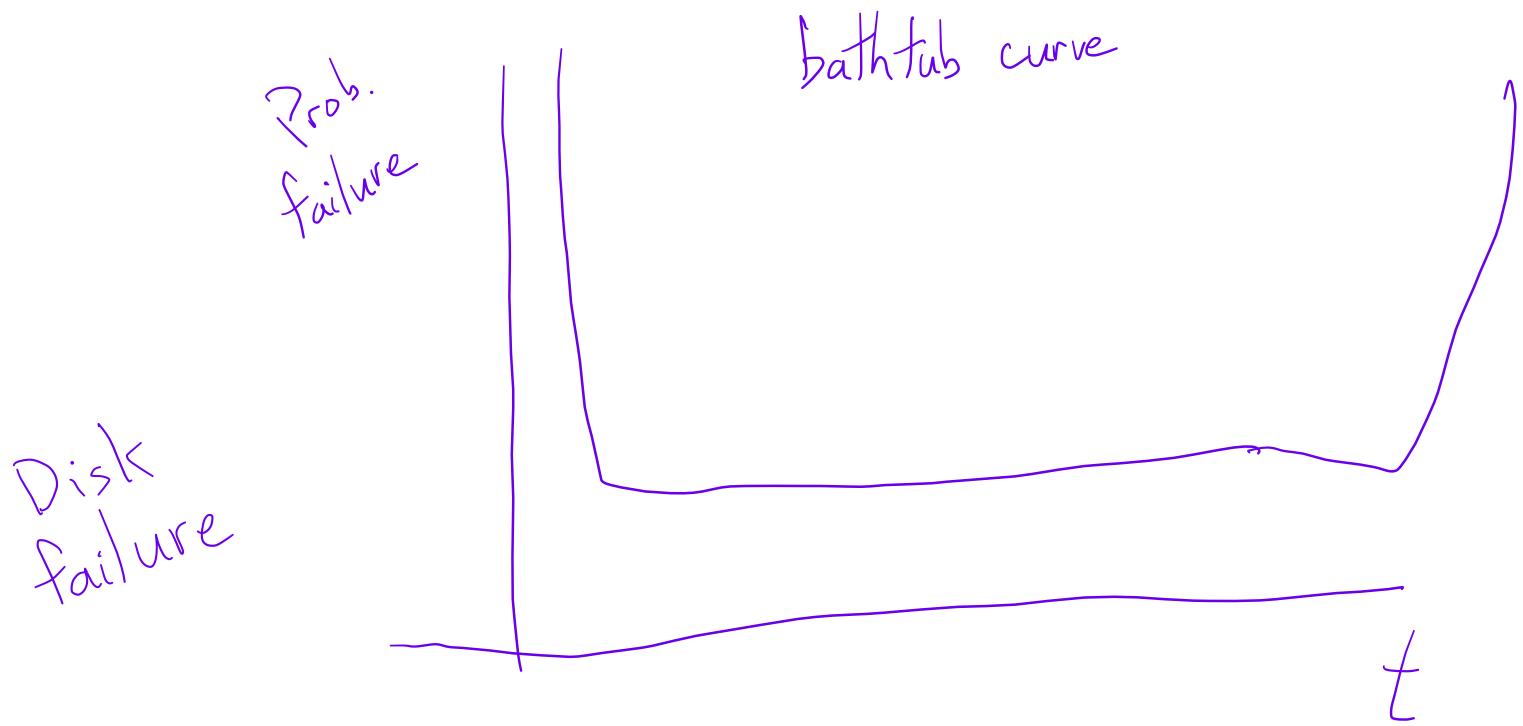


- ▢ 1. Last time
 - ▢ 2. Intro to file systems
 - ▢ 3. Files
 - ▢ 4. Implementing files
-



2.

Intro to file systems

What does a FS do?

- provide persistence

- create a way to name data on the disk

FS: can be implemented in lots of places

- We focus on the disk, generalize later

Note: disk is the 1^{st} thing we've seen that is both modifiable and persistent.

3.

Files

What is a file?

From user's view: a named, contiguous run of bytes

From FS's view: collection of disk blocks

Job of a FS:

map {file, offset in file} $\xrightarrow{\text{FS}}$ disk address

operations:

create(file), delete(file), read(), write()

Goal: operations have as few disk accesses as possible
and minimal space overhead

4. Implementing files

- A. Contiguous
- B. Linked files
- C. Indexed files

Assume for now that a given file's metadata is known to the system.

Access patterns to support:

- Sequential
- Random access

Ideal is good sequential + good random access performance

Candidate designs:

- A. Contiguous allocation

user pre-specifies length

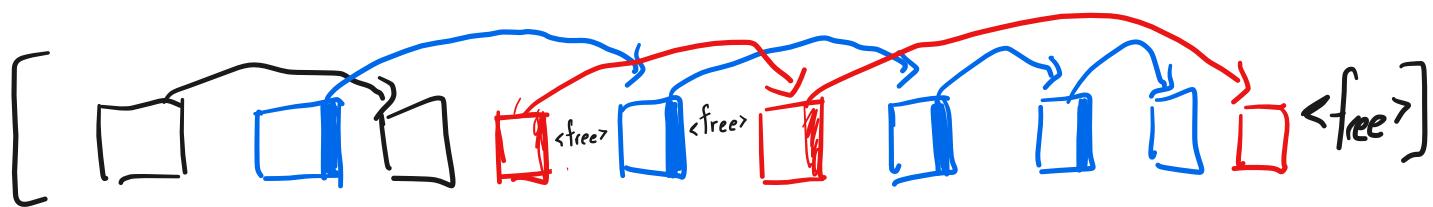
$[<1 \text{ free}> a_1 \ a_2 \ a_3 <5 \text{ free}> b_1 \ b_2 <1 \text{ free}>]$

- + fast access, both seq. and R.A.
- + simple

- fragmentation

B. Linked files

metadata is pointer (disk address) to file's first block



+ Seq. access easy + probably fast

+ no more fragmentation

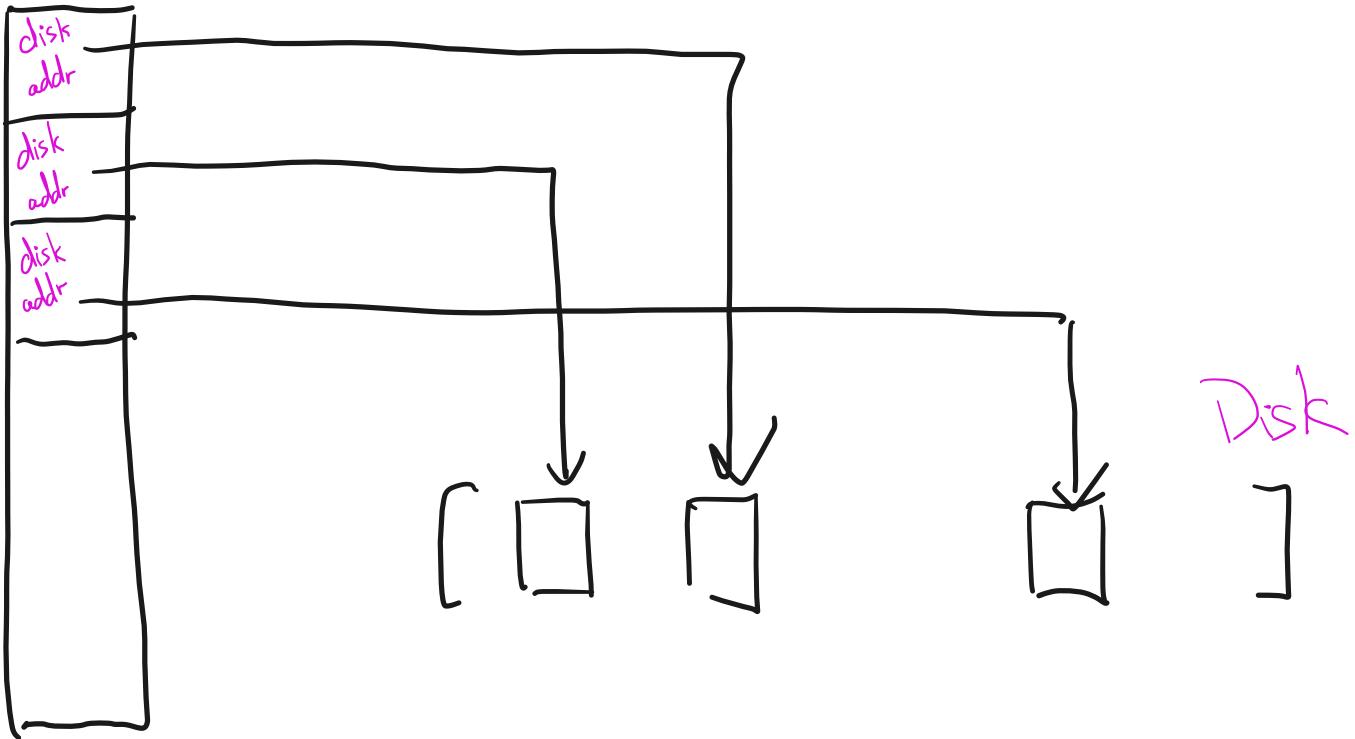
- R.A. is a disaster

- alignment of data can get messed up

C. Indexed files

1 1 1 : attempt 1

metadata.

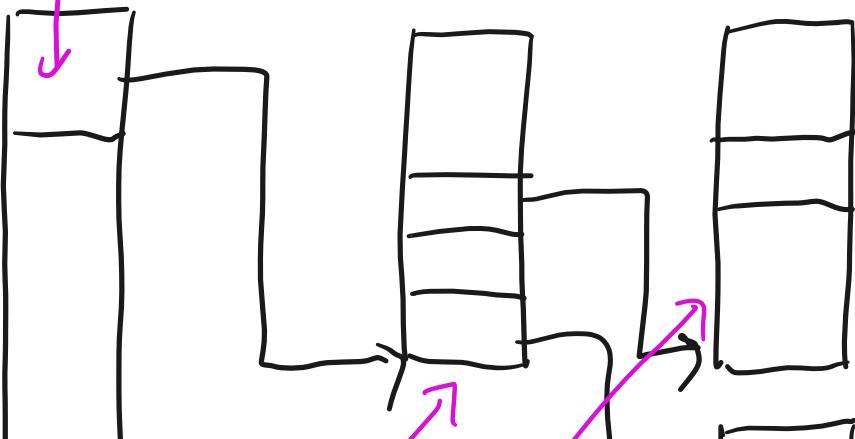


+ Seg, R.A. easy
+

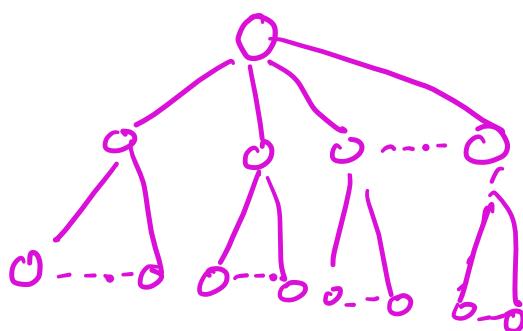
- storing this array is impractical

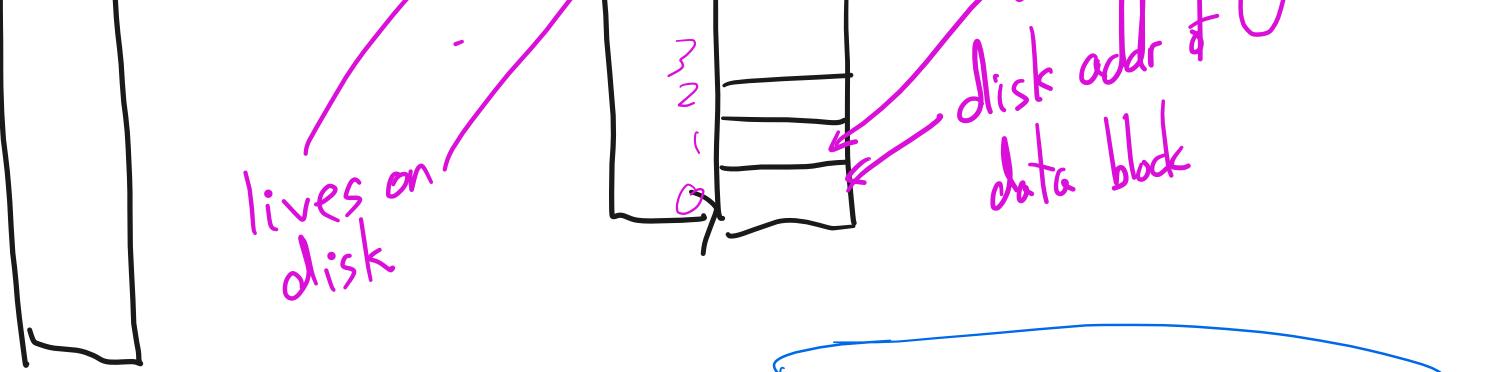
disk addr of
metadata

attempt 2

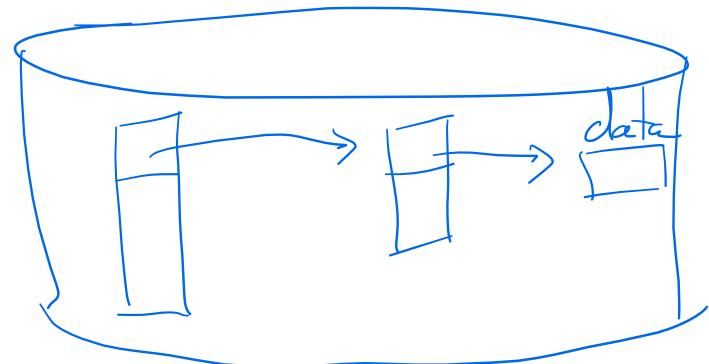


disk addr of 1st datablock
II 1st 2nd 3rd ... nth





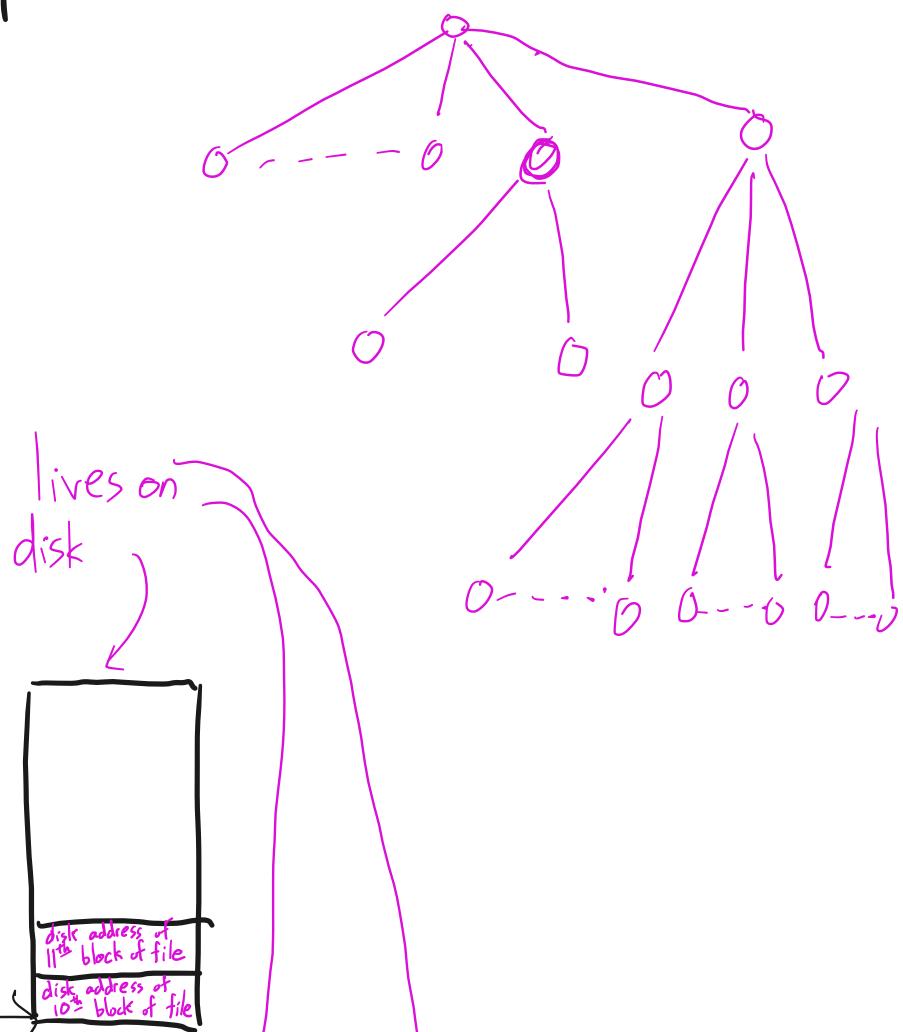
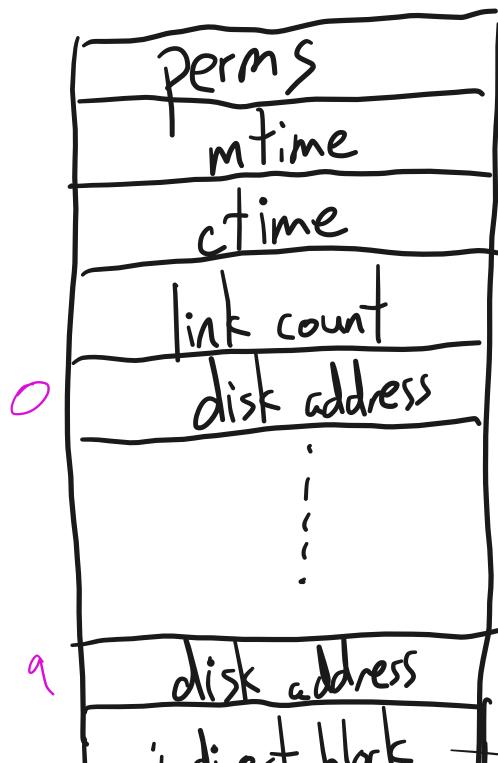
+ metadata is compact

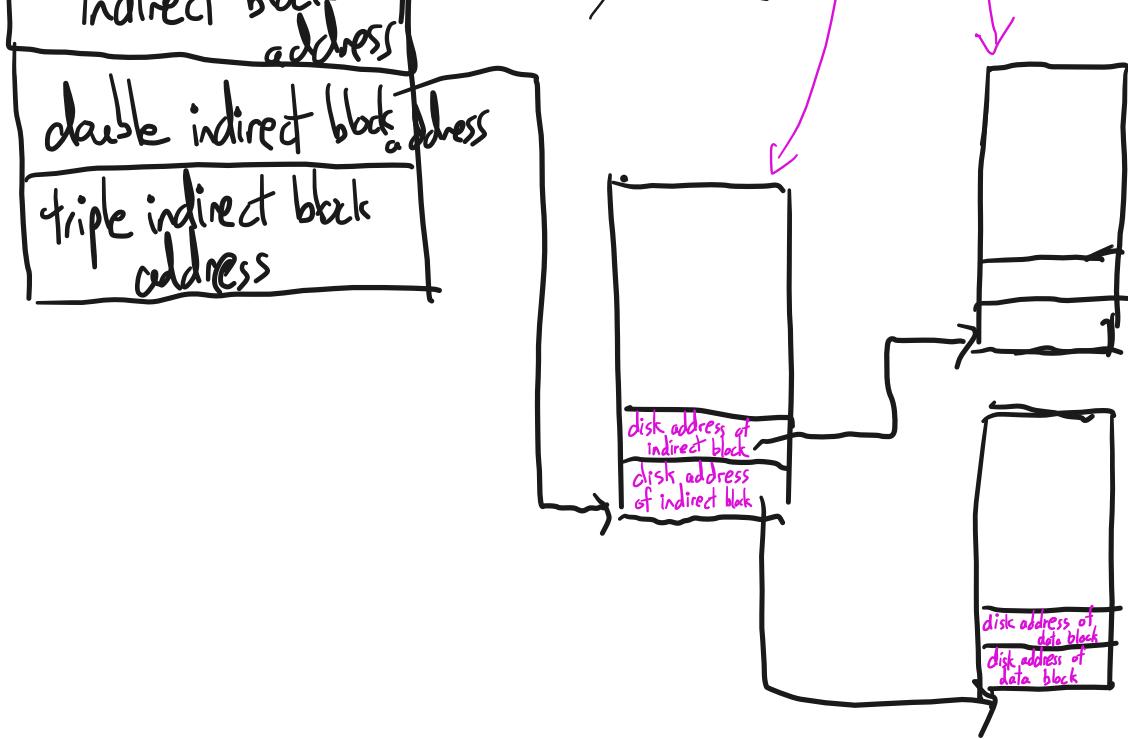


- looking up any block requires many disk accesses

attempt 3

Metadata: inode

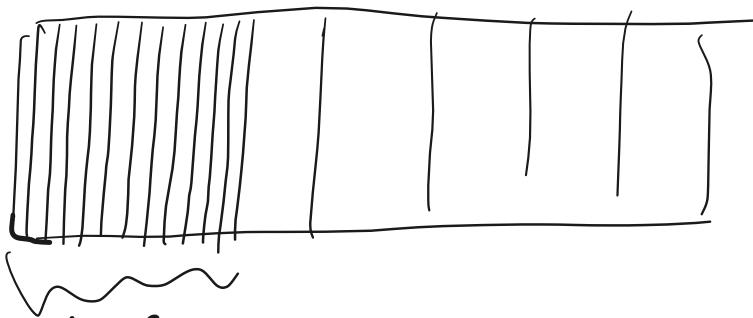




- + fast access to small files
- + Max length can be enormous
- + Simple, easy to build
- worst case # of accesses: pretty bad
- " " " space ovhd.
- locality issues

inodes: stored in a fixed-size array, known location

vocab: "inumber"



stat (&sb);

slots for
inodes

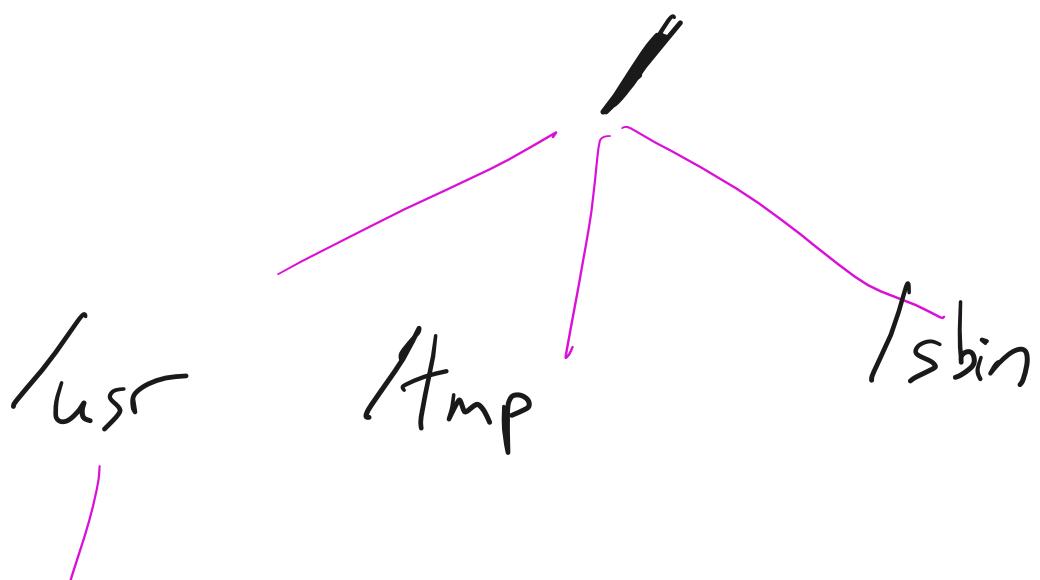
5. Directories (next time)

/bin /

/sbin

/usr

/tmp



/usr/mw

}

/lib

}

Kernel.c ..