

Random Graphs:

A general model of a social network.

→ ER (Erdős-Renyi) Random Graphs.

Two Ways of Describing Random Graphs:

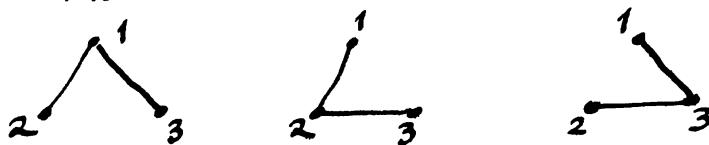
Closely related variants of ER-Random Graphs.

{ $G(n, m)$ Models:
 $G(n, p)$ Models:

(A) $G(n, m)$ Models:

A graph $G = (V, E)$ is chosen uniformly at random from the collection of all graphs, which have $|V| = n$ nodes and $|E| = m$ edges.

$G(3, 2)$ - Model.



$$P_r = \left\{ \frac{1}{3}, \frac{1}{3}, \frac{1}{3} \right\}$$

- ❖ Exactly three possible graphs on three vertices and two edges.
- ❖ Each is selected with equal probability = $\frac{1}{3}$.

(68)

(B) $G(n,p)$ Models

A graph $G = (V, E)$ is constructed by connecting every pair of nodes uniformly randomly with $\Pr[e \in E] = p$.

For every pair of vertices $u, v \in V$ an edge $e = (u, v) \in E$ is included in the graph with probability p independent from every other edge.

Equivalently,

All graphs with $|V| = n$ & $|E| = M$ have equal edge probability of

$$p = \frac{M}{\binom{n}{2}} = \text{Density of the graph.}$$

- At $p = \frac{1}{2}$, all graphs on n vertices are chosen with equal probability = $p^k (1-p)^{\binom{n}{2} - k} = 2^{-\binom{n}{2}} = 2^{-\binom{n}{2}}$

- As p increases from 0 to 1, the model produces denser graphs with higher probability (than sparser graphs).

- Expected Number of Edges:

$$\langle |E| \rangle = \binom{n}{2} p = \frac{n(n-1)p}{2}$$

- Expected Degree.

$$\bar{d} = \langle d \rangle = \frac{2 \langle |E| \rangle}{\langle |V| \rangle} = \frac{2 \binom{n}{2} p}{n} = (n-1)p$$

There are $(n-1)$ possible other vertices, of which each can be adjacent with probability $= p$.

$$d(v) \sim \text{Bin}(n-1, p)$$

$$P[d(v) = k] = \binom{n-1}{k} p^k (1-p)^{n-1-k}.$$

The degree of a vertex in a graph $G \in G(n, p)$ is distributed as a Binomial.

$$\mu[d(v)] = (n-1)p \quad \sigma^2[d(v)] = (n-1)pq.$$

~~~~~

Asymptotic Analysis:

$$|V| = n \rightarrow \infty$$

Random Graphs are often studied in the asymptotic case, as  $|V| = n$  (the number of vertices) tends to infinity.

⇒ Graphon, Graphlet.

~~~~~

If expected degree \bar{d} is held constant (independent of n)

$$\bar{d} = (n-1)p = \text{const} = \lambda$$

$$p = \frac{\lambda}{n-1}$$

$$P[d(v) = k] \approx \frac{(n-1)^{(k)}}{k!} p^k (1-p)^{n-1-k} p$$

$$\approx \frac{[(n-1)p]^k}{k!} e^{-(n-1)p}$$

$$= \frac{\lambda^k}{k!} e^{-\lambda}$$

$$d(v) \sim \text{Poisson}(\lambda)$$

$$\left. \begin{array}{l} \mu[d(v)] = \lambda \\ \sigma^2[d(v)] = \lambda \end{array} \right\} \text{Poisson Approximation.}$$

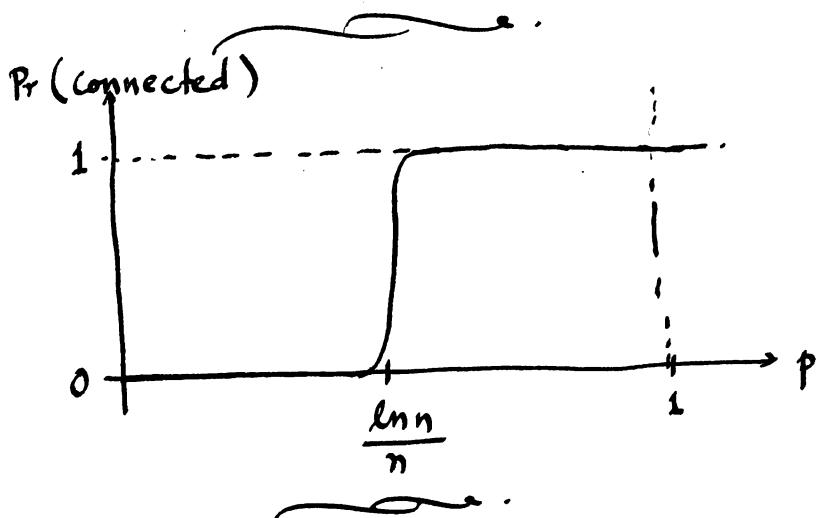
TIPPING POINT

(70)

$\left\{ \begin{array}{l} \text{Phase Transition} \\ \text{0-1 Law.} \end{array} \right.$

Small p : If $p < \frac{(1-\epsilon) \ln n}{n}$, then a graph in $G(n,p)$ will a.s. contain isolated vertices \Rightarrow DISCONNECTED.

Large p : If $p > \frac{(1+\epsilon) \ln n}{n}$, then a graph in $G(n,p)$ will a.s. be CONNECTED.



QUESTIONS ABOUT RANDOM SOCIAL NETWORK MODELS:

- 1) Does the network have isolated nodes? Cycles? Giant Component?
- 2) What are the probabilities of such events?
- 3) Asymptotic Analysis:
Compute Probabilities, as $n \rightarrow \infty$.

THRESHOLD FUNCTIONS FOR CONNECTIVITY: (Endős-Rényi 1961).

A threshold function for the connectivity of the Endős-Rényi model $G(n,p)$ is

$$t(n) = \frac{\ln n}{n}$$

(71)

For a graph $G \in G(n, \lambda \frac{\ln n}{n})$

$$\Pr[G = \text{connected}] = \begin{cases} 0, & \text{if } \lambda < 1; \\ 1, & \text{if } \lambda > 1. \end{cases}$$

$$\mathbb{1}_i = \begin{cases} 1, & \text{if node } i \text{ is isolated;} \\ 0, & \text{otherwise.} \end{cases}$$

$$\mathbb{1}_i \sim \text{Bernoulli}(\pi) \quad \left\{ \begin{array}{l} \pi = \Pr[\mathbb{1}_i = 1] = (1-p)^{n-1} \\ = (1-p)^{1/p} \cdot (n-1)p \\ = e^{-(n-1)\lambda \frac{\ln n}{n}} \\ \approx e^{-\lambda \ln n} = n^{-\lambda}. \end{array} \right.$$

$$\boxed{\mathbb{1}_i \sim \text{Bernoulli}(n^{-\lambda})} \leftrightarrow \boxed{\pi = n^{-\lambda}}$$

~~is Bernoulli~~

$$X = \sum \mathbb{1}_i = \text{Total \# of isolated nodes.}$$

~~$E[X] \approx \text{Var}[X] = \text{Var}[X]$~~

$$E[X] = n \cdot n^{-\lambda} = n^{1-\lambda} \rightarrow \begin{cases} \infty, & \text{if } \lambda < 1; \\ 0, & \text{if } \lambda > 1. \end{cases}$$

Problems

- (1) Need to show that $\Pr[X=0] = 0$, if $\lambda < 1$.
- (2) Note that $\Pr[X=0] > 0$ does not necessarily imply that the graph is connected.

(72)

Assume that $\lambda < 1$.

$$\text{Var}[x] = \sum_i \text{var}(\mathbf{1}_i) + \sum_i \sum_{j \neq i} \text{cov}(\mathbf{1}_i, \mathbf{1}_j)$$

$$= n \text{var}(\mathbf{1}_1) + n(n-1) \text{cov}(\mathbf{1}_1, \mathbf{1}_2)$$

$$\left\{ \begin{array}{l} \text{var}(\mathbf{1}_1) = \pi(1-\pi) = \pi - \pi^2 \\ \text{cov}(\mathbf{1}_1, \mathbf{1}_2) = E(\mathbf{1}_1 \mathbf{1}_2) - E(\mathbf{1}_1) E(\mathbf{1}_2) \\ = (1-p)^{2n-3} - (1-p)^{n-1} (1-p)^{n-1} \\ = \frac{\pi^2}{1-p} - \pi^2 \end{array} \right.$$

$$\text{Var}[x] = n\pi - n\pi^2 + \frac{n^2\pi^2}{1-p} - n^2\pi^2$$

$$\approx n\pi + n^2\pi^2 p$$

$$= n \cdot n^{-\lambda} + n^2 n^{-2\lambda} p$$

$$\approx n^{1-\lambda} = E[x] \quad (\text{if } \lambda < 1)$$

$$\text{Var}[x] = \int (x - E[x])^2 p(x) dx$$

$$\geq (0 - E[x])^2 \Pr[x = 0]$$

$$\Pr[x = 0] \leq \frac{\mathbb{E}[x]^2}{\text{Var}[x]} = \frac{\pi^2}{n^{1-\lambda}}$$

$$\frac{\text{Var}[x]}{E[x]^2} \approx \frac{1}{E[x]} \rightarrow 0$$

"Graph is disconnected"

$$\Rightarrow \exists v' \in V, |v'| = k \quad \text{Disconnect } [v', V \setminus v'] \\ k \leq n/2$$

$$\Pr[V' \text{ is not connected to } V \setminus V', |V'| = k] = (1-p)^{k(n-k)} \\ = (1-p)^{1/p} \frac{k(n-k)p}{e^{-\lambda k(n-k) \ln n/n}} \\ \approx e^{-\lambda k(n-k) \ln n/n}$$

$$\Pr[\exists v', |v'| = k \quad V' \text{ is not connected to } V \setminus V']$$

$$= \binom{n}{k} (1-p)^{k(n-k)} \\ = \frac{n!}{k!(n-k)!} e^{-\lambda k(n-k) \ln n/n} \\ k = \frac{n}{2} \left\{ \right. \propto \frac{\sqrt{2\pi n} \left(\frac{n}{e}\right)^n}{\pi n \left(\frac{n}{2e}\right)^{n/2} \left(\frac{n}{2e}\right)^{n/2}} e^{-\lambda \frac{n^2}{4} \ln n/n} \\ \approx \sqrt{\frac{2}{\pi n}} \cdot 2^n \cdot n^{-\lambda n/4}$$

$$\lambda > 1 \left\{ \right. \propto \sqrt{\frac{2}{\pi n}} \cdot 2^{+n - n \lg n / 4}$$

$$\Pr[\text{Graph is disconnected}]$$

$$\approx \sum_{k=1}^{n/2} \binom{n}{k} (1-p)^{k(n-k)} \\ < \frac{n}{2} \sqrt{\frac{2}{\pi n}} \cdot 2^{-[n \lg n / 4 - n]} \\ = \sqrt{\frac{n}{2\pi}} \cdot 2^{-[n \lg n / 4 - n]} \\ = \frac{1}{\sqrt{2\pi}} \cdot 2^{-[n \lg n / 4 - n + \lg n / 2]}$$

$\rightarrow 0$ as $n \rightarrow \infty$.

□