- [x] 1. Last time
- [x] 2. x86-64: addresses
- [x] 3. x86-64: page table structures
- [x] midterm logistics
- [x] 4. TLBs
- [ ] 5. Where does the OS live?

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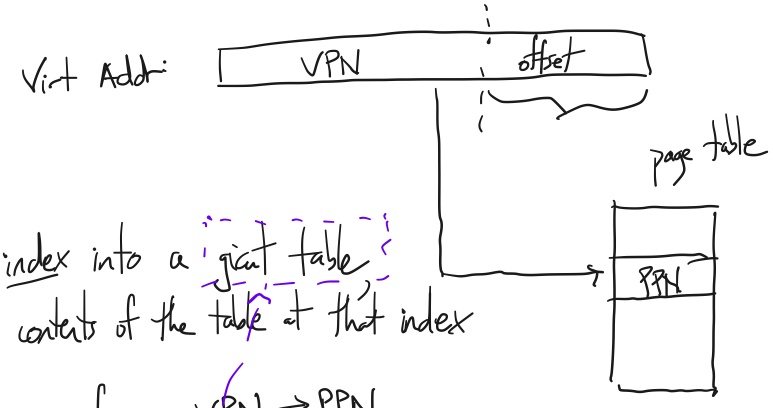
1. Last time



- purpose of virtual memory
- central mechanism: page table
- idealized page table: VPN is an index into a giant table, PPN is the contents of the table at that index
- thus, page table implements a map from VPN → PPN
- in reality, it's a map from VPN → PPN ∪ {∅}, because a VPN might not have a valid mapping in the page table
- NOTE: VPN + PPN do not necessarily have the same number of bits.
- Because the table would be gigantic, it's not materialized as a linear table. Instead, the architecture specifies multilevel page tables

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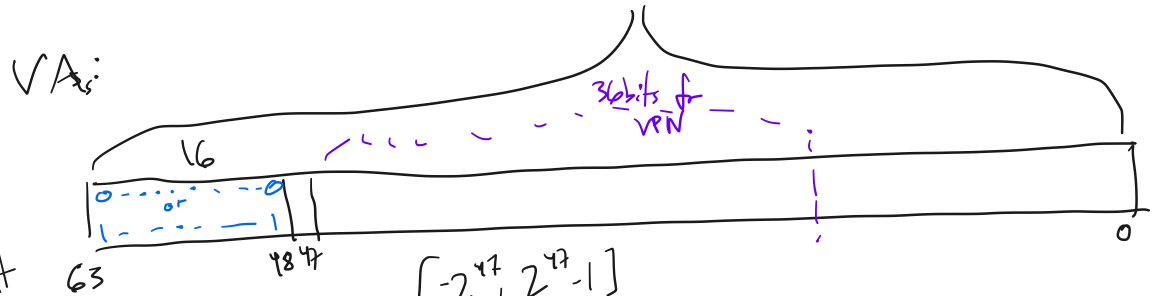
2. x86-64: addresses

movq 0x20000, %rax

0x0...0



"byte-addressed"

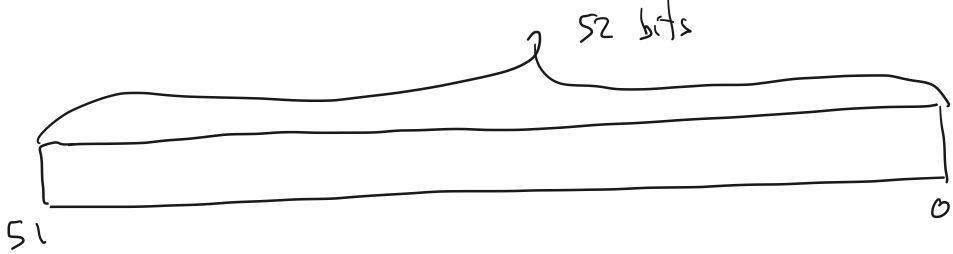
$[-2^{47}, 2^{47}-1]$

Address space has 48 usable bits. 2^{48} possible addresses (each addresses a byte).

Thus, 256 TB

$2^{48} = 2^{40} \cdot 2^8 = \text{tera} \times 256 = 256\text{TB}$

PAs:



Physical memory can be addressed w/ up to 52 bits.
How much physical memory can thus be supported?
4 PB $2^{50} \cdot 2^2 =$ Peta × 4

Mapping: going from 48-bit number (VA) to 52-bit number (PA)
at the granularity of ranges of 2^{12}.
So it's really a mapping from 36-bit numbers to 40-bit numbers.

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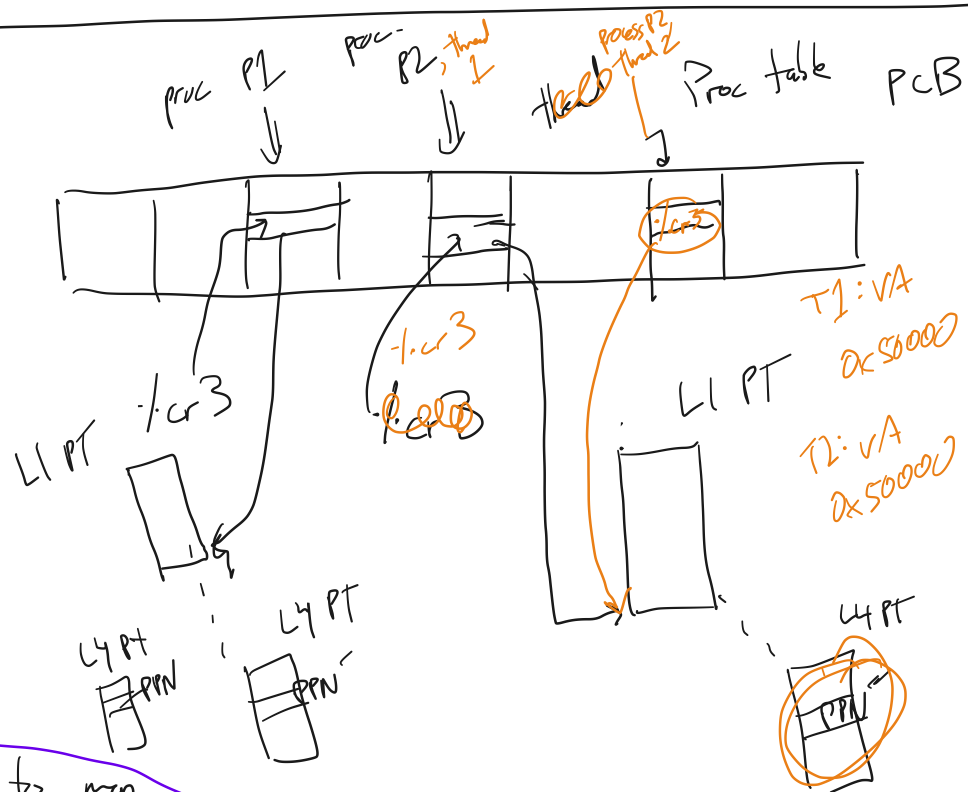
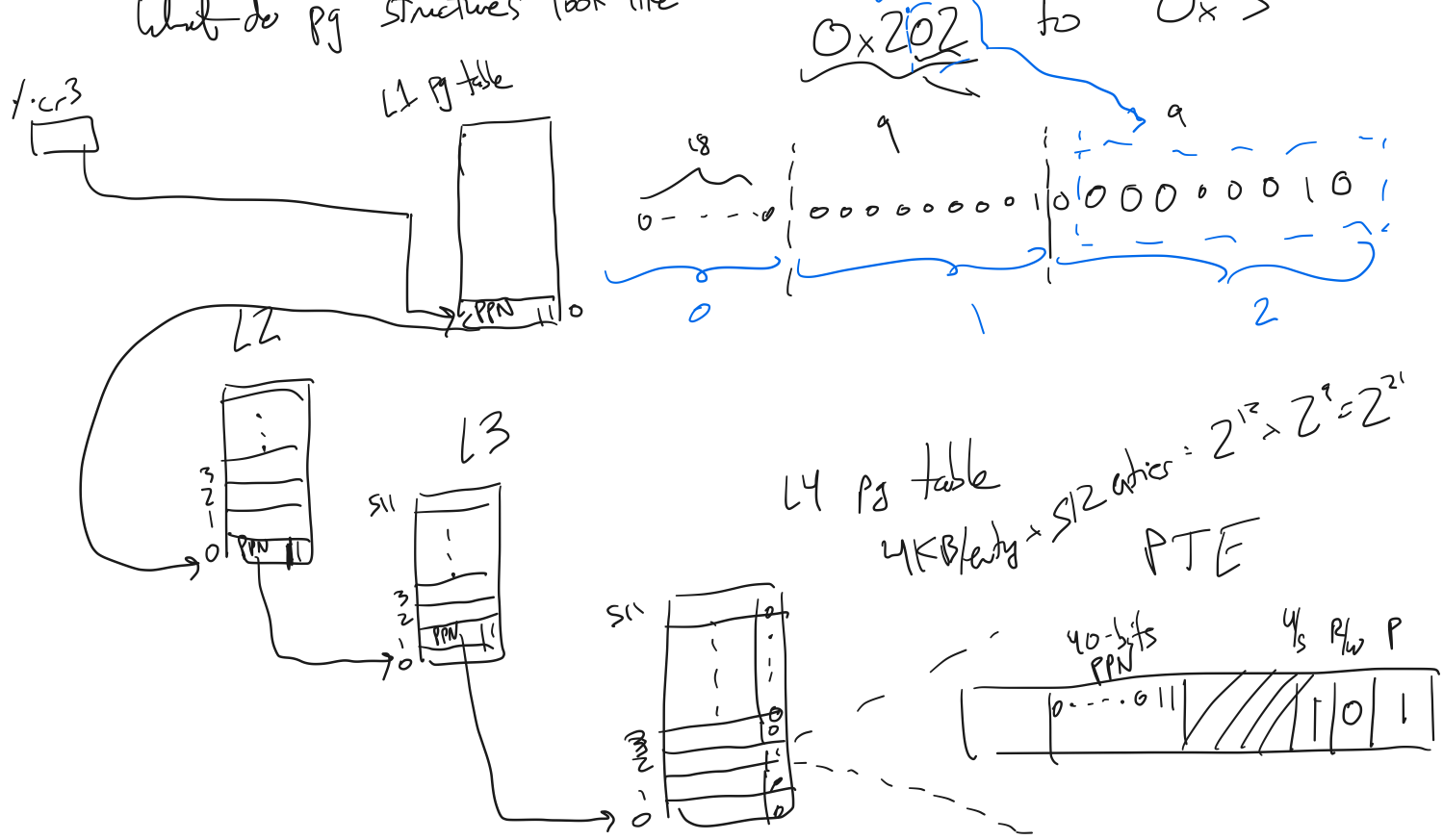
3. Page table structures.

[see handout]

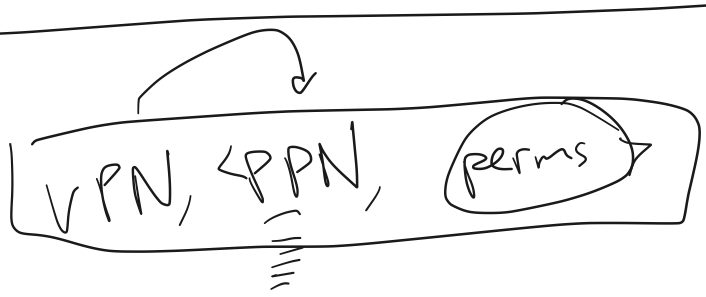


Exercise: OS wants to map
VA 0x0202000 to
PA 0x3000
and make it accessible to user-level but read-only

VPN PPN



4. TLB



H/w - managed x86, ARM

S/w - managed MIPS

TLB miss $\xrightarrow{?}$ Pg fault (No!)

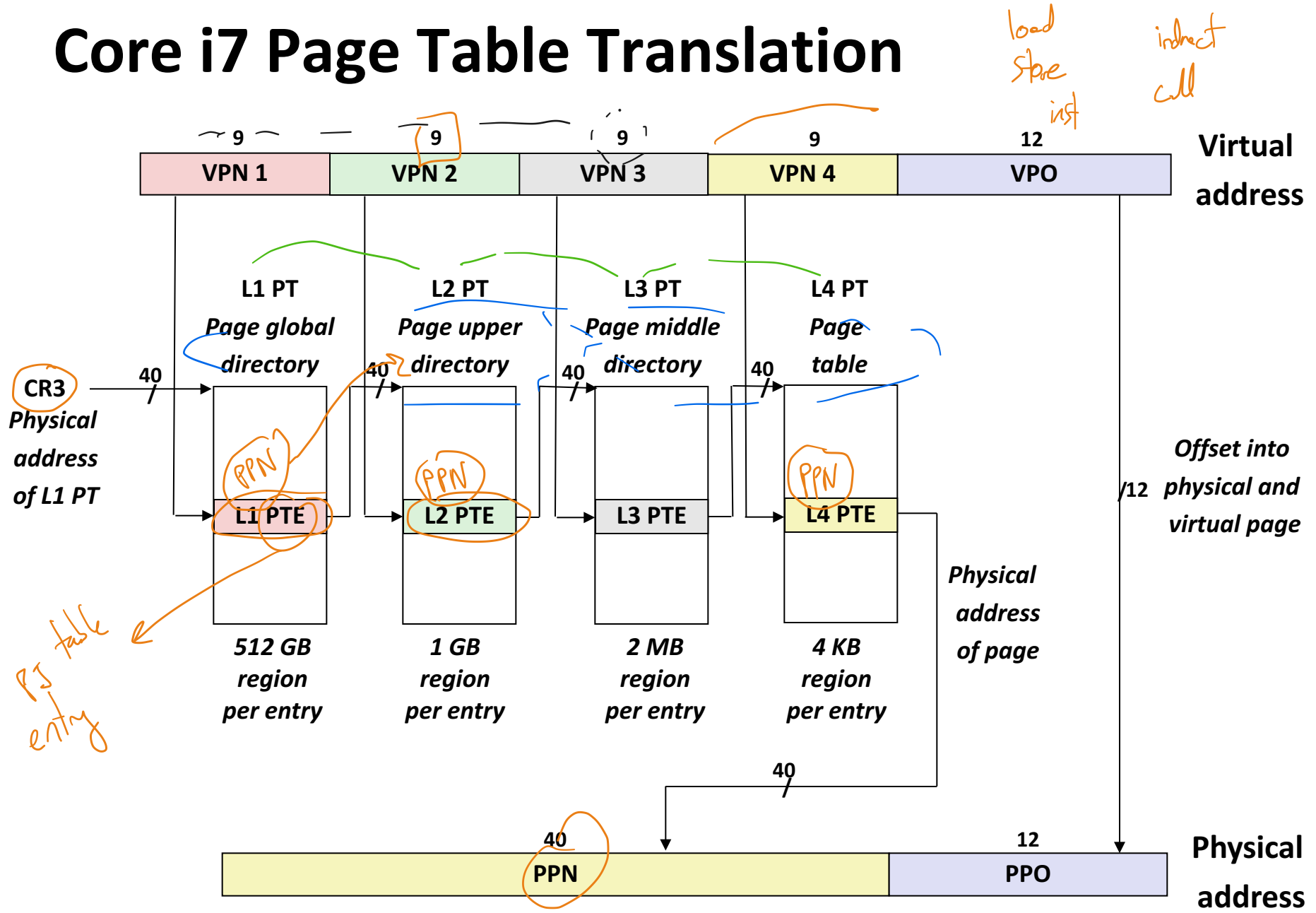
Pg fault $\xrightarrow{?}$ TLB miss (No!)

entry could be Read-only

Generalization/Simplification:

when %cr3 loaded, TLB flushed

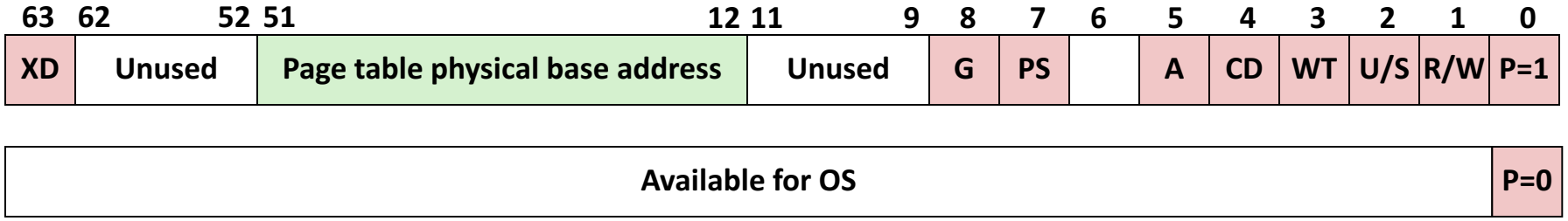
Core i7 Page Table Translation



Review of Symbols

- **Basic Parameters**
  - **N = 2^n** : Number of addresses in virtual address space
  - **M = 2^m** : Number of addresses in physical address space
  - **P = 2^p** : Page size (bytes)
- **Components of the virtual address (VA)**
  - **TLBI**: TLB index
  - **TLBT**: TLB tag
  - **VPO**: Virtual page offset
  - **VPN**: Virtual page number
- **Components of the physical address (PA)**
  - **PPO**: Physical page offset (same as VPO)
  - **PPN:** Physical page number
  - **CO**: Byte offset within cache line
  - **CI:** Cache index
  - **CT**: Cache tag

Core i7 Level 1-3 Page Table Entries



Each entry references a 4K child page table. Significant fields:

P: Child page table present in physical memory (1) or not (0).

R/W: Read-only or read-write access access permission for all reachable pages.

U/S: user or supervisor (kernel) mode access permission for all reachable pages.

WT: Write-through or write-back cache policy for the child page table.

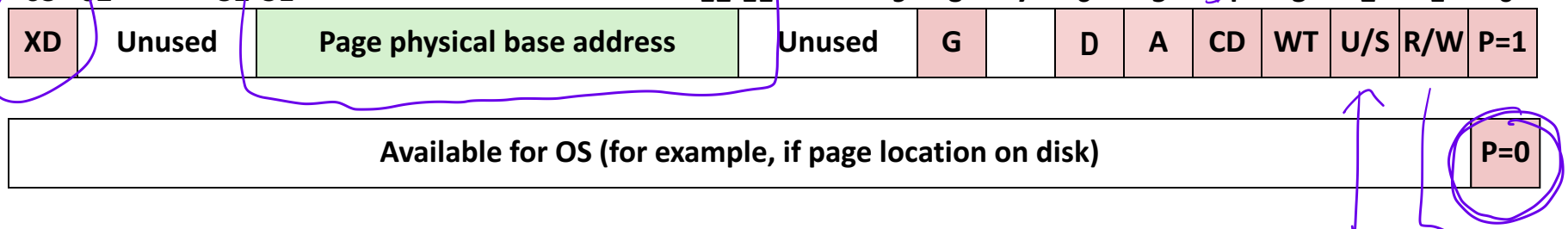
A: Reference bit (set by MMU on reads and writes, cleared by software).

PS: Page size: if bit set, we have 2 MB or 1 GB pages (bit can be set in Level 2 and 3 PTEs only).

Page table physical base address: 40 most significant bits of physical page table address (forces page tables to be 4KB aligned)

XD: Disable or enable instruction fetches from all pages reachable from this PTE.

Core i7 Level 4 Page Table Entries



Each entry references a 4K child page. Significant fields:

P: Child page is present in memory (1) or not (0)

R/W: Read-only or read-write access permission for this page

U/S: User or supervisor mode access

WT: Write-through or write-back cache policy for this page

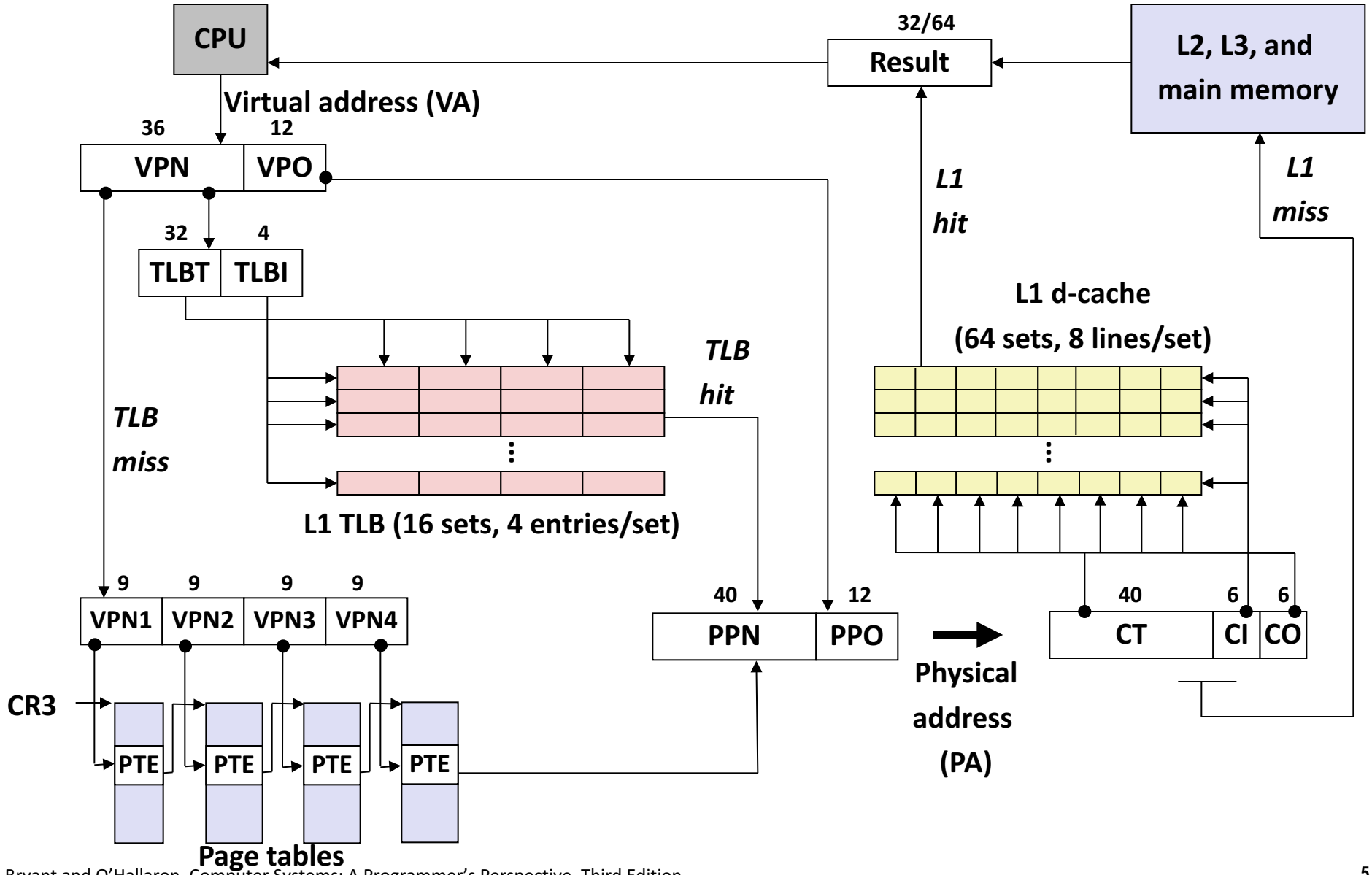
A: Reference bit (set by MMU on reads and writes, cleared by software)

D: Dirty bit (set by MMU on writes, cleared by software)

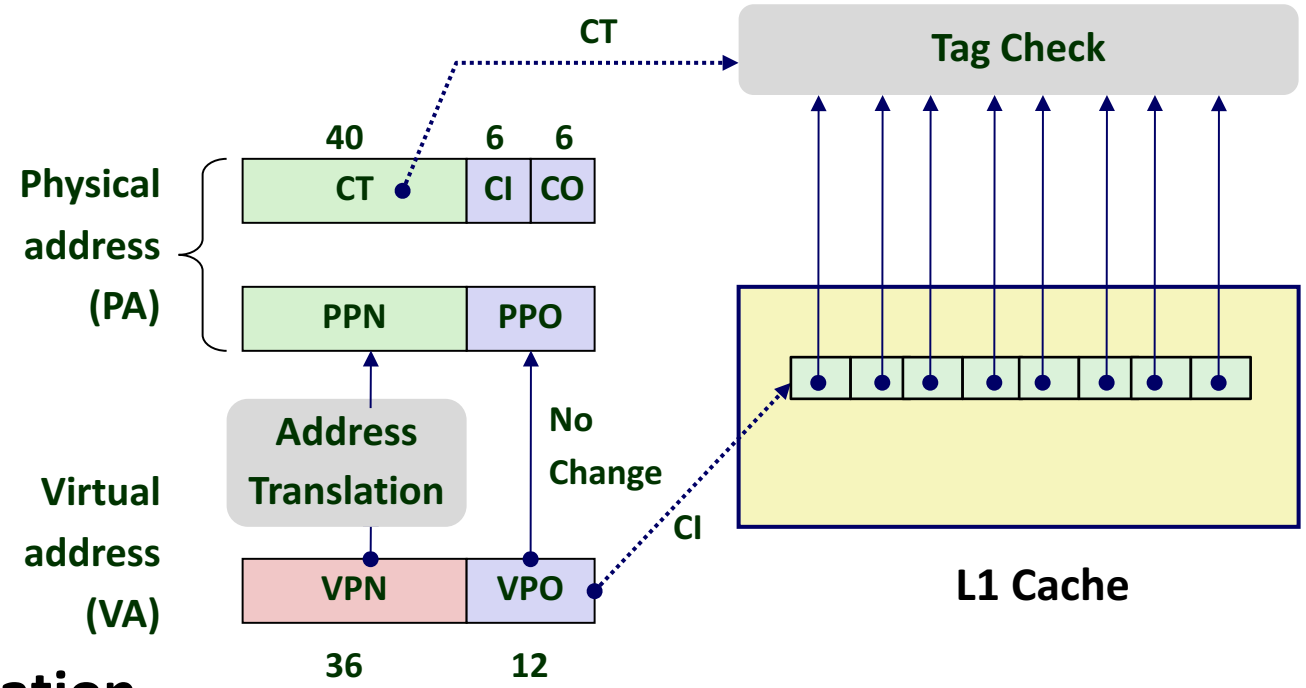
Page physical base address: 40 most significant bits of physical page address (forces pages to be 4KB aligned)

XD: Disable or enable instruction fetches from this page.

End-to-end Core i7 Address Translation

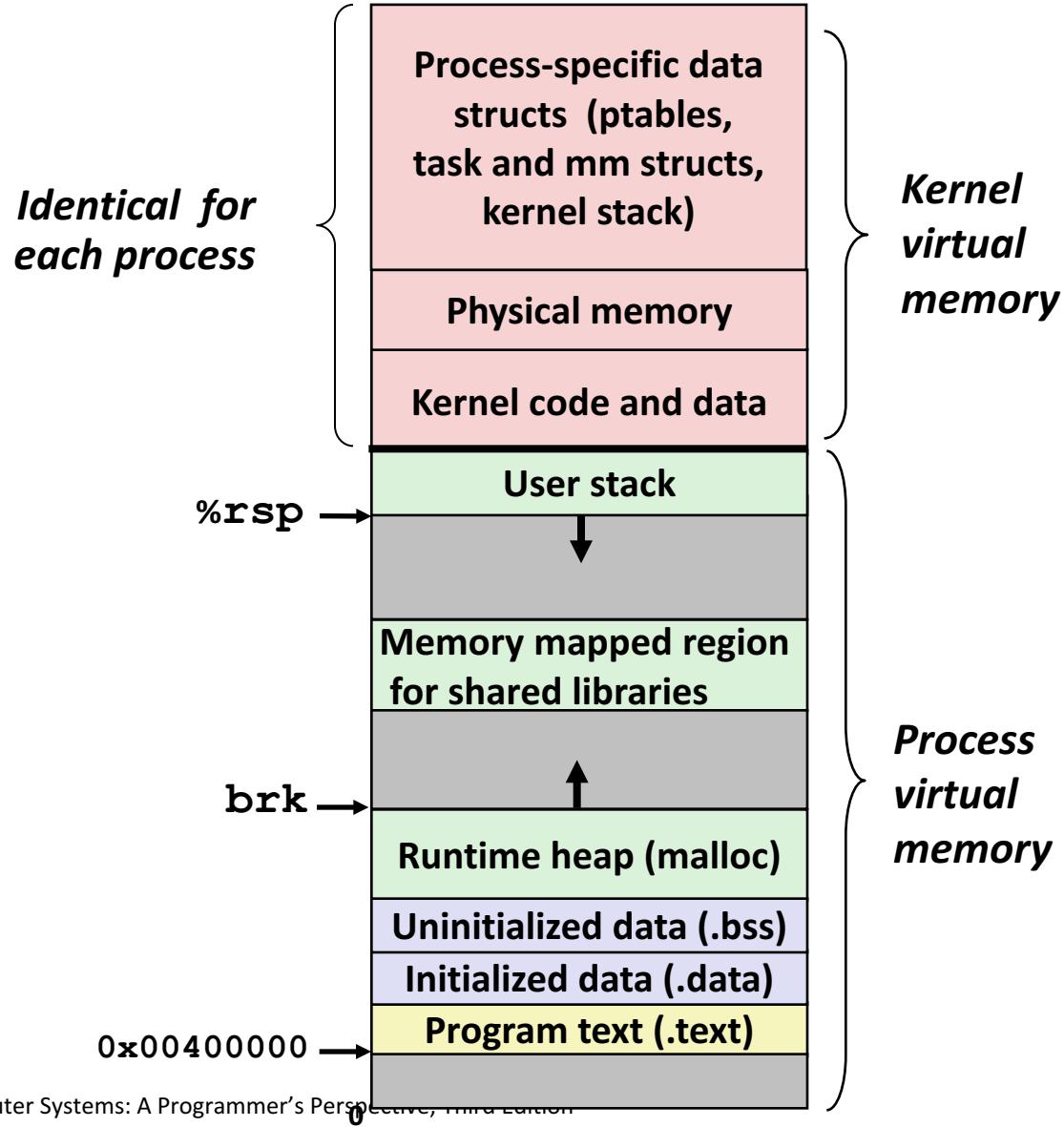


Cute Trick for Speeding Up L1 Access



- **Observation**
  - Bits that determine CI identical in virtual and physical address
  - Can index into cache while address translation taking place
  - Cache carefully sized to make this possible: 64 sets, 64-byte cache blocks
  - Means 6 bits for cache index, 6 for *cache* offset
  - That's 12 bits; matches *VPO, PPO* → One reason pages are 2^{12} bits = 4 KB

Virtual Address Space of a Linux Process



P	0	The fault was caused by a non-present page.
	1	The fault was caused by a page-level protection violation.
W/R	0	The access causing the fault was a read.
	1	The access causing the fault was a write.
U/S	0	A supervisor-mode access caused the fault.
	1	A user-mode access caused the fault.
RSVD	0	The fault was not caused by reserved bit violation.
	1	The fault was caused by a reserved bit set to 1 in some paging-structure entry.
I/D	0	The fault was not caused by an instruction fetch.
	1	The fault was caused by an instruction fetch.
PK	0	The fault was not caused by protection keys.
	1	There was a protection-key violation.
SGX	0	The fault is not related to SGX.
	1	The fault resulted from violation of SGX-specific access-control requirements.

Figure 4-12. Page-Fault Error Code