

□ 1. Last time

□ 2. Weensy OS walk through

□ 3. Weensy OS context switches

□ 4. gdb

lab 4 work here

↵ : thinks it is interacting with hardware



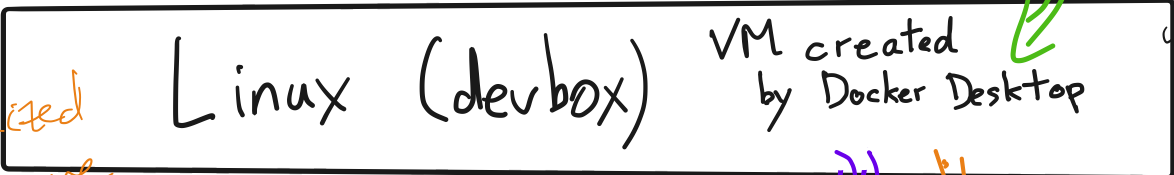
cpu, h/w

↵ : uses syscall interface (and possibly virtualization extensions)

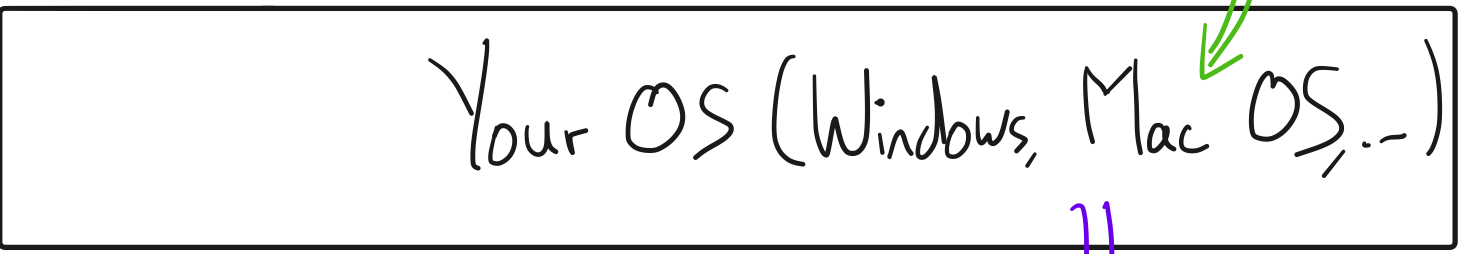


Process

↵ : paravirtualized (in between hardware and syscall)



cpu, h/w



Hints:

- processes: files matching p-*.c
- kernel code: files matching k-*.c, k-*.s
- system calls and returns // (*) /* cousin of mmap() */

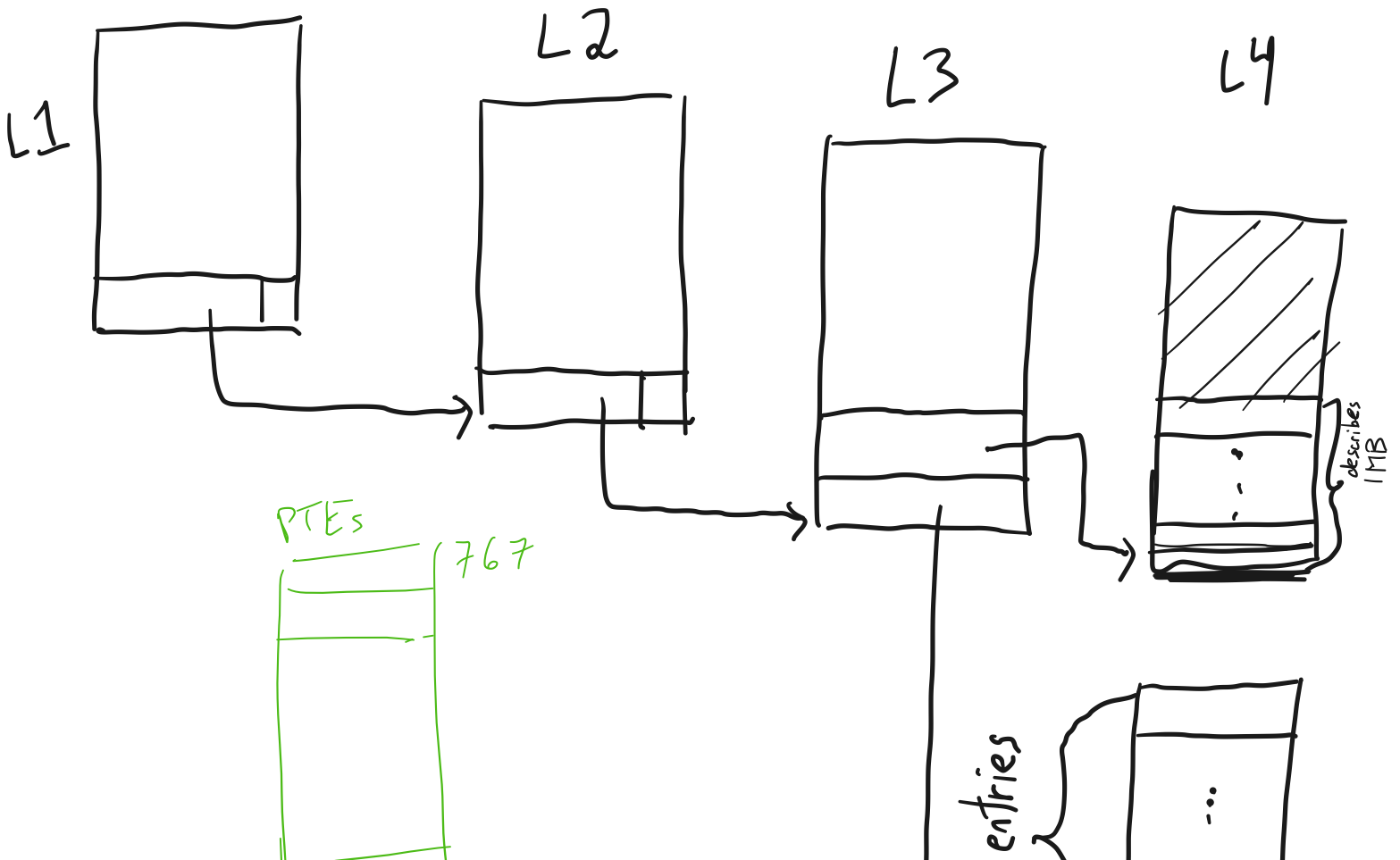
sys_page_alloc, lookin process.h

kernel returns: exception_return();
%rax contains return value of system call
(errno, 0)

%rdi:
arg to
syscall

- you'll use virtual-memory-map() = (NULL vs. non-NULL)
- pay attention to the allocator argument
- make sure your allocator initializes the page table
memset(addr, 0, len);

- a process's virtual address space: 3 MB. What's the page table structure?

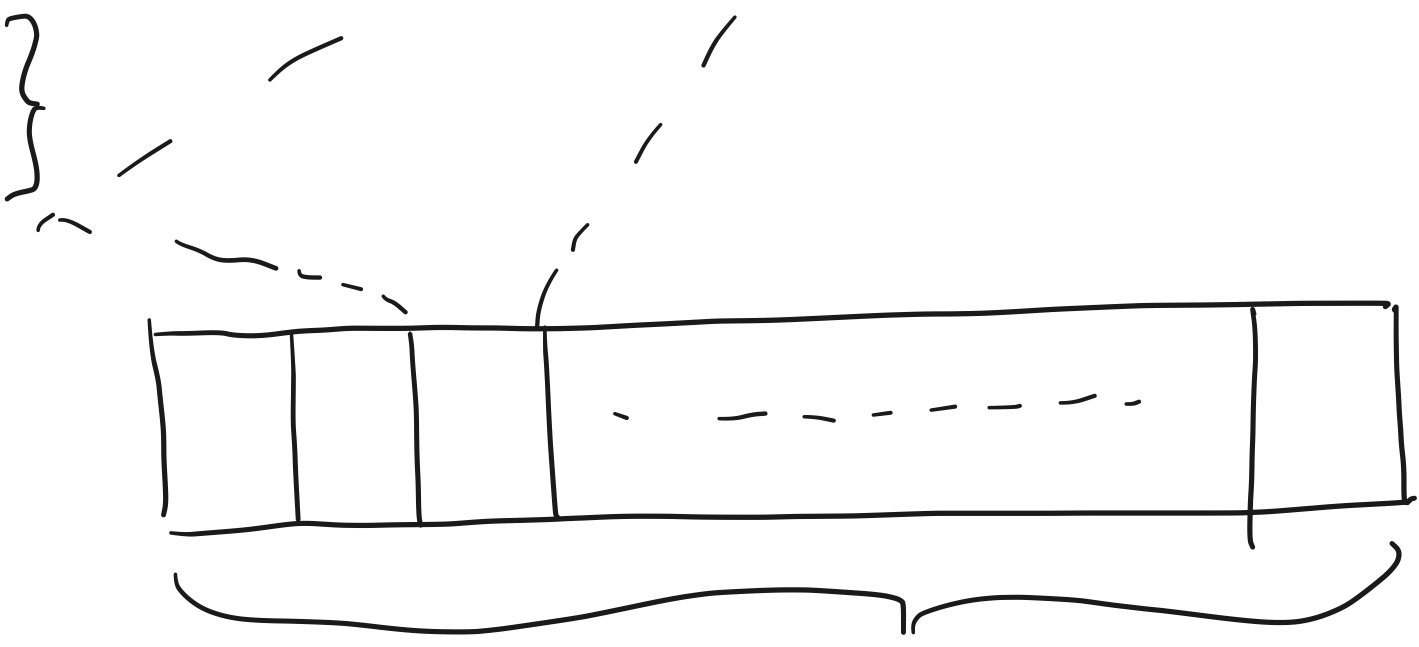




each entry
describes a mapping
for one page
(4KB)

PCB \equiv struct proc (in kernel.h)

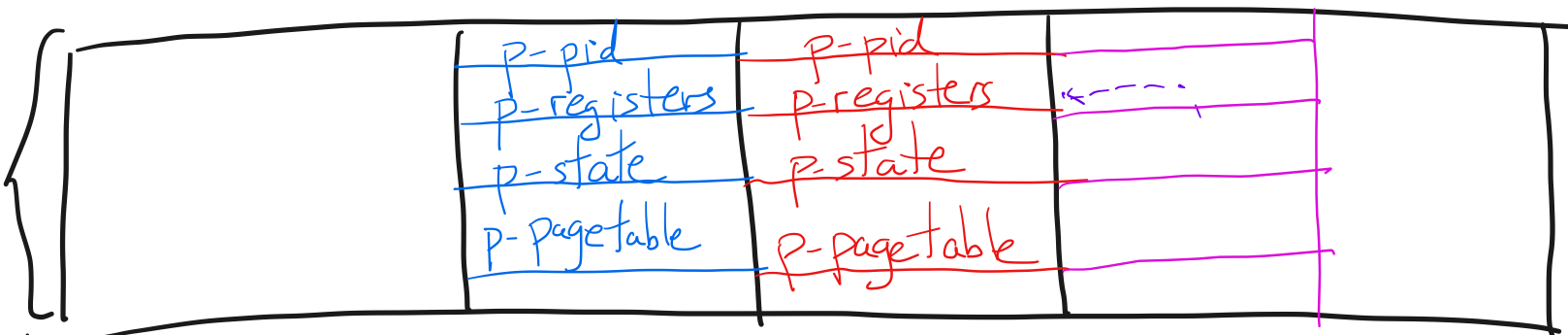
struct physical_pageinfo {
 int8_t owner;
 int8_t refcount;
}



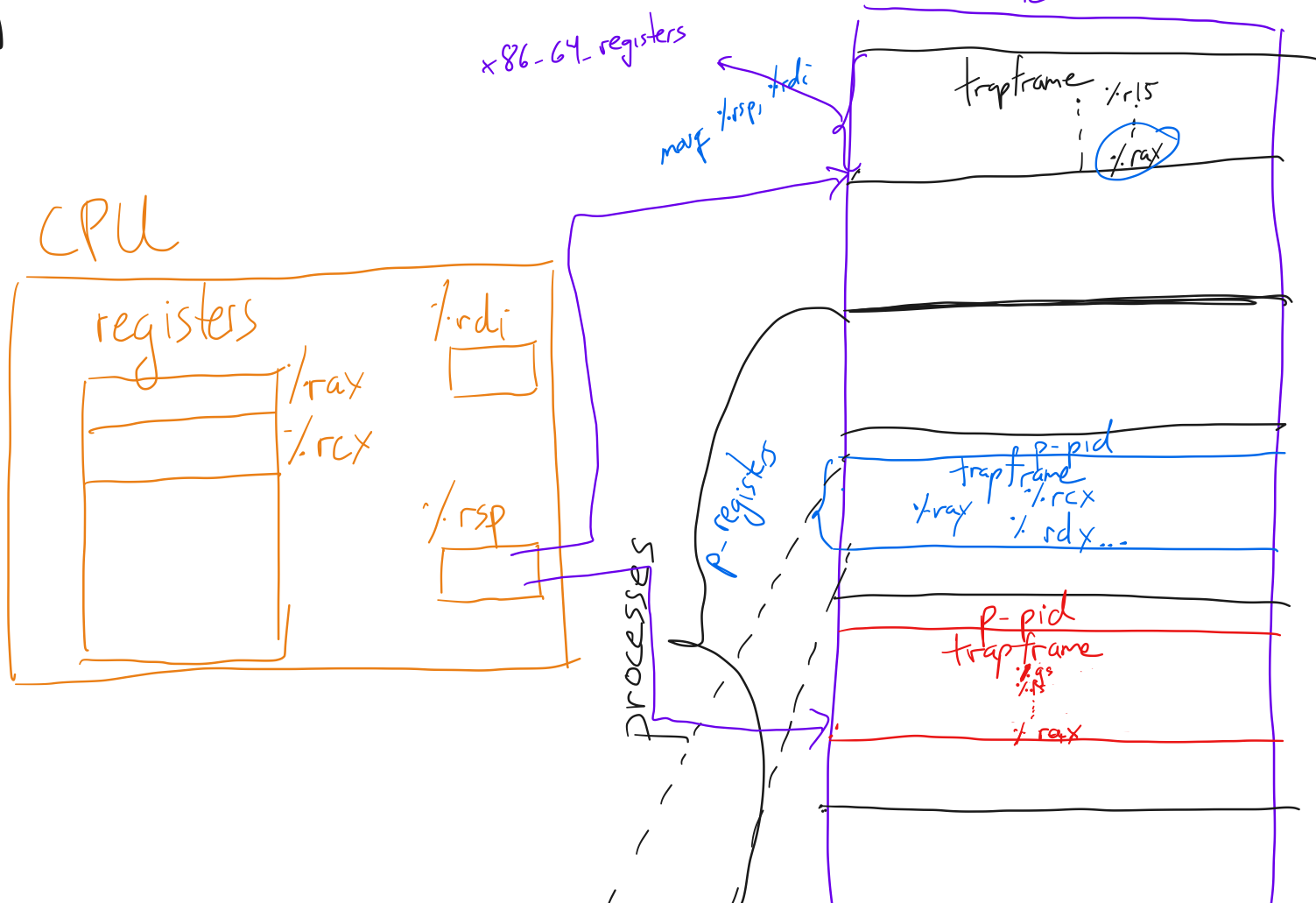
one per physical page

this is not a page table; it is bookkeeping.

Context switches in Weensy OS



processes



trapframe

