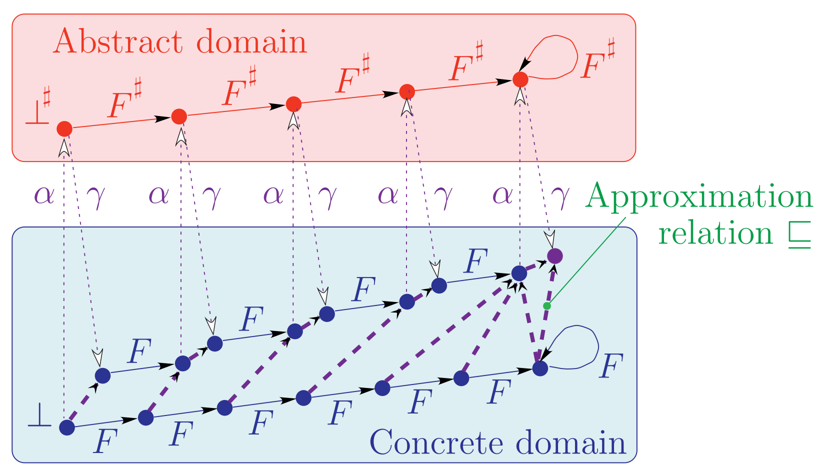
[Cousot and Cousot, 1979b] Patrick Cousot, Radhia Cousot: Systematic Design of Program Analysis Frameworks. POPL 1979: 269-282

Exact fixpoint abstraction



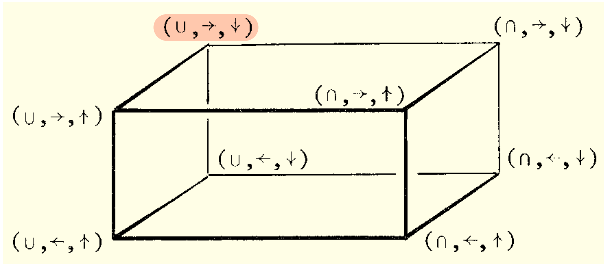
$\alpha \circ F = F^\sharp \circ \alpha \;\Rightarrow\; \alpha(\mathit{lfp}\, F) = \mathit{lfp}\, F^\sharp$

Approximate fixpoint abstraction



$\mathit{lfp}\, F \sqsubseteq \gamma(\mathit{lfp}\, F^\sharp)$

Duality



Order duality: join (∪) or meet (∩)

Inversion duality: forward (→) or backward ($\leftarrow = (\rightarrow)^{-1}$)

Fixpoint duality: least (↓) or greatest (↑)

Patrick Cousot, Radhia Cousot: **Abstract Interpretation: A Unified Lattice Model for Static Analysis of Programs by Construction or Approximation of Fixpoints.** POPL 1977: 238-252

Why abstracting properties of semantics, not semantics?

1. Abstract interpretation = a non-standard semantics (computations on values in the standard semantics are replaced by computations on abstract values) $\Longrightarrow$ extremely limited
2. Abstract interpretation = an abstraction of the standard semantics $\Longrightarrow$ limited
3. Abstract interpretation = an abstraction of properties of the standard semantics $\Longrightarrow$ more

i.e. (1) is an abstraction of (2), (2) is an abstraction of (3)

Example: trace semantics properties

Domain of [in]finite traces on states: Π

"Standard" trace semantics domain: $\mathcal{D} = \wp(\Pi)$

"Standard" trace semantics $S[\![\mathrm{P}]\!] \in \mathcal{D} = \wp(\Pi)$

Domain of semantics properties is $\wp(\mathcal{D}) = \wp(\wp(\Pi))$

Collecting semantics $C[\![\mathrm{P}]\!] \triangleq \{S[\![\mathrm{P}]\!]\} \in \wp(\mathcal{D}) = \wp(\wp(\Pi))$

How to abstract the standard semantics?

The join abstraction:

$$\langle \wp(\wp(\Pi)), \subseteq\rangle \underset{\alpha_\cup}{\overset{\gamma_\cup}{\leftrightarrows}} \langle \wp(\Pi), \subseteq\rangle$$

$$\alpha_\cup(X) \triangleq \bigcup X$$

$$\gamma_\cup(Y) \triangleq \wp(Y)$$

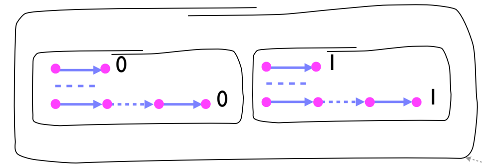
Join abstraction of the collecting semantics:

$$\alpha_\cup(C[\![\mathtt{P}]\!]) \triangleq \bigcup\{S[\![\mathtt{P}]\!]\} \triangleq S[\![\mathtt{P}]\!]$$

(*i.e.* the semantics is the join abstraction of its strongest property)

Loss of information

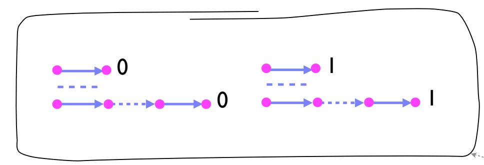
"Always terminate with the same value, either 0 or 1"

$P =$ 

$P \in \wp(\wp(\Pi))$

always the same result

Join abstraction:

$\alpha_\cup(P) =$ 

$\alpha_\cup(P) \in \wp(\Pi)$

results can be different

"Always terminate, either with 0 or 1"

Limitations of the union abstraction

Complete iff any property of the semantics $S[\![\mathtt{P}]\!]$ is also valid for any subset $\gamma(S[\![\mathtt{P}]\!]) = \wp(S[\![\mathtt{P}]\!])$:

- Examples: safety, liveness
- Counter-example: security (e.g. authentification using a random cryptographic nonce)

Exact abstractions

Exact abstractions

The concrete properties of the standard semantics $S[\![\mathtt{P}]\!]$ that you want to prove can always be proved in the abstract (which is simpler):

$$\forall Q \in \mathscr{A} : S[\![\mathtt{P}]\!] \in \gamma(Q) \iff S[\![\overline{\mathtt{P}}]\!] \sqsubseteq Q$$

where

$$S[\![\overline{\mathtt{P}}]\!] \triangleq \alpha \circ S[\![\mathtt{P}]\!] \circ \gamma$$

Example 1 of exact abstraction: grammars

Patrick Cousot, Radhia Cousot: Grammar semantics, analysis and parsing by abstract interpretation. Theor. Comput. Sci. 412(44): 6135-6192 (2011)

Example: Grammars

Context-free grammar on alphabet $A = Num \cup Var \cup \{+,-,(,),\ldots\}$:

$$E \ ::= \ Num \mid Var \mid E + E \mid -E \mid (E)$$

Chomsky-Schützenberger fixpoint semantics:

$$\mathcal{S}[\![E]\!] = \mathbf{lfp}^{\subseteq} \mathcal{F}[\![E]\!]$$

$$\begin{aligned}\mathcal{F}[\![E]\!]X &\triangleq \mathcal{S}[\![Num]\!] \cup \mathcal{S}[\![Var]\!] \\ &\cup \{e_1 + e_2 \mid e_1, e_2 \in X\} \\ &\cup \{-e \mid e \in X\} \cup \{(e) \mid e \in X\}\end{aligned}$$

Example: Grammars (cont'd)

FIRST abstraction of a language $X \in A^*$:

$$\alpha_F(X) \triangleq \{\ell \mid \exists \sigma \in A^* : \ell\sigma \in X\} \cup \{\epsilon \mid \epsilon \in X\}$$

Galois connection:

$$\langle \wp(A^*), \subseteq \rangle \xrightleftharpoons[\alpha_F]{\gamma_F} \langle \wp(A \cup \{\epsilon\}), \subseteq \rangle$$

where

$$\gamma_F(Y) \triangleq \{\ell\sigma \mid \ell \in Y \wedge \sigma \in A^*\} \cup \{\epsilon \mid \epsilon \in Y\}$$

Example: Grammars (cont'd)

Commutation:

$$\alpha_F \circ \mathcal{F}[\![E]\!] \quad = \quad \overline{\mathcal{F}}[\![E]\!] \circ \alpha_F$$

where for $\quad E \quad ::= \quad \mathit{Num} \mid \mathit{Var} \mid E + E \mid -E \mid (E)$

$$\overline{\mathcal{F}}[\![E]\!]Y \quad \triangleq \quad \mathcal{S}[\![\mathit{Num}]\!] \cup \mathcal{S}[\![\mathit{Var}]\!] \cup (Y \setminus \{\epsilon\}) \cup \{+ \mid \epsilon \in Y\} \cup \{-, (\}$$

FIRST abstract semantics:

$$\begin{aligned} \overline{\mathcal{S}}[\![E]\!] &\triangleq \alpha_F(\mathcal{S}[\![E]\!]) \\ &= \alpha_F(\mathbf{lfp}^{\subseteq}\, \mathcal{F}[\![E]\!]) && \text{(Chomsky-Schützenberger)} \\ &= \mathbf{lfp}^{\subseteq}\, \overline{\mathcal{F}}[\![E]\!] && \text{(fixpoint abstraction th.)} \end{aligned}$$

Machine-checkable calculational design

$$\begin{aligned}
& \alpha_F \circ \mathcal{F}[\![E]\!] \\
= \;& \boldsymbol{\lambda} X \bullet \alpha_F(\mathcal{F}[\![E]\!](X)) && \text{(def. } \circ\text{)} \\
= \;& \boldsymbol{\lambda} X \bullet \{\ell \mid \exists \sigma \in A^* : \ell\sigma \in \mathcal{F}[\![E]\!](X)\} \cup \{\epsilon \mid \epsilon \in \mathcal{F}[\![E]\!](X)\} && \text{(def. } \alpha_F\text{)} \\
= \;& \boldsymbol{\lambda} X \bullet \{\ell \mid \exists \sigma \in A^* : \ell\sigma \in \mathcal{F}[\![E]\!](X)\} && \text{(since. } \forall X : \epsilon \notin \mathcal{F}[\![E]\!](X)\text{)} \\
= \;& \boldsymbol{\lambda} X \bullet \{\ell \mid \exists \sigma \in A^* : \ell\sigma \in \mathcal{S}[\![\mathit{Num}]\!] \cup \mathcal{S}[\![\mathit{Var}]\!] \cup \{e_1 + e_2 \mid e_1, e_2 \in X\} \cup \{-e \mid e \in X\} \cup \{(e) \mid e \in X\}\} \\
& \quad \text{(def. } \mathcal{F}[\![E]\!]X \triangleq \mathcal{S}[\![\mathit{Num}]\!] \cup \mathcal{S}[\![\mathit{Var}]\!] \cup \{e_1 + e_2 \mid e_1, e_2 \in X\} \cup \{-e \mid e \in X\} \cup \{(e) \mid e \in X\}\text{)} \\
= \;& \boldsymbol{\lambda} X \bullet \mathcal{S}[\![\mathit{Num}]\!] \cup \mathcal{S}[\![\mathit{Var}]\!] \cup \{\ell \mid \exists \sigma \in A^* : \ell\sigma \in X\} \cup \{+ \mid \epsilon \in X\} \cup \{-\} \cup \{(\} \\
& && \text{(def. } \in \text{ and } \epsilon + e_2 = +e_2\text{)} \\
= \;& \boldsymbol{\lambda} X \bullet \mathcal{S}[\![\mathit{Num}]\!] \cup \mathcal{S}[\![\mathit{Var}]\!] \cup (\alpha_F(X) \setminus \{\epsilon\}) \cup \{+ \mid \epsilon \in \alpha_F(X)\} \cup \{-\} \cup \{(\} \\
& && \text{(def. } \alpha_F \text{ and } \epsilon \in X \iff \epsilon \in \alpha_F(X)\text{)} \\
= \;& \boldsymbol{\lambda} X \bullet \overline{\mathcal{F}}[\![E]\!](\alpha_F(X)) \\
& \quad \text{(by defining } \overline{\mathcal{F}}[\![E]\!]Y \triangleq \mathcal{S}[\![\mathit{Num}]\!] \cup \mathcal{S}[\![\mathit{Var}]\!] \cup (Y \setminus \{\epsilon\}) \cup \{+ \mid \epsilon \in Y\} \cup \{-, (\}\text{)} \\
= \;& \overline{\mathcal{F}}[\![E]\!] \circ \alpha_F && \text{(def. } \circ\text{)}
\end{aligned}$$

Algorithm

Read the grammar G, establish the system of equations $Y = \overline{\mathcal{F}}[\![G]\!](Y)$, solve by chaotic iterations

This is, up to [en]coding details, the classical algorithm:

```
for each α ∈ (T ∪ ε)
    FIRST(α) ← α
for each A ∈ NT
    FIRST(A) ← ∅
while (FIRST sets are still changing)
    for each p ∈ P, where p has the form A→β
        if β is β1β2 ... βk, where βi ∈ T ∪ NT, then
            FIRST(A) ← FIRST(A) ∪ (FIRST(β1) − {ε})
            i ← 1
            while (ε ∈ FIRST(βi) and i ≤ k-1)
                FIRST(A) ← FIRST(A) ∪ (FIRST(βi+1) − {ε})
                i ← i + 1
        if i = k and ε ∈ FIRST(βk)
            then FIRST(A) ← FIRST(A) ∪ {ε}
```

Keith D. Cooper, Linda Torczon: Engineering a Compiler. Morgan Kaufmann 2004

Hierarchies of abstractions

Comparison of abstractions

$$\langle P, \preccurlyeq \rangle \xleftrightarrow[\alpha_1]{\gamma_1} \langle Q, \sqsubseteq \rangle$$

is more precise than

$$\langle P, \preccurlyeq \rangle \xleftrightarrow[\alpha_2]{\gamma_2} \langle R, \lesssim \rangle$$

iff $\gamma_2(R) \subseteq \gamma_1(Q)$

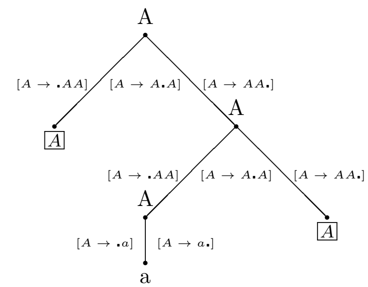
(every abstraction in R is exactly expressible by Q)

We say that Q is a refinement of R and R that is a abstraction of Q

A pre-order

Hierarchy of Grammar Semantics

Chomsky–Schützenberger terminal language



Example of proto-derivation tree

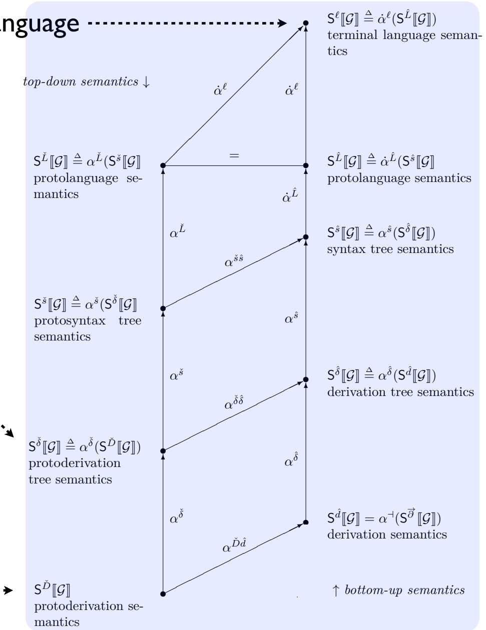
$$\vdash \xrightarrow{\boxed{A}} \dashv$$

$$\Rrightarrow_{\mathcal{G}} \vdash \xrightarrow{(\!|A} \dashv[A \to \bullet AA] \xrightarrow{\boxed{A}} \dashv[A \to A\bullet A] \xrightarrow{\boxed{A}} \dashv[A \to AA\bullet] \xrightarrow{A|\!)} \dashv$$

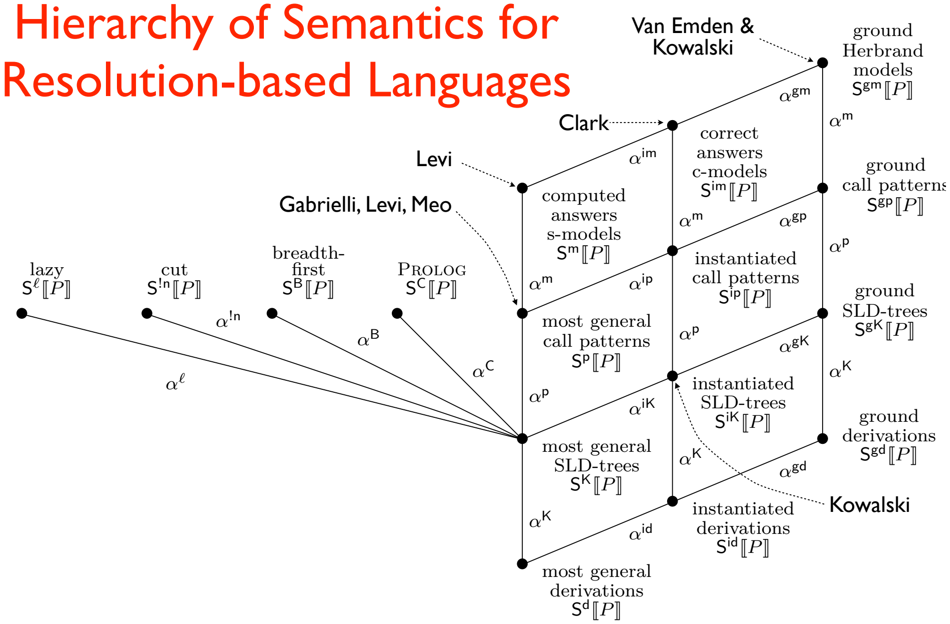
$$\Rrightarrow_{\mathcal{G}} \vdash \xrightarrow{(\!|A} \dashv[A \to \bullet AA] \xrightarrow{\boxed{A}} \dashv[A \to A\bullet A] \xrightarrow{(\!|A} \dashv[A \to AA\bullet][A \to \bullet a] \xrightarrow{a} \dashv[A \to AA\bullet][A \to a\bullet] \xrightarrow{A|\!)} \dashv[A \to AA\bullet] \xrightarrow{A|\!)} \dashv$$

$$\Rrightarrow_{\mathcal{G}} \vdash \xrightarrow{(\!|A} \dashv[A \to \bullet AA] \xrightarrow{(\!|A} \dashv[A \to A\bullet A][A \to \bullet a] \xrightarrow{a} \dashv[A \to A\bullet A][A \to a\bullet] \xrightarrow{A|\!)} \dashv[A \to A\bullet A] \xrightarrow{(\!|A} \dashv[A \to AA\bullet][A \to \bullet a] \xrightarrow{a} \dashv[A \to AA\bullet][A \to a\bullet] \xrightarrow{A|\!)} \dashv[A \to AA\bullet] \xrightarrow{A|\!)} \dashv . \qquad \square$$

Example of proto-derivation



Hierarchy of Semantics for Resolution-based Languages



Patrick Cousot, Radhia Cousot, Roberto Giacobazzi: Abstract interpretation of resolution-based semantics. Theor. Comput. Sci. 410(46): 4724-4746 (2009)

Example II of exact abstraction: graphs

Ilya Sergey, Jan Midtgaard, Dave Clarke: Calculating Graph Algorithms for Dominance and Shortest Path. MPC 2012: 132-156

Transition system

Transition system: $\langle \Sigma,\ \mathbb{A},\ \longrightarrow \rangle$

transition relation: $\longrightarrow\ \in \wp(\Sigma \times \mathbb{A} \times \Sigma)$

transitions/edges: $\sigma \xrightarrow{\mathrm{A}} \sigma'$

Example: non-negatively weighted graphs $\mathbb{A} \triangleq \mathbb{N}$

Finite paths

Finite paths:

$$\begin{aligned}\Theta^+ \ \triangleq\ \{\sigma_0 \xrightarrow{\mathrm{A}_0} \sigma_1 \ldots \sigma_{n-1} \xrightarrow{\mathrm{A}_{n-1}} \sigma_n \mid n \geqslant 0 \wedge \\ \forall i \in [0,n] : \sigma_i \in \Sigma \wedge \forall i \in [0,n) : \mathrm{A}_i \in \mathbb{A}\}\end{aligned}$$

Paths between two vertices:

$$\begin{aligned}\Pi &\in (\Sigma \times \Sigma) \mapsto \wp(\Theta^+) \\ \Pi(\sigma,\sigma') &\triangleq \{\sigma_0 \xrightarrow{\mathrm{A}_0} \sigma_1 \ldots \sigma_{n-1} \xrightarrow{\mathrm{A}_{n-1}} \sigma_n \mid \sigma = \sigma_0 \wedge n \geqslant 0 \wedge \\ &\qquad \forall i \in [0,n-1] : \sigma_i \xrightarrow{\mathrm{A}_i} \sigma_{i+1} \wedge \sigma_n = \sigma'\}\end{aligned}$$

destination

departure

Fixpoint characterization

Pointwise fixpoint charaterization:

$$\begin{aligned}\Pi &= \mathbf{lfp}^{\dot{\subseteq}}\, F \\ F &\in ((\Sigma \times \Sigma) \mapsto \wp(\Theta^+)) \mapsto ((\Sigma \times \Sigma) \mapsto \wp(\Theta^+)) \\ F(X)(\sigma,\sigma') &= \llbracket \sigma = \sigma' \ ?\ \{\sigma\} \ {\large\circ\!\!\circ}\ \bigcup_{\sigma'' \in \Sigma} \{\sigma \xrightarrow{\mathrm{A}} \sigma''\pi \mid \sigma \xrightarrow{\mathrm{A}} \sigma'' \wedge \sigma''\pi \in X(\sigma'',\sigma')\} \rrbracket\end{aligned}$$

(a path of n transitions is either a single vertex ($n = 0$) or an edge followed by a path of n - 1 transitions)

Minimal path length abstraction

Edges have non-negative lengths $\mathbb{A} = \mathbb{N}$

Abstraction:

$$\begin{aligned}\alpha &\in \Theta^+ \mapsto \mathbb{N} \\ \alpha(\sigma) &\triangleq 0 \\ \alpha(\sigma \xrightarrow{n} \sigma'\pi) &\triangleq n + \alpha(\sigma'\pi) \\ \alpha &\in \wp(\Theta^+) \mapsto \mathbb{N}^\infty \\ \alpha(X) &\triangleq \min\{\alpha(\pi) \mid \pi \in X\}\end{aligned}$$

where

$\min \emptyset = +\infty$

$\mathbb{N}^\infty \triangleq \mathbb{N} \cup \{+\infty\}$

$\langle \mathbb{N}^\infty,\ \geqslant,\ \min \rangle$ is a complete lattice

Galois connection

$$\langle \wp(\Theta^+),\ \subseteq\rangle \ \underset{\alpha}{\overset{\gamma}{\leftrightarrows}}\ \langle \mathbb{N}^\infty,\ \geqslant\rangle$$

Pointwise extension:

$$\begin{array}{rcl}\dot\alpha & \in & (\Sigma\times\Sigma \mapsto \wp(\Theta^+)) \mapsto (\Sigma\times\Sigma \mapsto \mathbb{N}^\infty)\\ \dot\alpha(X)(\sigma,\sigma') & \triangleq & \alpha(X(\sigma,\sigma'))\end{array}$$

Pointwise Galois connection:

$$\langle(\Sigma\times\Sigma)\mapsto\wp(\Theta^+),\ \dot\subseteq\rangle \ \underset{\dot\alpha}{\overset{\dot\gamma}{\leftrightarrows}}\ \langle(\Sigma\times\Sigma)\mapsto\mathbb{N}^\infty,\ \dot\geqslant\rangle$$

Shortest distance

Shortest distance $\Delta(\sigma,\sigma')$ between any two vertices

$$\begin{array}{rcl}\Delta & \in & (\Sigma\times\Sigma)\mapsto\mathbb{N}^\infty\\ \Delta & \triangleq & \dot\alpha(\Pi) \ = \ \dot\alpha(\mathbf{lfp}^{\dot\subseteq} F)\end{array}$$

Calculational design of the shortest distance algorithm

$\dot\alpha\circ F$

$= \boldsymbol{\lambda} X\bullet\dot\alpha(F(X))$ $\wr$def. $\circ\wr$

$= \boldsymbol{\lambda}(\sigma,\sigma')\bullet\boldsymbol{\lambda} X\bullet\dot\alpha(F(X))(\sigma,\sigma')$ $\wr$def. $\boldsymbol{\lambda} x\bullet e\wr$

$= \boldsymbol{\lambda}(\sigma,\sigma')\bullet\boldsymbol{\lambda} X\bullet\dot\alpha(\boldsymbol{\lambda}(\sigma,\sigma')\bullet\llbracket\sigma=\sigma' \mathrel{?} \{\sigma\} \mathrel{:} \bigcup_{\sigma''\in\Sigma}\{\sigma\xrightarrow{n}\sigma''\pi \mid \sigma\xrightarrow{n}\sigma''\wedge\sigma''\pi\in X(\sigma'',\sigma')\}\rrbracket)(\sigma,\sigma')$ $\wr$def. $F\wr$

$= \boldsymbol{\lambda}(\sigma,\sigma')\bullet\boldsymbol{\lambda} X\bullet\alpha(\llbracket\sigma=\sigma' \mathrel{?} \{\sigma\} \mathrel{:} \bigcup_{\sigma''\in\Sigma}\{\sigma\xrightarrow{n}\sigma''\pi \mid \sigma\xrightarrow{n}\sigma''\wedge\sigma''\pi\in X(\sigma'',\sigma')\}\rrbracket)$ $\wr$def. $\dot\alpha(X)(\sigma,\sigma')\triangleq\alpha(\Delta(\sigma,\sigma'))\wr$

$= \boldsymbol{\lambda}(\sigma,\sigma')\bullet\boldsymbol{\lambda} X\bullet\llbracket\sigma=\sigma' \mathrel{?} \alpha(\{\sigma\}) \mathrel{:} \alpha(\bigcup_{\sigma''\in\Sigma}\{\sigma\xrightarrow{n}\sigma''\pi \mid \sigma\xrightarrow{n}\sigma''\wedge\sigma''\pi\in X(\sigma'',\sigma')\})\rrbracket$ $\wr$def. conditional $\llbracket\ldots\mathrel{?}\ldots\mathrel{:}\ldots\rrbracket\wr$

$= \boldsymbol{\lambda}(\sigma,\sigma')\bullet\boldsymbol{\lambda} X\bullet\llbracket\sigma=\sigma' \mathrel{?} \alpha(\{\sigma\}) \mathrel{:} \min_{\sigma''\in\Sigma}\alpha(\{\sigma\xrightarrow{n}\sigma''\pi \mid \sigma\xrightarrow{n}\sigma''\wedge\sigma''\pi\in X(\sigma'',\sigma')\})\rrbracket$ $\wr$join preservation in Galois C.$\wr$

$= \boldsymbol{\lambda}(\sigma,\sigma')\bullet\boldsymbol{\lambda} X\bullet\llbracket\sigma=\sigma' \mathrel{?} \min\{\alpha(\pi)\mid\pi\in\{\sigma\}\} \mathrel{:} \min_{\sigma''\in\Sigma}\min\{\alpha(\pi)\mid\pi\in\{\sigma\xrightarrow{n}\sigma''\pi \mid \sigma\xrightarrow{n}\sigma''\wedge\sigma''\pi\in X(\sigma'',\sigma')\}\}\rrbracket$ $\wr$def. $\alpha(X)\triangleq\min\{\alpha(\pi)\mid\pi\in X\}\wr$

Calculational design of the shortest distance algorithm

$= \boldsymbol{\lambda}(\sigma,\sigma')\bullet\boldsymbol{\lambda} X\bullet\llbracket\sigma=\sigma' \mathrel{?} \min\{\alpha(\sigma)\} \mathrel{:} \min_{\sigma''\in\Sigma}\min\{\alpha(\sigma\xrightarrow{n}\sigma''\pi)\mid\sigma\xrightarrow{n}\sigma''\wedge\sigma''\pi\in X(\sigma'',\sigma')\}\rrbracket$ $\wr$def. $\in\wr$

$= \boldsymbol{\lambda}(\sigma,\sigma')\bullet\boldsymbol{\lambda} X\bullet\llbracket\sigma=\sigma' \mathrel{?} \min\{0\} \mathrel{:} \min_{\sigma''\in\Sigma}\min\{n+\alpha(\sigma''\pi)\mid\sigma\xrightarrow{n}\sigma''\wedge\sigma''\pi\in X(\sigma'',\sigma')\}\rrbracket$ $\wr$def. $\alpha(\sigma)\triangleq 0$ and $\alpha(\sigma\xrightarrow{n}\sigma'\pi)\triangleq n+\alpha(\sigma'\pi)\wr$

$= \boldsymbol{\lambda}(\sigma,\sigma')\bullet\boldsymbol{\lambda} X\bullet\llbracket\sigma=\sigma' \mathrel{?} 0 \mathrel{:} \min_{\sigma''\in\Sigma}\{n+\min\{\alpha(\sigma''\pi)\mid\sigma''\pi\in X(\sigma'',\sigma')\}\mid\sigma\xrightarrow{n}\sigma''\}\rrbracket$ $\wr$def. $\min\wr$

$= \boldsymbol{\lambda}(\sigma,\sigma')\bullet\boldsymbol{\lambda} X\bullet\llbracket\sigma=\sigma' \mathrel{?} 0 \mathrel{:} \min_{\sigma''\in\Sigma}\{n+\dot\alpha(X)(\sigma'',\sigma')\mid\sigma\xrightarrow{n}\sigma''\}\rrbracket$

$\wr$def. $\dot\alpha(X)(\sigma'',\sigma') \triangleq \alpha(X(\sigma'',\sigma')) = \min\{\alpha(\pi)\mid\pi\in X(\sigma'',\sigma')\} = \min\{\alpha(\sigma''\pi)\mid\sigma''\pi\in X(\sigma'',\sigma')\}$ where $\pi=\sigma''\pi'$ and π' can be empty $\wr$

$= \boldsymbol{\lambda}(\sigma,\sigma')\bullet\boldsymbol{\lambda} X\bullet G(\dot\alpha(X))(\sigma,\sigma')$

by defining

$$G(X)(\sigma,\sigma') \ = \ \llbracket\sigma=\sigma' \mathrel{?} 0 \mathrel{:} \min_{\sigma''\in\Sigma}\{n+X(\sigma'',\sigma')\mid\sigma\xrightarrow{n}\sigma''\}\rrbracket$$

$= \boldsymbol{\lambda} X\bullet G(\dot\alpha(X))$ $\wr$def. $\boldsymbol{\lambda} x\bullet e\wr$

$= G\circ\dot\alpha$ $\wr$def. $\circ\wr$

Shortest distance in fixpoint form

By the fixpoint abstraction theorem

$$\begin{aligned}\Delta &= \alpha(\mathbf{lfp}^{\dot\subseteq}\, F)\\ &= \mathbf{lfp}^{\dot\geqslant}\, G\\ &= \min_{n\in\mathbb{N}} G^n(\boldsymbol{\lambda}\,(\sigma,\sigma')\bullet\,+\infty)\end{aligned}$$

where the iterates are

1. $G^0(X) = X$

   $G^{n+1} = G \circ G^n,\ n \in \mathbb{N}$

Shortest distance algorithm

```
forall σ ∈ Σ do
  forall σ′ ∈ Σ do
    Δ(σ, σ′) := if σ = σ′ then 0 else +∞;
repeat
  change := false;
  forall σ ∈ Σ do
    forall σ′ ∈ Σ do
      forall σ″ ∈ Σ do
        if (σ ≠ σ′ ∧ σ -n-> σ″ ∧ Δ(σ, σ′) > n + Δ(σ″, σ′)) then
          { Δ(σ, σ′) := n + Δ(σ″, σ′);
            change := true }
until ¬change;
```

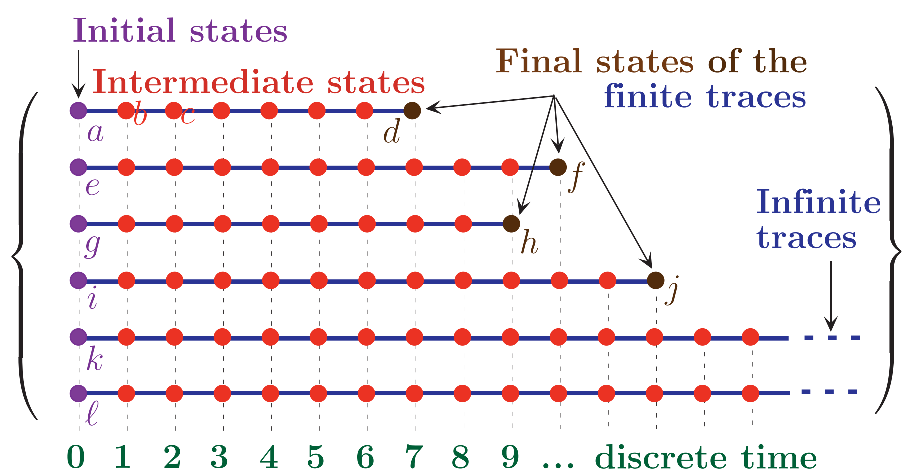
Not Floyd-Warshall? Take instead:

$$\begin{aligned}\alpha(\sigma) &\triangleq 0\\ \alpha(\sigma \xrightarrow{n} \sigma') &\triangleq n\\ \alpha(\pi\sigma\pi') &\triangleq \alpha(\pi\sigma) + \alpha(\sigma\pi')\end{aligned}$$

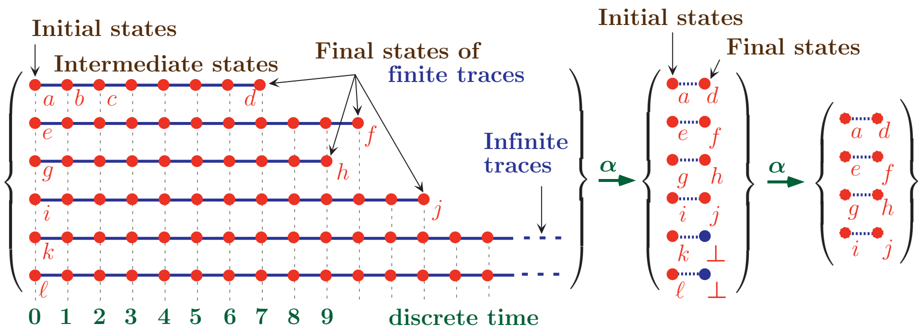
Example III of exact abstractions: semantics

Patrick Cousot: Constructive design of a hierarchy of semantics of a transition system by abstract interpretation. Theor. Comput. Sci. 277(1-2): 47-103 (2002)

Trace semantics



Abstraction to denotational/natural semantics

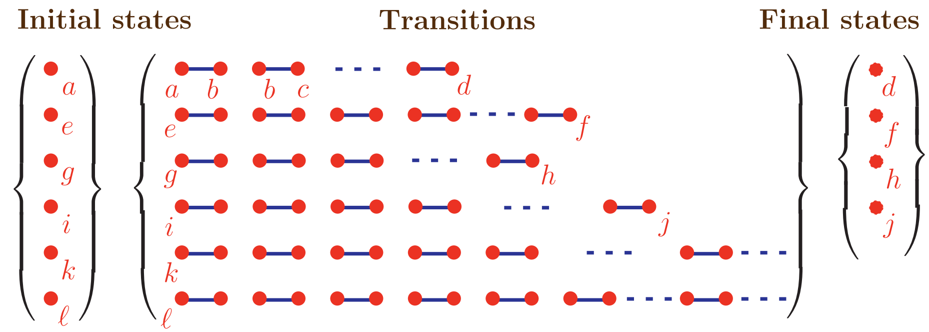


Trace semantics

Denotational semantics

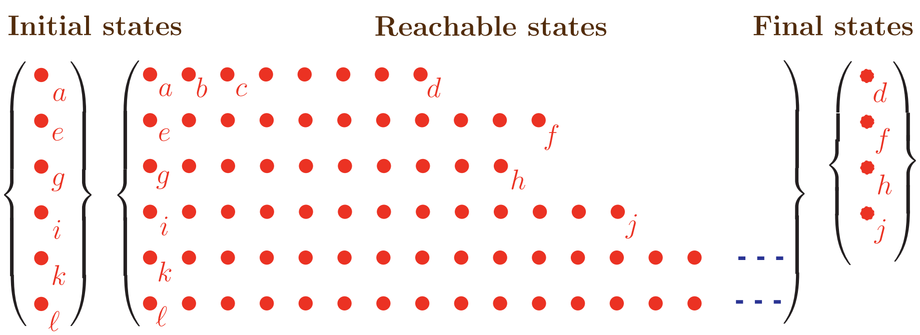
Natural semantics

Abstraction to small-steps operational semantics



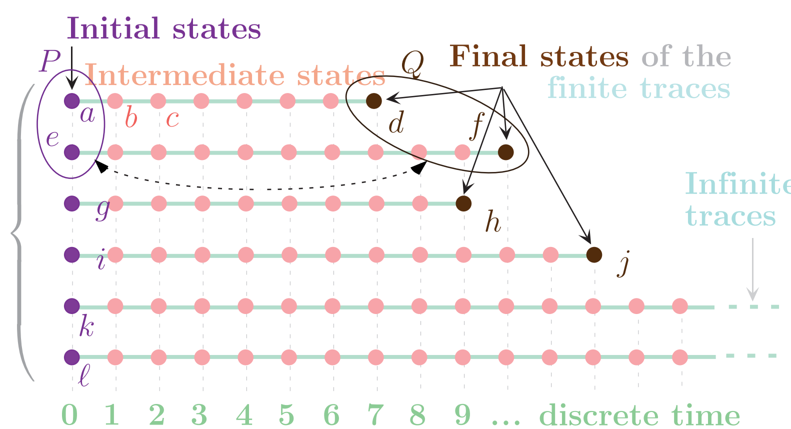
(Small-Step) Operational Semantics

Abstraction to reachability/invariance



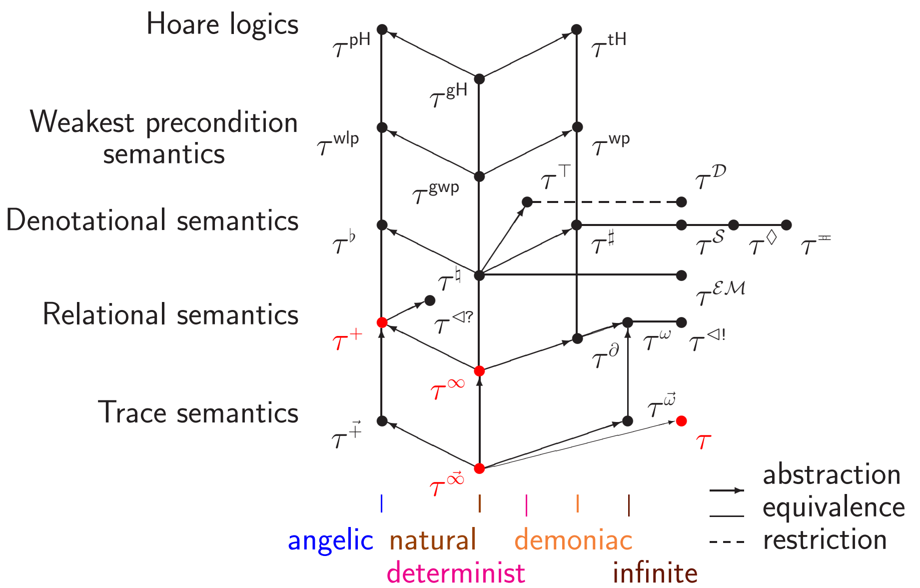
Partial Correctness / Invariance Semantics

Abstraction to Hoare logic



$$\{P\}C\{Q\} \Leftrightarrow \{\bullet \mid \bullet \in P \wedge \bullet\!-\!\bullet\!-\!\bullet\!-\ldots-\!\bullet \in [\![C]\!]\} \subseteq Q$$

Poset of semantics



Analysis & Verification

Verification/static analysis by abstract interpretation

Define the syntax of programs $\mathtt{P} \in \mathbb{L}$

Define the concrete semantics of programs:

- $\mathscr{D}[\![\mathtt{P}]\!]$ concrete semantic domain
- $\forall\, \mathtt{P} \in \mathbb{L} : \mathcal{S}[\![\mathtt{P}]\!] \in \mathscr{D}[\![\mathtt{P}]\!]$ concrete semantics

Concrete/semantic properties: $\wp(\mathscr{D}[\![\mathtt{P}]\!])$

Collecting semantics: $\{\mathcal{S}[\![\mathtt{P}]\!]\} \in \wp(\mathscr{D}[\![\mathtt{P}]\!])$

(the strongest property of the semantics, which implies all other semantic properties)

Verification/static analysis by abstract interpretation

Define the abstraction:

- $\langle \wp(\mathscr{D}[\![\mathtt{P}]\!]),\, \subseteq \rangle \underset{\alpha[\![\mathtt{P}]\!]}{\overset{\gamma[\![\mathtt{P}]\!]}{\leftrightarrows}} \langle \mathscr{A}[\![\mathtt{P}]\!],\, \sqsubseteq \rangle$

Calculate the abstract semantics:

- $\mathcal{S}^{\#}[\![\mathtt{P}]\!] = \alpha[\![\mathtt{P}]\!](\{\mathcal{S}[\![\mathtt{P}]\!]\})$ exact abstraction
- $\mathcal{S}^{\#}[\![\mathtt{P}]\!] \sqsupseteq \alpha[\![\mathtt{P}]\!](\{\mathcal{S}[\![\mathtt{P}]\!]\})$ approximate abstraction

Soundness (by construction):

$$\forall\, \mathtt{P} \in \mathbb{L} : \forall\, Q \in \mathscr{A} : \mathcal{S}^{\#}[\![\mathtt{P}]\!] \sqsubseteq Q \Longrightarrow \mathcal{S}[\![\mathtt{P}]\!] \in \gamma[\![\mathtt{P}]\!](Q)$$

Verification/static analysis by abstract interpretation

Completeness (for exact abstractions only)

$\forall \mathtt{P} \in \mathbb{L} : \forall Q \in \mathscr{A}[\![\mathtt{P}]\!] : S[\![\mathtt{P}]\!] \in \gamma[\![\mathtt{P}]\!](Q) \Longrightarrow S^{\#}[\![\mathtt{P}]\!] \sqsubseteq Q$

Methodolody:

- Structural induction on programs $\mathtt{P}$
- Compositional definition(*) of $\mathscr{A}[\![\mathtt{P}]\!]$ and $\alpha[\![\mathtt{P}]\!]/\gamma[\![\mathtt{P}]\!]$
- Fixpoint abstraction/approximation for recursion

Verification for fixpoints is the main problem:

$\mathsf{lfp}^{\sqsubseteq} F^{\#}[\![\mathtt{P}]\!] \sqsubseteq Q$

(*) Patrick Cousot, Radhia Cousot: A Galois connection calculus for abstract interpretation. POPL 2014: 3-4 + Aux. mat. 15p.

Verification/static analysis by abstract interpretation

Method: find $I \in \mathscr{A}[\![\mathtt{P}]\!]$ such that $F^{\#}[\![\mathtt{P}]\!] I \sqsubseteq I \wedge I \sqsubseteq Q$
(so that $\mathsf{lfp}^{\sqsubseteq} F^{\#}[\![\mathtt{P}]\!] \sqsubseteq Q$, by Tarski)

- Verification/deductive/proof methods:
  - ask the end-user for the inductive argument I
- Static analysis:
  1. compute I knowing $F^{\#}[\![\mathtt{P}]\!]$ and Q
  2. compute I knowing $F^{\#}[\![\mathtt{P}]\!]$ (and later given any Q check that $I \sqsubseteq Q$)

Approximate abstractions

Approximate abstractions

The concrete properties of the standard semantics $S[\![\mathtt{P}]\!]$ that you want to prove may not always be provable in the abstract:

$$\forall Q \in \mathscr{A} : S[\![\mathtt{P}]\!] \in \gamma(Q) \;\substack{\Longleftarrow \\ \not\Longrightarrow}\; \overline{S}[\![\mathtt{P}]\!] \sqsubseteq Q$$

where

$$\overline{S}[\![\mathtt{P}]\!] \stackrel{\Delta}{\sqsupseteq} \alpha \circ S[\![\mathtt{P}]\!] \circ \gamma$$

Why abstraction may be approximate?

Example

$\{ x = y \wedge 0 \leqslant x \leqslant 10 \}$

```
x := x - y;
```

$\{ x = 0 \wedge 0 \leqslant y \leqslant 10 \}$

Interval abstraction:

$\{ x \in [0, 10] \wedge y \in [0, 10] \}$

```
x := x - y;
```

$\{ x \in [-10, 10] \wedge y \in [0, 10] \}$

(but for constants, the interval abstraction can't express equality)

Refinement

Refinement: good news

Problem: how to prove a valid abstract property $\alpha(\{\mathrm{lfp}\ F[\![\mathtt{P}]\!]\}) \sqsubseteq Q$ when $\alpha \circ F \sqsubseteq F^{\#} \circ \alpha$ but $\mathrm{lfp}\ F^{\#}[\![\mathtt{P}]\!] \not\sqsubseteq Q$?

It is always possible to refine $\langle \mathscr{A}, \sqsubseteq \rangle$ into a most abstract more precise abstraction $\langle \mathscr{A}', \sqsubseteq' \rangle$ such that

$$\langle \wp(\mathscr{D}), \subseteq \rangle \underset{\alpha'}{\overset{\gamma'}{\leftrightarrows}} \langle \mathscr{A}', \sqsubseteq' \rangle$$

and $\alpha' \circ F = F' \circ \alpha$ with $\mathrm{lfp}\ F'[\![\mathtt{P}]\!] \sqsubseteq' \alpha' \circ \gamma\ (Q)$

(thus proving $\mathrm{lfp}\ F[\![\mathtt{P}]\!] \in \gamma'(Q)$ which implies $\mathrm{lfp}\ F[\![\mathtt{P}]\!] \in \gamma(Q)$)

Roberto Giacobazzi, Francesco Ranzato, Francesca Scozzari: Making abstract interpretations complete. J. ACM 47(2): 361-416 (2000)

Refinement: bad news

But, refinements of an abstraction can be intrinsically incomplete

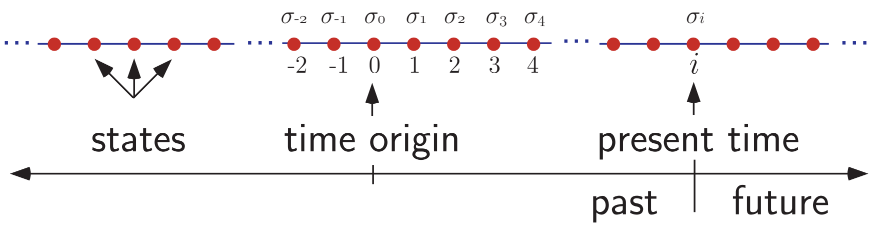
The only complete refinement of that abstraction for the collecting semantics is :

- the identity (i.e. no abstraction at all)

In that case, the only complete refinement of the abstraction is to the collecting semantics and any other refinement is always imprecise

Example of intrinsic approximate refinement

Consider executions traces $\langle i, \sigma \rangle$ with infinite past and future:



Patrick Cousot, Radhia Cousot: Temporal Abstract Interpretation. POPL 2000: 12-25

Example of intrinsic approximate refinement

Consider the temporal specification language $\mu^{\star}$ (containing LTL, CTL, CTL*, and Kozen's μ-calculus as fragments):

$\varphi ::=$	$\boldsymbol{\sigma}_S$	$S \in \wp(\mathbb{S})$	state predicate
$\mid$	$\boldsymbol{\pi}_t$	$t \in \wp(\mathbb{S} \times \mathbb{S})$	transition predicate
$\mid$	$\oplus \varphi_1$		next
$\mid$	$\varphi_1^{\curvearrowleft}$		reversal
$\mid$	$\varphi_1 \vee \varphi_2$		disjunction
$\mid$	$\neg \varphi_1$		negation
$\mid$	X	$X \in \mathbb{X}$	variable
$\mid$	$\boldsymbol{\mu} X \cdot \varphi_1$		least fixpoint
$\mid$	$\boldsymbol{\nu} X \cdot \varphi_1$		greatest fixpoint
$\mid$	$\forall \varphi_1 : \varphi_2$		universal state closure

Patrick Cousot, Radhia Cousot: Temporal Abstract Interpretation. POPL 2000: 12-25

Example of intrinsic approximate refinement

Consider universal model-checking abstraction:

$$\mathrm{MC}^{\forall}_{M}(\phi) = \alpha^{\forall}_{M}(\llbracket \phi \rrbracket) \in \wp(\mathit{Traces}) \rightarrow \wp(\mathit{States})$$

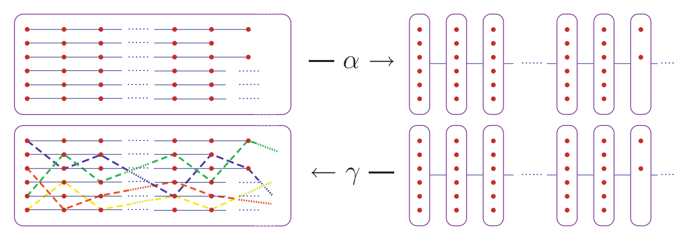
$$= \{ s \in \mathit{States} \mid \forall \langle i, \sigma \rangle \in \mathit{Traces}_M \,.\, (\sigma_i = s) \Rightarrow \langle i, \sigma \rangle \in \llbracket \phi \rrbracket \}$$

where M is defined by a transition system

(and dually the existential model-checking abstraction)

Example of intrinsic approximate refinement

The abstraction from a set of traces to a trace of sets is sound but *incomplete*, even for finite systems (*)

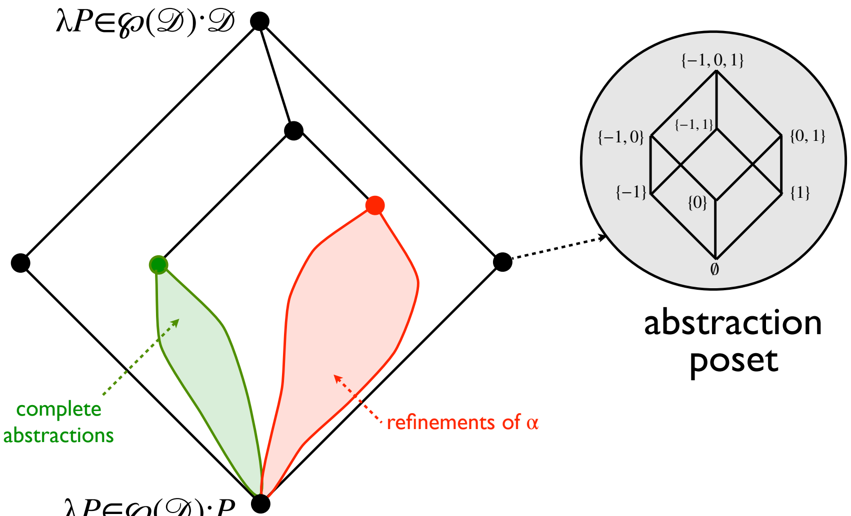


Any refinement of this abstraction is *incomplete* (but to the infinite past/future trace semantics itself) (**)

(*) Patrick Cousot, Radhia Cousot: Temporal Abstract Interpretation. POPL 2000: 12-25

(**) Roberto Giacobazzi, Francesco Ranzato: Incompleteness of states w.r.t. traces in model checking. Inf. Comput. 204(3): 376-407 (2006)

Intrinsic approximate refinement



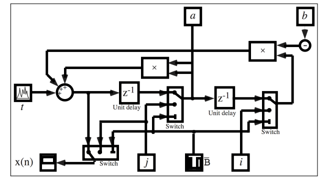
abstraction poset

poset of abstractions

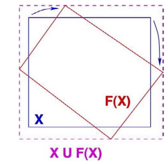
In general refinement does not terminate

- Example:filter invariant abstraction:

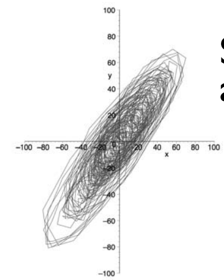
2nd order filter:



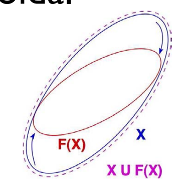
Unstable polyhedral abstraction:



Counter-example guided refinement will indefinitely add missing points according to the execution trace:



Stable ellipsoidal abstraction:



Julien Bertrane, Patrick Cousot, Radhia Cousot, Jérôme Feret, Laurent Mauborgne, Antoine Miné, & Xavier Rival. Static Analysis and Verification of Aerospace Software by Abstract Interpretation. In *AIAA Infotech@@Aerospace 2010*, Atlanta, Georgia. American Institute of Aeronautics and Astronautics, 20—22 April 2010. © AIAA.

In general refinement does not terminate

Narrowing is needed to stop infinite iterated automatic refinements:

e.g. SLAM stops refinement after 20mn, now abandoned

Intelligence is needed for refinement:

e.g. human-driven refinement of Astrée

Thomas Ball, Vladimir Levin, Sriram K. Rajamani: A decade of software model checking with SLAM. Commun. ACM 54(7): 68-76 (2011)

Julien Bertrane, Patrick Cousot, Radhia Cousot, Jérôme Feret, Laurent Mauborgne, Antoine Miné, & Xavier Rival. Static Analysis and Verification of Aerospace Software by Abstract Interpretation. In *AIAA Infotech@@Aerospace 2010*, Atlanta, Georgia. American Institute of Aeronautics and Astronautics, 20—22 April 2010. © AIAA.

Finite versus infinite abstractions

[In]finite abstractions

Given a program $\mathtt{P}$ and a program property Q which holds (*i.e.* $\mathsf{lfp}\ F[\![\mathtt{P}]\!] \in Q$) there exists a most abstract abstraction in a finite domain $\mathscr{A}[\![\mathtt{P}]\!]$ to prove it (*)

Example:

`x=0; while x<1 do x++` $\longrightarrow \{\bot, [0,0], [0,1], [-\infty,\infty]\}$

`x=0; while x<2 do x++` $\longrightarrow \{\bot, [0,0], [0,1], [0,2], [-\infty,\infty]\}$

…

`x=0; while x<`n` do x++` $\longrightarrow \{\bot, [0,0], [0,1], [0,2], [0,3], \ldots, [0,n], [-\infty,\infty]\}$

…

(*) Patrick Cousot: Partial Completeness of Abstract Fixpoint Checking. SARA 2000: 1-25

[In]finite abstractions

No such domain exists for infinitely many programs

1. $\bigcup_{\mathtt{P} \in \mathbb{L}} \mathscr{A}[\![\mathtt{P}]\!]$ is infinite

   Example: $\{\bot, [0,0], [0,1], [0,2], [0,3], \ldots, [0,n], [0,n+1], \ldots, [-\infty,\infty]\}$

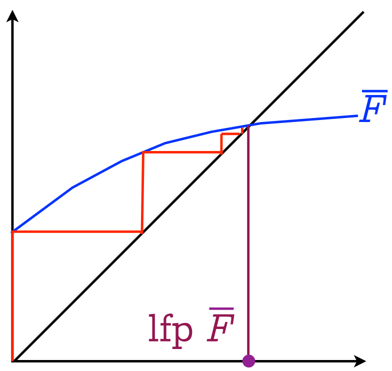
2. $\lambda \mathtt{P} \in \mathbb{L} \cdot \mathscr{A}[\![\mathtt{P}]\!]$ is not computable (for undecidable properties)

$\Longrightarrow$ finite abstractions will fail infinitely often while infinite abstractions will succeed!

Fixpoint approximation in infinite abstractions

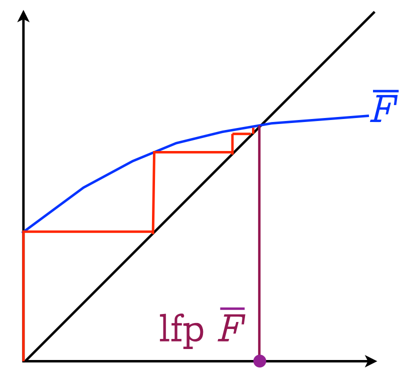
Abstract Induction (in non-Noetherian domains)

Convergence acceleration

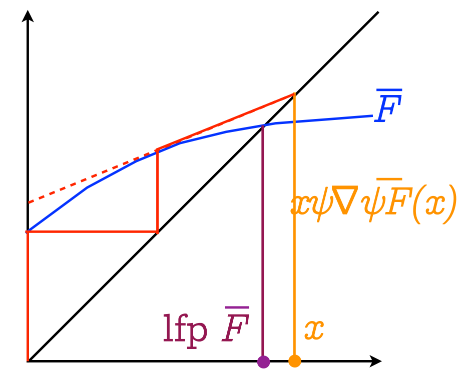


Infinite iteration

Convergence acceleration



Infinite iteration



Accelerated iteration with widening
(e.g. with a widening based on the derivative as in Newton-Raphson method(*))

(*) Javier Esparza, Stefan Kiefer, Michael Luttenberger: Newtonian program analysis. J. ACM 57(6): 33 (2010)

Problem with infinite abstractions

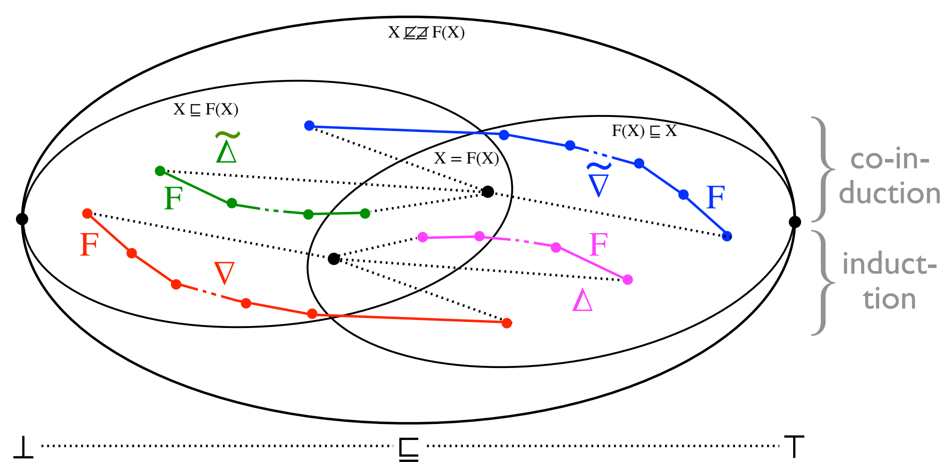
For non-Noetherian iterations, we need

- finitary abstract induction, and
- finitary passage to the limit

$X^0=\perp, \ldots, X^{n+1} = \mathfrak{F}(X^0, \ldots, X^n, F(X^0), \ldots, F(X^n)), \ldots, \lim_{n\to\infty} X^n$

		iteration converging	
	$\mathfrak{F}$	above the limit	below the limit
Iteration starting from	below the limit	widening $\triangledown$	dual narrowing $\widetilde{\triangle}$
	above the limit	narrowing $\triangle$	dual widening $\widetilde{\triangledown}$

[Semi-]dual abstract induction methods



(separate from termination conditions)

Examples of widening/narrowing

Abstract induction for intervals:

- a widening [1,2]

$(x \bar{\nabla} y) = \underline{\text{cas}}\ x \in V_a,\ y \in V_a\ \underline{\text{dans}}$
$\quad \vdash \square,\ ? \Longrightarrow y\ ;$
$\quad \vdash ?,\ \square \Longrightarrow x\ ;$
$\quad \vdash [n_1, m_1], [n_2, m_2] \Longrightarrow$
$\qquad [\underline{\text{si}}\ n_2 < n_1\ \underline{\text{alors}}\ -\infty\ \underline{\text{sinon}}\ n_1\ \underline{\text{fsi}}\ ;$
$\qquad\ \underline{\text{si}}\ m_2 > m_1\ \underline{\text{alors}}\ +\infty\ \underline{\text{sinon}}\ m_1\ \underline{\text{fsi}}]\ ;$
$\underline{\text{fincas}}\ ;$

$[a_1, b_1]\ \bar{\nabla}\ [a_2, b_2] =$
$\quad [\underline{\text{if}}\ a_2 < a_1\ \underline{\text{then}}\ -\infty\ \underline{\text{else}}\ a_1\ \underline{\text{fi}},$
$\qquad \underline{\text{if}}\ b_2 > b_1\ \underline{\text{then}}\ +\infty\ \underline{\text{else}}\ b_1\ \underline{\text{fi}}]$

- a narrowing [2]

$[a_1, b_1]\ \bar{\Delta}\ [a_2, b_2] =$
$\quad [\underline{\text{if}}\ a_1 = -\infty\ \underline{\text{then}}\ a_2\ \underline{\text{else}}\ \text{MIN}\ (a_1, a_2),$
$\qquad \underline{\text{if}}\ b_1 = +\infty\ \underline{\text{then}}\ b_2\ \underline{\text{else}}\ \text{MAX}\ (b_1, b_2)]$

[1] Patrick Cousot, Radhia Cousot: Vérification statique de la cohérence dynamique des programmes, Rapport du contrat IRIA-SESORI No 75-032, 23 septembre 1975.
[2] Patrick Cousot, Radhia Cousot: Abstract Interpretation: A Unified Lattice Model for Static Analysis of Programs by Construction or Approximation of Fixpoints. POPL 1977: 238-252

On widening/narrowing/and their duals

Because the abstract domain is non-Noetherian, *any* widening/narrowing/duals can be *strictly* improved infinitely many times (*i.e.* no best widening)

E.g. widening with thresholds [1]

$\forall x \in \bar{L}_2,\ \bot\ \nabla_2(j)\ x = x\ \nabla_2(j) \bot = x$
$[l_1, u_1]\ \nabla_2(j)\ [l_2, u_2]$
$\quad = [\mathit{if}\ 0 \leq l_2 < l_1\ \mathit{then}\ 0\ \mathit{elsif}\ l_2 < l_1\ \mathit{then}\ -b-1\ \mathit{else}\ l_1\ \mathit{fi},$
$\qquad \mathit{if}\ u_1 < u_2 \leq 0\ \mathit{then}\ 0\ \mathit{elsif}\ u_1 < u_2\ \mathit{then}\ b\ \mathit{else}\ u_1\ \mathit{fi}]$

Any *terminating* widening is not increasing (in its first parameter)

Any abstraction done with Galois connections *can be done* with widenings (*i.e.* a widening calculus)

[1] Patrick Cousot, Semantic foundations of program analysis, Ch. 10 of Program flow analysis: theory and practice, N. Jones & S. Muchnich (eds), Prentice Hall, 1981.

Infinitary static analysis with abstract induction

Widening

$\langle \mathscr{A}, \sqsubseteq \rangle$ poset

$\nabla \in \mathscr{A} \times \mathscr{A} \longrightarrow \mathscr{A}$

Sound widening (upper bound):

$\forall x, y \in \mathscr{A}: x \sqsubseteq x \nabla y \ \wedge\ y \sqsubseteq x \nabla y$

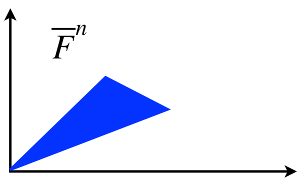
Terminating widening: for any $\langle x^n \in \mathscr{A}, n \in \mathbb{N} \rangle$, the sequence $y^0 \triangleq x^0, \ldots, y^{n+1} \triangleq y^n \nabla x^n, \ldots$ is *ultimately stationary* $(\exists \varepsilon \in \mathbb{N}: \forall n \geqslant \varepsilon: y^n = y^\varepsilon)$

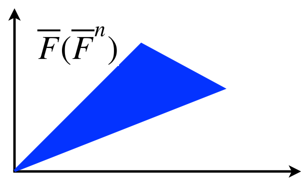
(Note: sound and terminating are independent properties)

Patrick Cousot, Radhia Cousot: Abstract Interpretation: A Unified Lattice Model for Static Analysis of Programs by Construction or Approximation of Fixpoints. POPL 1977: 238-252

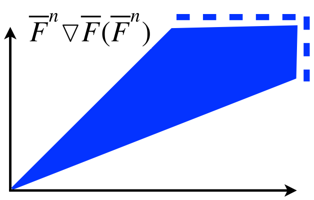
Example: (simple) widening for polyhedra

Iterates





Widening



Patrick Cousot. Méthodes itératives de construction et d'approximation de points fixes d'opérateurs monotones sur un treillis, analyse sémantique des programmes. *Thèse És Sciences Mathématiques*, Université Joseph Fourier, Grenoble, France, 21 March 1978.
Patrick Cousot, Nicolas Halbwachs: Automatic Discovery of Linear Restraints Among Variables of a Program. POPL 1978: 84-96

Iteration with widening for static analysis

Problem: compute I such that $\mathrm{lfp}^{\sqsubseteq} F \sqsubseteq I \sqsubseteq Q$

Compute I as the limit of the iterates:

- $X^0 \triangleq \bot$,
- $X^{n+1} \triangleq X^n$ when $F(X^n) \sqsubseteq X^n$ so $I = X^n$
- $X^{n+1} \triangleq (X^n \triangledown F(X^n)) \vartriangle Q$ otherwise

I can be improved by an iteration with narrowing $\vartriangle$

Check that $F(I) \sqsubseteq Q$

Example: Astrée

Patrick Cousot, Radhia Cousot: Abstract Interpretation: A Unified Lattice Model for Static Analysis of Programs by Construction or Approximation of Fixpoints. POPL 1977: 238-252

Dual narrowing

$\langle \mathscr{A}, \sqsubseteq \rangle$ poset

$\widetilde{\vartriangle} \in \mathscr{A} \times \mathscr{A} \longrightarrow \mathscr{A}$

Sound dual narrowing (interpolation):

$\forall x, y \in \mathscr{A}: x \sqsubseteq y \implies x \sqsubseteq x \,\widetilde{\vartriangle}\, y \sqsubseteq y$

Terminating dual narrowing: for any $\langle x^n \in \mathscr{A}, n \in \mathbb{N} \rangle$, the sequence $y^0 \triangleq x^0, \ldots, y^{n+1} \triangleq y^n \,\widetilde{\vartriangle}\, x^n, \ldots$ is ultimately stationary ($\exists \varepsilon \in \mathbb{N}: \forall n \geqslant \varepsilon: y^n = y^\varepsilon$)

(Note: sound and terminating are independent properties)

Cousot, P. Méthodes itératives de construction et d'approximation de points fixes d'opérateurs monotones sur un treillis, analyse sémantique de programmes (in French). Thèse d'État ès sciences mathématiques, Université scientifique et médicale de Grenoble, France 1978.

Iteration with dual narrowing for static checking

Problem: find I such that $\mathrm{lfp}^{\sqsubseteq} F \sqsubseteq I \sqsubseteq Q$

Compute I as the limit of the iterates:

- $X^0 \triangleq \bot$,
- $X^{n+1} \triangleq X^n$ when $F(X^n) \sqsubseteq X^n$ so $I = X^n$
- $X^{n+1} \triangleq F(X^n) \,\widetilde{\vartriangle}\, Q$, otherwise

Check that $F(I) \sqsubseteq Q$

Example: First-order logic + Graig interpolation (with some choice of one of the solutions, control of combinatorial explosion, and convergence enforcement)

Kenneth L. McMillan: Applications of Craig Interpolants in Model Checking. TACAS 2005: 1-12

Industrialization

Daniel Kästner, Christian Ferdinand, Stephan Wilhelm, Stefana Nevona, Olha Honcharova, Patrick Cousot, Radhia Cousot, Jérôme Feret, Laurent Mauborgne, Antoine Miné, Xavier Rival, and Élodie-Jane Sims. Astrée: Nachweis der Abwesenheit von Laufzeitfehlern. In *Workshop "Entwicklung zuverlässiger Software-Systeme"*, Regensburg, Germany, June 18th, 2009.

Olivier Bouissou, Éric Conquet, Patrick Cousot, Radhia Cousot, Jérôme Feret, Khalil Ghorbal, Éric Goubault, David Lesens, Laurent Mauborgne, Antoine Miné, Sylvie Putot, Xavier Rival, & Michel Turin. Space Software Validation using Abstract Interpretation. In *Proc. of the Int. Space System Engineering Conf., Data Systems in Aerospace (DASIA 2009)*. Istambul, Turkey, May 2009, 7 pages. ESA.

Jean Souyris, David Delmas: Experimental Assessment of Astrée on Safety-Critical Avionics Software. SAFECOMP 2007: 479-490

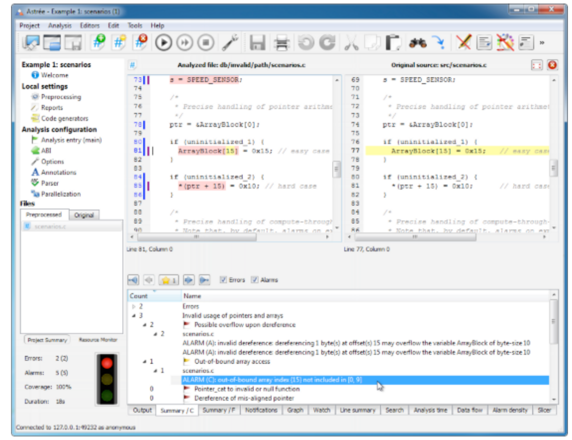
David Delmas, Jean Souyris: Astrée: From Research to Industry. SAS 2007: 437-451

Jean Souyris: Industrial experience of abstract interpretation-based static analyzers. IFIP Congress Topical Sessions 2004: 393-400

Stephan Thesing, Jean Souyris, Reinhold Heckmann, Famantanantsoa Randimbivololona, Marc Langenbach, Reinhard Wilhelm, Christian Ferdinand: An Abstract Interpretation-Based Timing Validation of Hard Real-Time Avionics Software. DSN 2003: 625-632

Astrée

Commercially available: `www.absint.com/astree/`

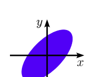


Effectively used in production to qualify truly large and complex software in transportation, communications, medicine, *etc*

Bruno Blanchet, Patrick Cousot, Radhia Cousot, Jérôme Feret, Laurent Mauborgne, Antoine Miné, David Monniaux, Xavier Rival: **A static analyzer for large safety-critical software.** *PLDI 2003*: 196-207

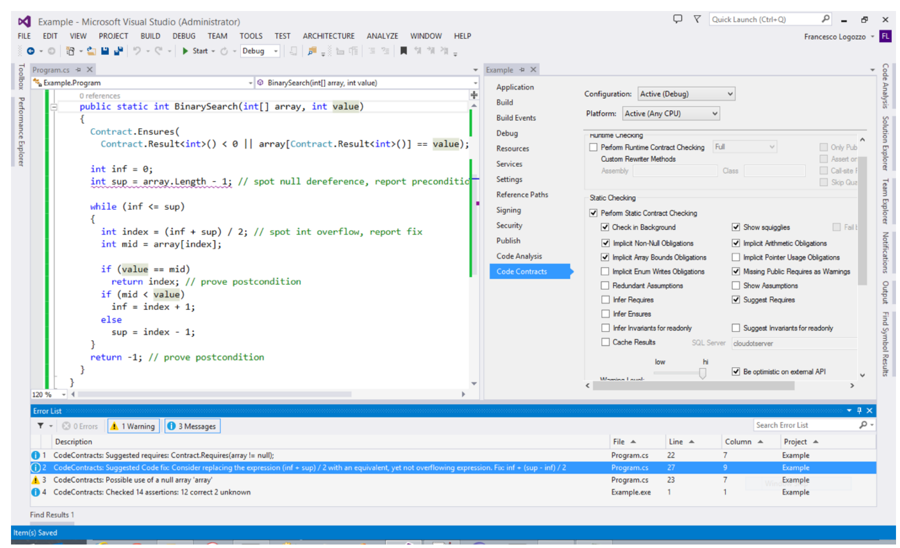
Example of domain-specific abstraction: ellipses

```
typedef enum {FALSE = 0, TRUE = 1} BOOLEAN;
BOOLEAN INIT; float P, X;
void filter () {
  static float E[2], S[2];
  if (INIT) { S[0] = X; P = X; E[0] = X; }
  else { P = (((((0.5 * X) - (E[0] * 0.7)) + (E[1] * 0.4))
              + (S[0] * 1.5)) - (S[1] * 0.7)); }
  E[1] = E[0]; E[0] = X; S[1] = S[0]; S[0] = P;
  /* S[0], S[1] in [-1327.02698354, 1327.02698354] */
}
void main () { X = 0.2 * X + 5; INIT = TRUE;
  while (1) {
    X = 0.9 * X + 35; /* simulated filter input */
    filter (); INIT = FALSE; }
}
```



Code Contract Static Checker (cccheck)

Was available within MS Visual Studio, now public domain



Manuel Fähndrich, Francesco Logozzo: **Static Contract Checking with Abstract Interpretation.** *FoVeOOS 2010*: 10-30

Comments on screenshot (courtesy Francesco Logozzo)

1. A screenshot from Clousot/`cccheck` on the classic binary search.
2. The screenshot shows from left to right and top to bottom
    1. C# code + CodeContracts with a buggy BinarySearch
    2. `cccheck` integration in VS (right pane with all the options integrated in the VS project system)
    3. `cccheck` messages in the VS error list
3. The features of `cccheck` that it shows are:
    1. basic abstract interpretation:
        1. the loop invariant to prove the array access correct and that the arithmetic operation may overflow is inferred fully automatically
        2. different from deductive methods as e.g. ESC/Java or Boogie where the loop invariant must be provided by the end-user
    2. inference of necessary preconditions:
        1. Clousot finds that array may be null (message 3)
        2. Clousot suggests and propagates a necessary precondition invariant (message 1)
    3. array analysis (+ disjunctive reasoning):
        1. to prove the postcondition should infer property of the content of the array
        2. please note that the postcondition is true even if there is no precondition requiring the array to be sorted.
    4. verified code repairs:
        1. from the inferred loop invariant does not follow that index computation does not overflow

Conclusion

Abstract interpretation

Intellectual tool (not to be confused with its specific application to iterative static analysis with $\triangledown$ & $\vartriangle$)

No cathedral would have been built without plumb-line and square, certainly not enough for skyscrapers:

Powerful tools are needed for progress and applicability of formal methods

Abstract interpretation

Varieties of researchers in formal methods:

(i) explicitly use abstract interpretation, and are happy to extend its scope and broaden its applicability

(ii) implicitly use abstract interpretation, and hide it

(iii) pretend to use abstract interpretation, but misuse it

(iv) don't know that they use abstract interpretation, but would benefit from it

Never too late to upgrade

The End

The End
Thank You