# **Tunable Automation**

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#### Verus

- Verus is a semi-automated program verifier for Rust programs
- Lots of systems and security work using Verus
  - https://verus-lang.github.io/verus/publications-and-projects/
- Recent effort at AWS to adopt Verus
- Intrinsic Verification, proof and executable code are intermixed, no need for external tools to translate and verify

## Spec/Proof/Exec Mode

```
spec fn divides(n: int, k: nat) -> bool { n % (k as int) == 0 }
spec fn is prime(n: nat) \rightarrow bool { forall|k: nat| 2 <= k < n ==> !divides(n as int, k) }
spec fn is even(i: int) -> bool { divides(i, 2) }
proof fn even gt 2 isnt prime(i: nat)
  requires i > 2 && is even(i as int)
  ensures !is prime(i) { }
fn is prime impl(n: u64) -> (result: bool)
  requires n \ge 2,
  ensures result == is prime(n as nat)
{ /* ... implementation and proof ... */}
```

#### Quantification

Quantifier Instantiation is the primary source of incompleteness, also the primary source of automation.

In auto-active theorem provers, we often use user-level *triggers* to guide quantifier instantiations.

```
spec fn is even(i: int) -> bool {
   i % 2 == 0
proof fn seq trigger example(s: Seq<int>)
   requires
                                            only instantiating for is_even(s[3]).
       5 <= s.len(),</pre>
       forall|i: int| 0 <= i < s.len() ==> #[trigger] is_even(s[i])
   assert(s[3] % 2 == 0); // FAILS
```

```
spec fn is even(i: int) -> bool {
   i % 2 == 0
proof fn seq trigger example(s: Seq<int>)
   requires
       5 <= s.len(),</pre>
       forall|i: int| 0 <= i < s.len() ==> # [trigger] is ever
                                               is even(#[trigger] s[i]),
   assert(s[3] % 2 == 0); // OK
                                               then we can instantiate the forall:
                                                       0 \le 3 \le s.len() => is even(s[3]),
```

```
spec fn is even(i: int) -> bool {
   i % 2 == 0
proof fn seq trigger example(s: Seq<int>)
   requires
       5 <= s.len(),</pre>
       forall|i: int| 0 <= i < s.len() ==> #[trigger] is even(s[i]),
   assert(is even(s[3])); // OK then we can instantiate the forall:
                                          0 <= 3 < s.len() ==> is_even(s[3]),
   assert(s[3] % 2 == 0); // OK
```

## Implicit Context

In Dafny/Verus, there's also an prelude of axioms, for example

```
proof fn seq_axiom_usage(s1: Seq<nat>, s2: Seq<nat>)
    requires
        s1.len() > 10 && s2.len() > 20,
    ensures
        s1.add(s2).len() > 30,
{}
```

#### uses this axiom:

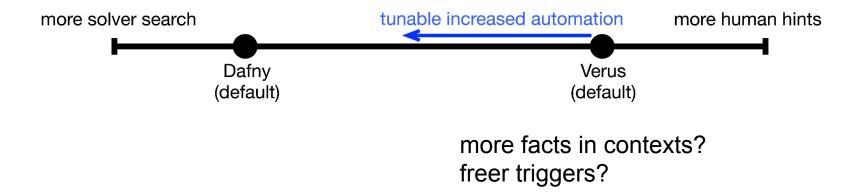
```
pub broadcast axiom fn axiom_seq_add_len<A>(s1: Seq<A>, s2: Seq<A>)
    ensures
    #[trigger] s1.add(s2).len() == s1.len() + s2.len(),
;
```

#### Verus

- Verus has a "conservative" design for verification performance, which leads to more manual proofs
  - No preconditions for spec functions (total math functions)
  - Limited Prelude
     (~280 axioms in `DafnyPrelude.bpl` vs ~96 in `vir/src/prelude.rs`, and ~111 of them in vstd)
  - No non-linear reasoning by default

All of this is because Verus wants more "native" encoding to SMT level

## Automation - Performance spectrum



A mechanism for fine-grained control of quantified facts.

You can import quantified facts at any level:

- after any `verus!` macro
- modules
- proof functions (`proof fn`)
- assert (expr) by { /\* proof \*/ }
- <u>calculational</u> proofs

```
pub proof fn push_contains(a: Seq<int>)
{
   let b = a.push(3);
   assert(b.contains(3)); // FAILS
}
```

because Verus can't infer the following from the default Verus context:

```
pub proof fn lemma_seq_contains_after_push<A>(s: Seq<A>, v: A, x: A)
    ensures
        s.push(v).contains(x) <==> v == x || s.contains(x),
{/* manual proof ... */}
```

```
pub broadcast proof fn lemma_seq_contains_after_push<A>(s: Seq<A>, v: A, x: A)
ensures
#[trigger] s.push(v).contains(x) <==> v == x || s.contains(x),
```



```
pub proof fn lemma_seq_contains_after_push<A>()
    ensures
    forall|s: Seq<A>, v: A, x: A|
    #[trigger] s.push(v).contains(x) <==> v == x || s.contains(x),
```

In VeriFast/Dafny, this kind of lemmas are named "auto" lemmas

```
pub proof fn push_contains(a: Seq<int>)
{
    broadcast use vstd::seq_lib::lemma_seq_contains_after_push;
    let b = a.push(3);
    assert(b.contains(3)); // PASSES because of the increased context
}
```

## Verus Standard Library

In 2023, we provide some sort of automation by explicitly universally quantifying the input parameters. Now it's just a broadcast group.

```
/// Properties of sequences from the Dafny prelude (which were axioms in Dafny, but proven here in Verus)
#[deprecated = "Use `broadcast use group seg properties` instead"]
pub proof fn lemma seg properties<A>()
    ensures
        forall|s: Seq<A>, x: A|
            s.contains(x) <==> exists|i: int| 0 <= i < s.len() && #[trigger] s[i] == x, //from lemma seg contains(s, x),</pre>
        forall|x: A| !(#[trigger] Seg::<A>::empty().contains(x)), //from lemma seg empty contains nothing(x),
        forallis: Seg<A>| #[trigger] s.len() == 0 ==> s =~= Seg::<A>::emptv(), //from lemma seg emptv equalitv(s),
        forallix: Seg<A>, v: Seg<A>, elt: Al #[trigger]
            (x + y).contains(elt) <==> x.contains(elt) || y.contains(elt), //from lemma_seq_concat_contains_all_elements(x, y, elt),
        forall|s: Seg<A>, v: A, x: A| #[trigger] s.push(v).contains(x) <==> v == x || s.contains(x), //from lemma_seq_contains_after_push(s, v, x)
        forall|s: Seg<A>, start: int, stop: int, x: A|
            (0 <= start <= stop <= s.len() && #[trigger] s.subrange(start, stop).contains(x)) <==> (
            exists|i: int| 0 <= start <= i < stop <= s.len() && #[trigger] s[i] == x). //from lemma seg subrange elements(s. start. stop. x).
        forall|s: Seq<A>, n: int| 0 <= n <= s.len() ==> #[trigger] s.take(n).len() == n, //from lemma_seq_take_len(s, n)
        forall|s: Seq<A>, n: int, x: A|
            (\#[trigger] s.take(n).contains(x) \&\& 0 <= n <= s.len()) <==> (exists|i: int|)
                0 \ll i \ll n \ll s.len() \&\& \#[trigger] s[i] == x), //from lemma seg take contains(s, n, x),
        forall|s: Seq<A>, n: int, j: int| 0 <= j < n <= s.len() ==> #[trigger] s.take(n)[j] == s[j], //from lemma_seq_take_index(s, n, j),
        forallis: Seq<A>, n: int| 0 <= n <= s.len() ==> #[trigger] s.skip(n).len() == s.len() - n. //from lemma seq skip len(s. n).
        forall|s: Seq<A>, n: int, x: A|
            (\#[trigger] s.skip(n).contains(x) \&\& 0 <= n <= s.len()) <==> (exists|i: int|)
                0 \ll n \ll i \ll s.len() \&\& \#[trigger] s[i] == x), //from lemma seq skip contains(s, n, x),
```

```
// include all the Dafny prelude lemmas
pub broadcast group group seg properties {
    lemma_seq_contains,
    lemma_seq_empty_contains_nothing,
    lemma_seq_empty_equality,
    lemma seg concat contains all elements.
    lemma seg contains after push,
    lemma seg subrange elements,
    lemma_seq_take_len,
    lemma_seq_take_contains,
    lemma_seq_take_index,
    lemma seg skip len.
    lemma seg skip contains,
    lemma seg skip index,
    lemma_seq_skip_index2,
    lemma_seq_append_take_skip,
    lemma_seq_take_update_commut1,
    lemma_seq_take_update_commut2,
    lemma seg skip update commut1,
    lemma seg skip update commut2,
    lemma_seq_skip_build_commut,
    lemma_seq_skip_nothing,
    lemma_seq_take_nothing,
    // Removed the following from group due to bad verification performance
    // for `lemma merge sorted with ensures`
   // lemma seg skip of skip,
    group_to_multiset_ensures,
```

```
pub proof fn push_contains(a: Seq<int>)
{
    broadcast use vstd::seq_lib::group_seq_properties;
    let b = a.push(3);
    assert(b.contains(3));
}
```

Worried about verification performance for larger proofs?

```
pub proof fn push contains(a: Seq<int>)
                             lemma seq contains after push
  broadcast use vstd::seq lib:: ______
  let b = a.push(3);
  assert(b.contains(3));
> verus broadcast.rs -V axiom-usage-info
note: checking this function used these broadcasted lemmas and broadcast groups:
       - (group) vstd::seq lib::group seq properties,
       - vstd::seq lib::lemma seq contains after push
 --> broadcast rs:5:1
5 | pub proof fn push contains(a: Seq<int>)
   ^^^^^
verification results:: 1 verified, 0 errors
```

#### **Modularized Proof Libraries**

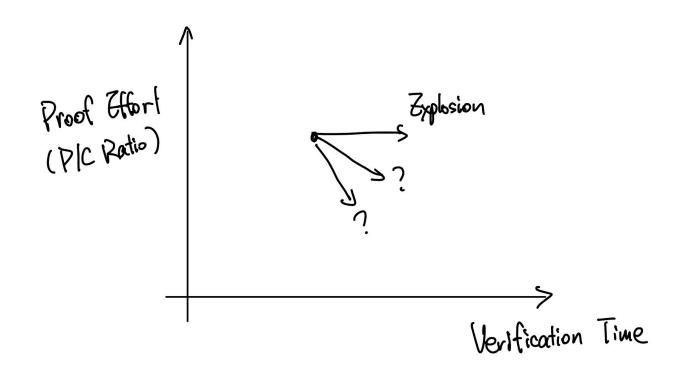
```
pub trait Key {
   . . .
   proof fn key obligations()
       ensures // ... conditions necessary for the type to be a valid key
   broadcast proof fn trans lt lt(a:Self, b:Self, c:Self)
       ensures a < b \&\& b < c \Longrightarrow a < c
     /* justified thanks to the ensures of `key obligations` */ }
    / ... additional properties ...
pub broadcast group group key cmp properties {
   Key::trans lt lt,
   // ... additional `broadcast` proofs from the trait
```

## **Exploring the Automation Tradeoff**

With all the Dafny Prelude lemma in Verus (with broadcast proofs), we can see what the impact is for increased implicit context:

Collection datatype	Seq	Map	Set	MultiSet	Total
$\#$ of group_ <type>_properties lemmas</type>	25	2	10	12	49

- 1. Does increasing the number of quantified facts in context result in more automation, i.e. fewer manual user-provided hints (in the form of asserts)?
- 2. Does increasing the number of quantified facts in context hinder verification performance or the verification experience?



## Projects under Study

- IronKV (SOSP24) is a distributed key-value store
- Splinter is an ongoing work on a key-value store designed around a Bε-tree.
- Anvil (OSDI24) is a framework for building and formally verifying Kubernetes controllers.
- CapybaraKV (OSDI25) is a storage system targeting persistent memory devices

#### **Minimization**

- To quantify "automation", we use the number of asserts as a metric, but most projects don't have a minimized number of assertions.
- We ran VerusMinimizer to compare the number of asserts w/ and w/o ambient facts

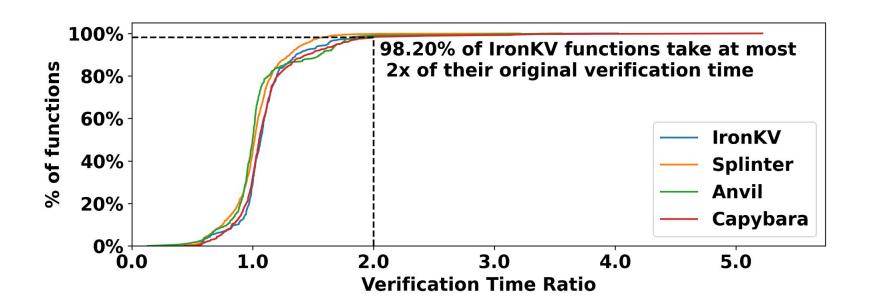
- Limitations:
  - Only a linear scan
  - Only target assertions (since detecting lemma calls can't be done syntactically)

### **Effect of Ambient Facts**

System	Original #	Minimized	Minimized with	
	of asserts	(baseline)	Ambient Facts	
IronKV	646	268	245(-23, -8.6%)	
Splinter	2678	1158	1130(-28, -2.4%)	
Anvil	701	343	336(-7, -2.0%)	
CapybaraKV	941	449	415 (-34, -7.6%)	

Table 1: Assertion counts for verification systems after minimization and importing ambient facts

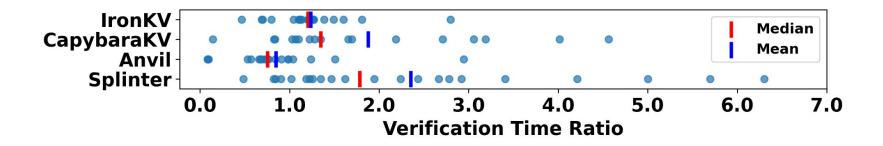
## Impact on Verification Failure Time



#### Verification Failure Time

- Once the solver finds a proof, it stops, which does not happen if the proof is incomplete or incorrect.
- As we dramatically increased search space, we might also have dramatic increase in verification failure time.

## Impact on Verification Failure Time



default safe triggers, or more aggressive multiple trigger groups?

```
forall|x: int, y: int| f(x) & g(y) & h(x,y)
           default: (0)
                                    ^^^^
     all candidates: (1)
                       ^^^
                             ^^^
                   (2)
                       ^^^
                                  ^^^^^ X <= 5
                             ^^^^ X <= 5
                   (3)
                   (4)
                       ^^^^ X <= 5
                   (5)
                                    ^^^^
                       ^^^
       all_triggers: (1)
                              ^^^
                   (5)
                                    ^^^^
```

## More Automation with More Triggers (Axioms?)

Is it possible to obtain a more extensive automation (i.e. a smaller number of required hints (asserts) to the solver) by changing how triggers are selected for the default set of quantified facts imported when using the Verus standard library?

#### You can't really "free" the triggers without explosion

```
pub broadcast axiom fn axiom_seq_add_len<A>(s1: Seq<A>, s2: Seq<A>)
    ensures
    #[trigger] s1.add(s2).len() == s1.len() + s2.len(),
;
```

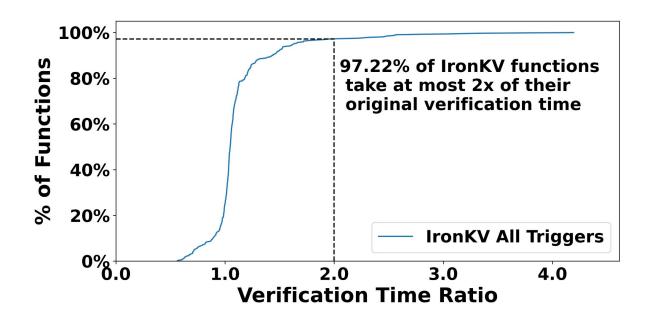
## More Automation with More Triggers (User Code?)

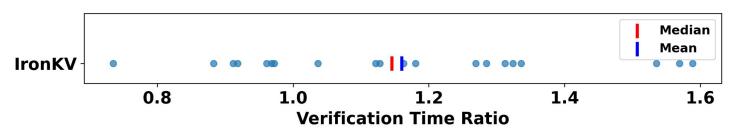
We added 206 #![all\_triggers] annotations in the 261 quantifiers in IronKV.

One proof broke.

System	Original	Min	All Triggers	Ambient Facts
IronKV	646	268	235(-33, -12.3%)	245(-23, -8.6%)

Table 2: Assertion counts for IronKV after minimization with all\_triggers





## Lessons from all\_triggers?

- all\_triggers has a non-trivial impact, though it is more susceptible to verification failure
- Managing triggers are trickier, it's hard to manage in bulk
- Experiment with multiple trigger groups for different level of automation?

#### Conclusion

## Thank you! ayb5065@nyu.edu

- Broadcast as a way of managing quantified facts
  - Customized axioms for user built abstractions
  - unsat-core for pinpointing useful lemmas
- VerusMinimizer to check "minimized" assertions
- More ambient facts in context:
  - reduced 2-9% assertions
  - only have impact on a handful of functions
- Turning on `all\_triggers`
  - Preliminary result showed it's more impactful than ambient facts
  - harder to automatically use all\_triggers