

Mehryar Mohri
 Advanced Machine Learning 2024
 Courant Institute of Mathematical Sciences
 Homework assignment 1
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1 Short proof of Hoeffding's lemma

Let X be a random variable with $\mathbb{E}[X] = 0$ and $|X| \leq 1$. Show that for any $t > 0$, $\mathbb{E}[e^{tX}] \leq e^{t^2/2}$. *Hint:* use the convexity of exponential to derive $\mathbb{E}[e^{tX}] \leq \cosh(t)$.

Solution: By convexity of the exponential function, for any $t > 0$,

$$e^{tX} \leq \frac{1-X}{2}e^{-t} + \frac{1+X}{2}e^t,$$

thus, using $(2n)! = (2n)(2n-1)\cdots \geq (2n)(2(n-1))\cdots = 2^n n!$, we have

$$\mathbb{E}[e^{tX}] \leq \frac{1}{2}(e^{-t} + e^t) = \cosh(t) = \sum_{n=0}^{+\infty} \frac{t^{2n}}{(2n)!} \leq \sum_{n=0}^{+\infty} \frac{t^{2n}}{2^n n!} = e^{t^2/2}.$$

□

2 Small loss bound

In lecture 3, we used the bound on the regret of Randomized Weighted Majority (RWM): $R_T \leq O(\sqrt{L_T^{\min} \log N})$, assuming $L_T^{\min} \neq 0$. Here, you are asked to give a proof. Using the proof given in Lecture 1 for the regret of RWM and the inequality $\mathcal{L}_T \leq \frac{\log N}{1-\beta} + (2-\beta)\mathcal{L}_T^{\min}$, show that for a suitable choice of β , we have $R_T \leq 4\sqrt{\mathcal{L}_T^{\min} \log N}$.

Solution: In view of the inequality, we can write

$$R_T \leq \frac{\log N}{1-\beta} + (2-\beta)\mathcal{L}_T^{\min},$$

for $\beta \in [1/2, 1)$, that is $(1-\beta) \in (0, 1/2]$. The right-hand side is a convex function of $(1-\beta)$. Assuming $\mathcal{L}_T^{\min} \neq 0$, the minimizer over \mathbb{R}_+ is $\sqrt{\frac{\log N}{\mathcal{L}_T^{\min}}}$,

thus the minimizer over $(0, 1/2]$ is achieved for $1 - \beta = \min\left\{\frac{1}{2}, \sqrt{\frac{\log N}{\mathcal{L}_T^{\min}}}\right\}$.

Plugging in this value gives: $R_T \leq 4\sqrt{\mathcal{L}_T^{\min} \log N}$. \square

3 Weighted online-to-batch

Let ℓ be a loss function convex with respect to its first argument and bounded by one. Let h_1, \dots, h_T be the hypotheses returned by an on-line learning algorithm \mathcal{A} with regret R_T when sequentially processing $(x_t, y_t)_{t=1}^T$, drawn i.i.d. according to some distribution \mathcal{D} .

1. Fix some arbitrary non-negative weights q_1, \dots, q_T summing to one. Then, show that with probability at least $1 - \delta$, the hypothesis $h = \sum_{t=1}^T q_t h_t$ satisfies each of the following inequalities:

$$\begin{aligned} \mathbb{E}_{(x,y) \sim \mathcal{D}}[\ell(h(x), y)] &\leq \sum_{t=1}^T q_t \ell(h_t(x_t), y_t) + \|q\|_2 \sqrt{2 \log(1/\delta)} \\ \mathbb{E}_{(x,y) \sim \mathcal{D}}[\ell(h(x), y)] &\leq \inf_{h \in \mathcal{H}} \mathbb{E}_{(x,y) \sim \mathcal{D}}[\ell(h(x), y)] + \frac{R_T}{T} \\ &\quad + \|q - u\|_1 + 2\|q\|_2 \sqrt{2 \log(1/\delta)}, \end{aligned}$$

where q is the vector with components q_t and u the uniform vector with all components equal to $1/T$.

Solution: Let $R(h_t) = \mathbb{E}_{(x,y) \sim \mathcal{D}}[\ell(h_t(x), y)]$. Let Z_t be the random variable defined by $Z_t = q_t \ell(h_t(x_t), y_t) - q_t R(h_t)$ and let \mathcal{F}_t denote the filtration associated to the sample process. We have $|Z_t| \leq q_t$ and $\mathbb{E}[Z_t | \mathcal{F}_{t-1}] = \mathbb{E}[q_t \ell(h_t(x_t), y_t) | h_t] - q_t R(h_t) = q_t R(h_t) - q_t R(h_t) = 0$. Thus, by Azuma's inequality, for any $\delta > 0$, with probability at least $1 - \delta$, the following holds:

$$\sum_{t=1}^T q_t R(h_t) \leq \sum_{t=1}^T q_t \ell(h_t(x_t), y_t) + \|q\|_2 \sqrt{2 \log(1/\delta)} \quad (1)$$

$$\sum_{t=1}^T q_t \ell(h_t(x_t), y_t) \leq \sum_{t=1}^T q_t R(h_t) + \|q\|_2 \sqrt{2 \log(1/\delta)} \quad (2)$$

Since ℓ is convex with respect to its first argument, by Jensen's inequality, we have $\mathbb{E}_{(x,y) \sim \mathcal{D}}[\ell(h(x), y)] = \mathbb{E}_{(x,y) \sim \mathcal{D}}[\ell(\sum_{t=1}^T q_t h_t(x), y)] \leq \sum_{t=1}^T q_t R(h_t)$. Thus, by (1), for any $\delta > 0$, the following holds with probability at least $1 - \delta$,

$$\mathbb{E}_{(x,y) \sim \mathcal{D}}[\ell(h(x), y)] \leq \sum_{t=1}^T q_t \ell(h_t(x_t), y_t) + \|q\|_2 \sqrt{2 \log(1/\delta)}. \quad (3)$$

This proves the first statement. To prove the second claim, we will bound the empirical error in terms of the regret. For any $h^* \in \mathcal{H}$, we can write using $\inf_{h \in \mathcal{H}} \frac{1}{T} \sum_{t=1}^T \ell(h(x_t), y_t) \leq \frac{1}{T} \sum_{t=1}^T \ell(h^*(x_t), y_t)$:

$$\begin{aligned}
& \sum_{t=1}^T q_t \ell(h_t(x_t), y_t) - \sum_{t=1}^T q_t \ell(h^*(x_t), y_t) \\
&= \sum_{t=1}^T \left(q_t - \frac{1}{T} \right) [\ell(h_t(x_t), y_t) - \ell(h^*(x_t), y_t)] + \frac{1}{T} \sum_{t=1}^T [\ell(h_t(x_t), y_t) - \ell(h^*(x_t), y_t)] \\
&\leq \|q - u\|_1 + \frac{1}{T} \sum_{t=1}^T \ell(h_t(x_t), y_t) - \inf_{h \in \mathcal{H}} \frac{1}{T} \sum_{t=1}^T \ell(h(x_t), y_t) \\
&\leq \|q - u\|_1 + \frac{R_T}{T}
\end{aligned}$$

Now, by definition of the infimum, for any $\epsilon > 0$, there exists $h^* \in \mathcal{H}$ such that $\mathbb{E}_{(x,y) \sim \mathcal{D}}[\ell(h^*(x), y)] \leq \inf_{h \in \mathcal{H}} \mathbb{E}_{(x,y) \sim \mathcal{D}}[\ell(h(x), y)] + \epsilon$. For that choice of h^* , in view of (3), with probability at least $1 - \delta/2$, the following holds:

$$\mathbb{E}_{(x,y) \sim \mathcal{D}}[\ell(h(x), y)] \leq \sum_{t=1}^T q_t \ell(h^*(x_t), y_t) + \|q - u\|_1 + \frac{R_T}{T} + \|q\|_2 \sqrt{2 \log(1/\delta)}.$$

By (2), for any $\delta > 0$, with probability at least $1 - \delta/2$,

$$\sum_{t=1}^T q_t \ell(h^*(x_t), y_t) \leq \mathbb{E}_{(x,y) \sim \mathcal{D}}[\ell(h^*(x), y)] + \|q\|_2 \sqrt{2 \log(1/\delta)}.$$

Combining these last two inequalities, by the union bound, with probability at least $1 - \delta$, the following holds:

$$\begin{aligned}
\mathbb{E}_{(x,y) \sim \mathcal{D}}[\ell(h(x), y)] &\leq \inf_{h \in \mathcal{H}} \mathbb{E}_{(x,y) \sim \mathcal{D}}[\ell(h(x), y)] + \epsilon + \frac{R_T}{T} \\
&\quad + \|q - u\|_1 + 2\|q\|_2 \sqrt{2 \log(1/\delta)}.
\end{aligned}$$

Since this inequality holds for all $\epsilon > 0$, it implies the second statement. \square

2. Here, we seek to prove a bound that holds uniformly for all weight vectors q in some set. To do so, we consider a weight vector p that serves as a *reference*. A natural reference in this context could be for example the uniform distribution. Show that, for any $\delta > 0$, the following holds with

probability at least $1 - \delta$ for all $q \in \{q: \|q - p\|_1 < 1\}$:

$$\begin{aligned} \mathbb{E}_{(x,y) \sim \mathcal{D}} [\ell(h(x), y)] &\leq \sum_{t=1}^T q_t \ell(h_t(x_t), y_t) + 2\|q - p\|_1 \\ &\quad + (\|q\|_2 + 2\|q - p\|_1) \left[2\sqrt{\log \log_2 \frac{2}{1 - \|q - p\|_1}} + \sqrt{2 \log \frac{2}{\delta}} \right]. \end{aligned}$$

Hint: consider the first inequality proven above for a fixed weight vector q^k and approximation error ϵ_k , for any $k \geq 0$. Show that the inequality can be extended to hold uniformly for all $k \geq 0$ if you choose $\epsilon_k = \epsilon + \sqrt{2 \log(k+1)}$.

Solution: Consider two sequences $(\epsilon_k)_{k \geq 0}$ and $(q^k)_{k \geq 0}$. By the first part, for any fixed $k \geq 0$, we have

$$\mathbb{P} \left[\mathbb{E}_{(x,y) \sim \mathcal{D}} [\ell(h(x), y)] > \sum_{t=1}^T q_t^k \ell(h_t(x_t), y_t) + \|q^k\|_2 \sqrt{2\epsilon_k} \right] \leq e^{-\epsilon_k^2}.$$

Choose $\epsilon_k = \epsilon + \sqrt{2 \log(k+1)}$. Then, by the union bound, we can write:

$$\begin{aligned} &\mathbb{P} \left[\exists k \geq 1: \mathbb{E}_{(x,y) \sim \mathcal{D}} [\ell(h(x), y)] > \sum_{t=1}^T q_t^k \ell(h_t(x_t), y_t) + \|q^k\|_2 \sqrt{2\epsilon_k} \right] \\ &\leq \sum_{k=0}^{+\infty} e^{-\epsilon_k^2} \leq \sum_{k=0}^{+\infty} e^{-\epsilon^2 - \log((k+1)^2)} = e^{-\epsilon^2} \sum_{k=1}^{+\infty} \frac{1}{k^2} = \frac{\pi^2}{6} e^{-\epsilon^2} \leq 2e^{-\epsilon^2}. \end{aligned} \quad (4)$$

We can choose q^k such that $\|q^k - p\|_1 = 1 - \frac{1}{2^k}$. Then, for any $q \in \{q: \|q - p\|_1 < 1\}$, there exists $k \geq 0$ such that $\|q^k - p\|_1 \leq \|q - p\|_1 < \|q^{k+1} - p\|_1$ and thus such that

$$\begin{aligned} \sqrt{2 \log(k+1)} &= \sqrt{2 \log \log_2 \frac{1}{1 - \|q^{k+1} - p\|_1}} = \sqrt{2 \log \log_2 \frac{2}{1 - \|q^k - p\|_1}} \\ &\leq \sqrt{2 \log \log_2 \frac{2}{1 - \|q - p\|_1}} \end{aligned}$$

Furthermore, for that k , the following inequalities hold:

$$\begin{aligned}
\|q^k\|_2 &\leq \|q\|_2 + \|q^k - q\|_2 \\
&\leq \|q\|_2 + \|q^k - q\|_1 \\
&\leq \|q\|_2 + \|q^k - p\|_1 + \|q - p\|_1 \\
&\leq \|q\|_2 + \|q - p\|_1 + \|q - p\|_1 \\
&\leq \|q\|_2 + 2\|q - p\|_1.
\end{aligned}$$

and

$$\begin{aligned}
\sum_{t=1}^T q_t^k \ell(h_t(x_t), y_t) &\leq \sum_{t=1}^T q_t \ell(h_t(x_t), y_t) + \|q^k - q\|_1 \\
&\leq \sum_{t=1}^T q_t \ell(h_t(x_t), y_t) + 2\|q - p\|_1.
\end{aligned}$$

Plugging in these inequalities in (4) concludes the proof. \square

4 Coarse correlated equilibrium

Consider a finite normal form game with $p < +\infty$ players and finite action sets \mathcal{A}_k , $k \in [1, p]$. Show that if each player plays an external regret minimization strategy that has regret at most ϵ , then the empirical average of the players (product) distributions: $\mathbf{p} = \frac{1}{T} \sum_{t=1}^T \mathbf{p}^t$, where $\mathbf{p}^t = \prod_{k=1}^p \mathbf{p}_k^t$, is an ϵ -approximate coarse correlated equilibrium, that is, for all $k \in [1, p]$, for all $a_k \in \mathcal{A}_k$ and all $a'_k \in \mathcal{A}_k$,

$$\mathbb{E}_{\mathbf{a} \sim \mathbf{p}} [u_k(a'_k, a_{-k})] \leq \mathbb{E}_{\mathbf{a} \sim \mathbf{p}} [u_k(a_k, a_{-k})] + \epsilon.$$

Solution: By definition, for each player $k \in [1, p]$, the following holds:

$$\begin{aligned}
\mathbb{E}_{\mathbf{a} \sim \mathbf{p}} [u_k(a_k, a_{-k})] &= \frac{1}{T} \sum_{t=1}^T \mathbb{E}_{\mathbf{a} \sim \mathbf{p}^t} [u_k(a_k, a_{-k})] \\
\mathbb{E}_{\mathbf{a} \sim \mathbf{p}} [u_k(a'_k, a_{-k})] &= \frac{1}{T} \sum_{t=1}^T \mathbb{E}_{\mathbf{a} \sim \mathbf{p}^t} [u_k(a'_k, a_{-k})].
\end{aligned}$$

Since each player has regret at most ϵ , for all $k \in [1, p]$,

$$\sup_{a'_k \in \mathcal{A}_k} \mathbb{E}_{\mathbf{a} \sim \mathbf{p}^t} [u_k(a'_k, a_{-k})] \leq \mathbb{E}_{\mathbf{a} \sim \mathbf{p}^t} [u_k(a_k, a_{-k})] + \epsilon$$

This implies the ϵ -approximate coarse correlated equilibrium condition. \square