

- 1. Last time
- 2. Process's view of memory (and registers)
- 3. Stack frames
- 4. System calls
- 5. (if time) Process/os control transfers

{ next time}

Today: use the "process's view of the world" to:

- demystify functional scope
- demystify pointers
- describe how programmers get work out of the system

{ next time}

## 2. Process's view of memory and registers

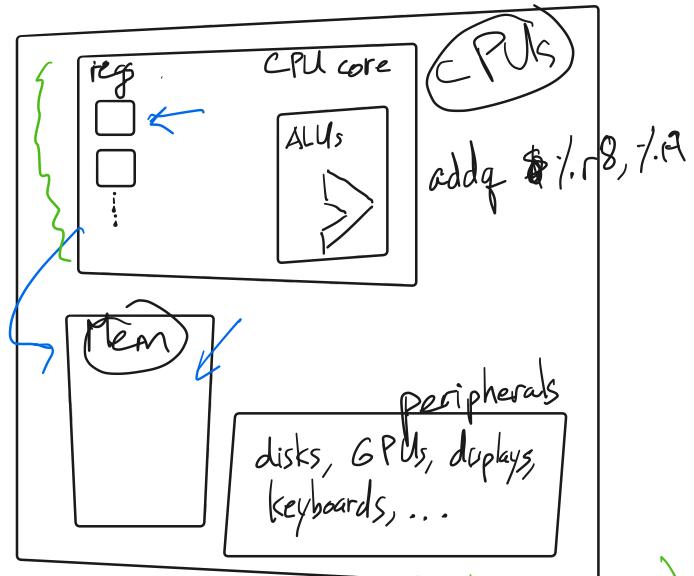
Background:

Registers (x86-64 architecture):

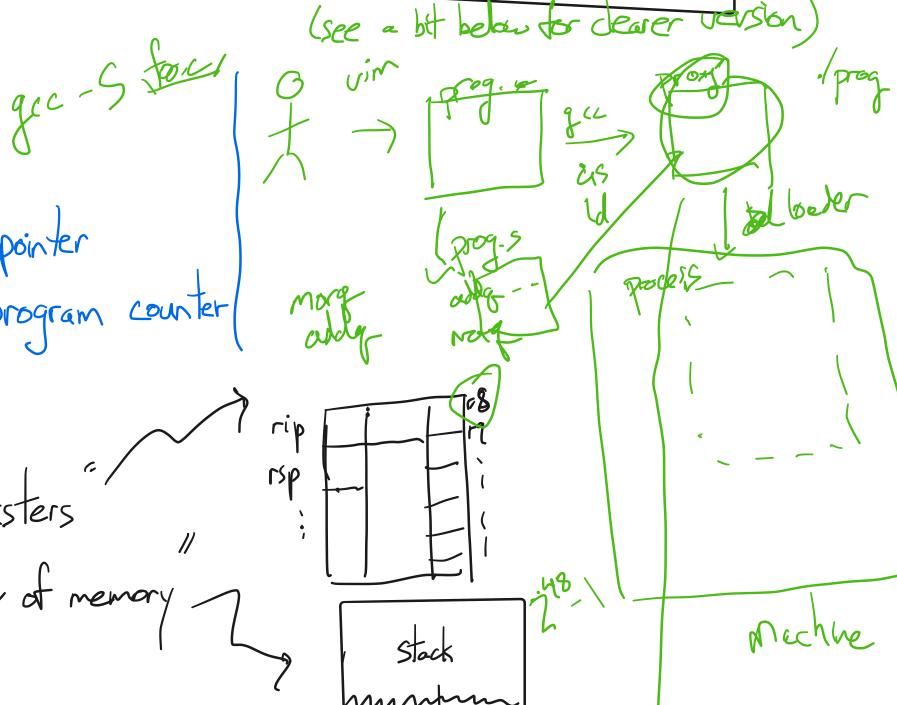
general-purpose:

`%rax, %rbx, %rcx, %rdx, %rsi, %rdi,  
%r8 - %r15, %rsp, %rbp`

special-purpose:

`%rip`

Three special registers:

`%rsp`: stack pointer`%rbp`: base pointer, or frame pointer`%rip`: instruction pointer, or program counter

Three aspects to a process:

(i) "each process has its own registers"

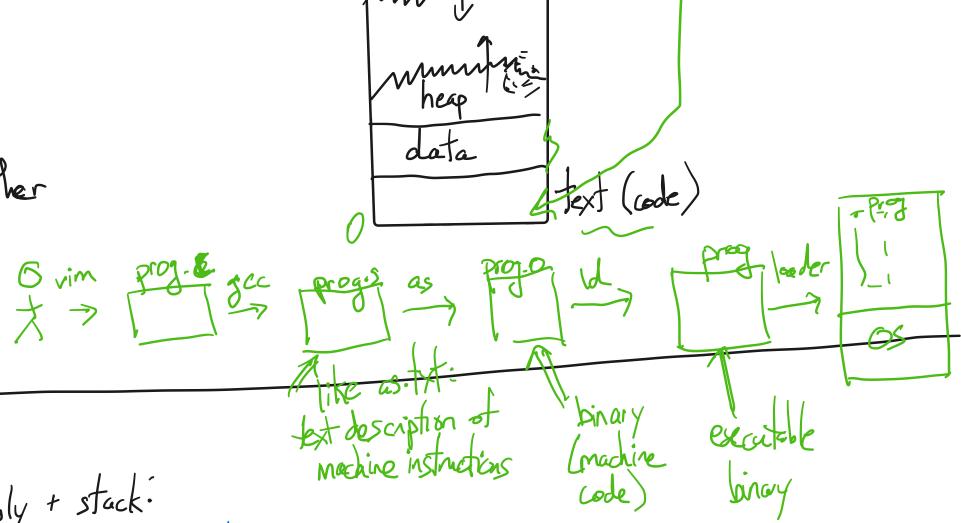
(ii) "each process has its own view of memory"

(iii) very little else needed!

some associated info:

- signal state

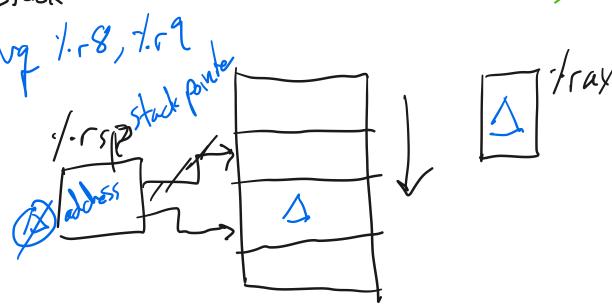
- (UI), signal mask, whether being debugged



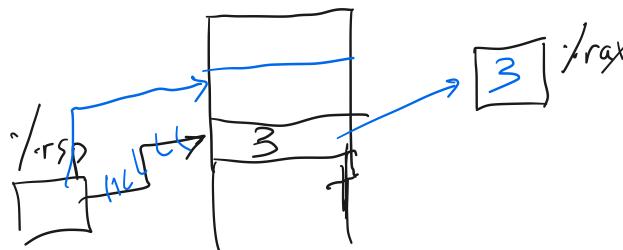
### 3. stack frames

crash course in x86-64 assembly + stack:

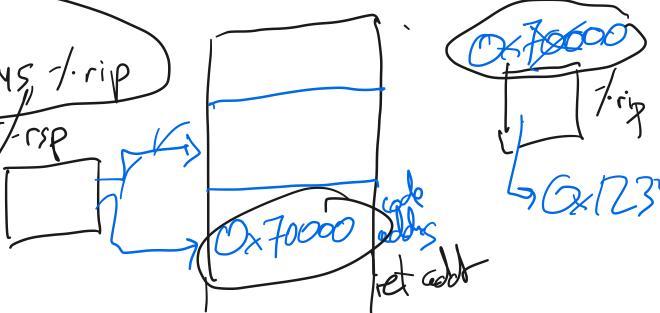
`movq FROM, TO`      `movq %rax, %rsp`  
  `pushq %rax`       $\equiv$  `subq $8, %rsp`  
                         $\equiv$  `movq %rax, (%rsp)`



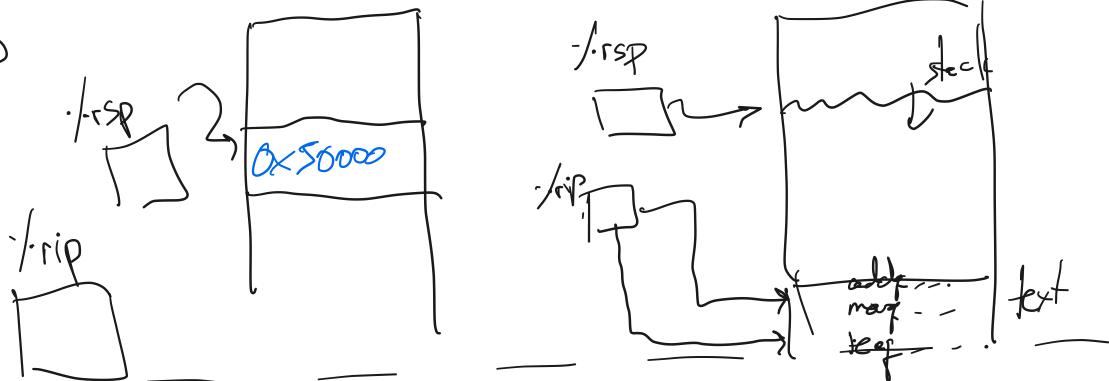
`popq %rax`       $\equiv$  `movq (%rsp), %rax`  
  `addq $8, %rsp`



`call 0x12345`       $\equiv$  `pushq %rip`  
                        `movq $0x12345, %rip`



`ret`  $\equiv$  `popq %rip`



main()

main():

`pushq %rbp`

?

f();  
:  
:



marq %rsp, %rbp

# push call-clobbered, to save them

pushq %r8

pushq %r9

call f

# restore call-clobbered

popq %r9

popq %r8

next time

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## example.c

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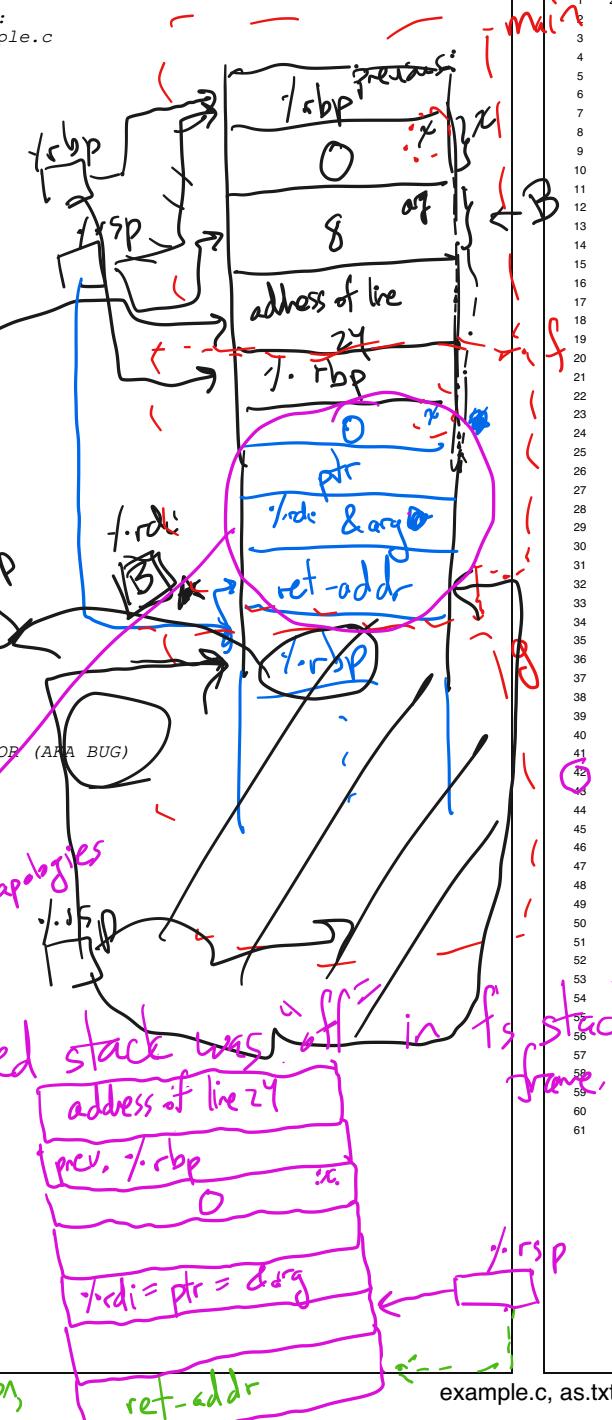
```

1  /* CS202 -- handout 1
2   * compile and run this code with:
3   * $ gcc -g -Wall -o example example.c
4   * $ ./example
5   *
6   * examine its assembly with:
7   * $ gcc -O0 -S example.c
8   * $ [editor] example.s
9  */
10
11 #include <stdio.h>
12 #include <stdint.h>
13
14 uint64_t f(uint64_t* ptr);
15 uint64_t g(uint64_t a);
16 uint64_t* q;
17
18 int main(void)
19 {
20     uint64_t x = 0;
21     uint64_t arg = 8;
22     x = f(&arg);
23
24     printf("x: %lu\n", x);
25     printf("dereference q: %lu\n", *q);
26
27     return 0;
28 }
29
30 uint64_t f(uint64_t* ptr)
31 {
32     uint64_t x = 0;
33     x = g(*ptr);
34     return x + 1;
35 }
36
37 uint64_t g(uint64_t a)
38 {
39     uint64_t x = 2*a;
40     q = &x; // <-- THIS IS AN ERROR (AKA BUG)
41     return x;
42 }

```

*h ( )  
y = \*q -*

*Correction: the depicted stack was off in f's stack frame.  
It should be as follows:  
(this is the state  
of the stack just  
before line 42 in  
as.txt).*



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## as.txt

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```

1  2. A look at the assembly...
2
3   To see the assembly code that the C compiler (gcc) produces:
4   $ gcc -O0 -S example.c
5   (then look at example.s.)
6   NOTE: what we show below is not exactly what gcc produces. We have
7   simplified, omitted, and modified certain things.
8
9
10 . main:
11    pushq  %rbp
12    movq   %rsp, %rbp
13
14    subq   $16, %rsp
15
16    movq   $0, -8(%rbp)
17    movq   $8, -16(%rbp)
18    leaq   -16(%rbp), %rdi
19
20    call   f
21
22    movq   %rax, -8(%rbp)
23
24    # eliding the rest of main()
25
26
27 f:
28    pushq  %rbp
29    movq   %rsp, %rbp
30
31    subq   $32, %rsp
32    movq   %rdi, -24(%rbp)
33
34    movq   $0, -8(%rbp)
35
36    movq   -24(%rbp), %r8
37    movq   (%r8), %r9
38    movq   %r9, %rdi
39
40
41    call   g
42
43    movq   %rax, -8(%rbp)
44    movq   -8(%rbp), %r10
45    addq   $1, %r10
46    movq   %r10, %rax
47
48
49    movq   %rbp, %rsp
50    popq   %rbp
51    ret
52
53
54 g:
55    pushq  %rbp
56    movq   %rsp, %rbp
57
58    ....
59
60    movq   %rbp, %rsp
61    popq   %rbp
62    ret

```

