

## Graph Theory:

- a) Combinatorial Structure
- b) Algebraic Structure ~ Spectral Properties
- c) Probabilistic Structure ~

## Random Graphs & Their Evolution.

### ① Social Interactions (Pair-wise)

- ↳ Graph Theory (Interaction Choices)
- ↳ Game Theory (Strategic Choices)

We wish to model Strategic Interactions among Rational Agents.

### Ingredients:

$V \equiv$  Set of Actors

$E \subseteq V \times V \equiv$  Set of Links

$S_v \equiv$  Strategy space  $v \in V$

②  $u_v: \prod S_v \rightarrow \mathbb{R}_+$  Pay-off functions of  $v \in V$

$(V, E, S_v|_{v \in V}, u_v|_{v \in V})$

determine

A Social Network and how its agents behave.

## Defn GRAPHS (Networks)

A graph  $G = (V, E)$  consists of a set of vertices  $V$  together with a set of edges  $E \subseteq V \times V$ .

- ⇒ A mathematical object describing an irreflexive, symmetric binary relation on a discrete set, which need not be finite.

Example: Friendship.

IRREFLEXIVE: One is not his own friend.

$$\langle v, v \rangle \notin E \text{ (no self-loop)}$$

SYMMETRIC: One is a friend to a friend.

$$\langle v, w \rangle \in E \Leftrightarrow \langle w, v \rangle \in E$$

NON TRANSITIVE: One is not necessarily a friend to a friend's friend.

$$\langle u, v \rangle \in E \wedge \langle v, w \rangle \in E \not\Rightarrow \langle u, w \rangle \in E.$$

- Friendship relation in a Social Network can be described by an undirected graph.

- An edge  $e = (u, v) \in E$  (where  $E \subseteq V \times V$ ) is described by the unordered pair of vertices (players), which serve as edge's end-points.

- Two vertices  $u$  and  $v$  are adjacent if  $\exists_e e = (u, v)$  connecting  $u$  and  $v$ .

- Notation:  $|V| = n = \# \text{ vertices}$   
 $|E| = m = \# \text{ edges.}$

$$m \leq \binom{n}{2} = \frac{n(n-1)}{2} \quad \left\{ \begin{array}{l} n = \# \text{ ways to choose } u \\ n-1 = \# \text{ ways to choose } v \\ \text{identifying edge } (u, v) \equiv (v, u) \text{ [Symm]} \end{array} \right.$$

### STRICT GRAPH:

No self-loop  $(u, u) \notin E \quad \forall u$

No multi-edge  $e_1 = (u_1, v_1) \neq e_2 = (u_2, v_2)$

$\Rightarrow (u_1 = u_2) \Rightarrow (v_1 \neq v_2)$

$\wedge (u_1 = v_2) \Rightarrow (v_1 \neq u_2)$

$\forall e_1, e_2$

$\sim$

Two vertices  $u$ , and  $v$  are adjacent

$\exists_e e = (u, v)$

Two edges  $e$  and  $f$  are incident

$\exists_u e = (u, v) \wedge f = (u, w)$

$$d(v) = |\{u \mid (u, v) \in E\}| = \text{Degree}.$$

The number of vertices adjacent to a given vertex  $v$  is called the degree of the vertex.

$$\sum_{v \in V} d(v) = 2|E| = 2m$$

Average degree of the graph

$$\bar{d} = \frac{\sum_{v \in V} d(v)}{V} = \frac{2m}{n}$$

$\sim$

Density = Ratio of the number of edges to the number of possible total

$$= \frac{|E|}{\binom{n}{2}} = \frac{2m}{n(n-1)} = \frac{\sum d(v)}{n(n-1)} = \frac{\bar{d}}{n-1}$$

① Density = 1  $\Rightarrow$  Graph is complete.

$$\bar{d} = (n-1) \quad \forall v \quad d(v) \leq n-1$$

$$\Rightarrow \sum_v d(v) = n-1$$

A graph is complete if all of its vertices are adjacent to all others.

- If a social network has many "well-connected individuals" then the network is "dense," since density =  $\bar{d}/n-1$
- ⇒ Finding and connecting to "well-connected" subnetworks is beneficial.

CLIQUE                    CLUB

CLAN.

Distance between two individuals:

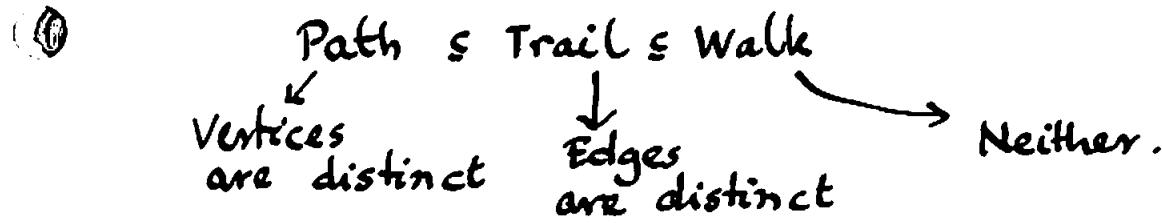
Path: A sequence of adjacent, <sup>distinct</sup> vertices

$$v_0, v_1, \dots, v_n$$

$$\forall i \quad (v_i, v_{i+1}) \in E, \quad 0 \leq i < n$$

$$\forall i, j \quad v_i \neq v_j \quad 0 \leq i, j \leq n$$

in which no vertex occurs more than once.



- The (geodesic) - distance between two vertices is the length of the shortest path connecting them.

$d(u, v)$  = Geodesic distance between  $u$  and  $v$ .

- ◊ The maximal geodesic-distance in a graph is its diameter,  $\delta(G)$

$$\forall u, v \quad d(u, v) \leq \delta(G).$$

$G_1(V_1, E_1) \subseteq G(V, E)$  (subgraph)

iff  $V_1 \subseteq V \wedge E_1 \subseteq E$ .

- ◊ A subgraph of a graph  $G$  is a graph whose vertices and edges are contained in  $G$ .

$$G_1 \subseteq G, |V_1| = k, |E_1| = \binom{k}{2}$$

- ◊ A clique is a maximal complete subgraph.

- ◊ The subgraph  $G(S)$  of a graph  $G$  induced by the set of nodes  $S$  is defined as the maximal subgraph of  $G$  that has vertex set  $S$ .