Page 2/2

```
swtch.txt
Sep 04, 2023 10:32
                                                                             Page 1/2
   CS 202, Spring 2023
   Handout 10 (Class 17)
2
   1. User-level threads and swtch()
       We'll study this in the context of user-level threads.
       Per-thread state in thread control block:
            typedef struct tcb {
10
11
                unsigned long saved_rsp;
                                             /* Stack pointer of thread */
                                             /* Bottom of thread's stack */
12
                char *t_stack;
                /* ... */
13
                };
14
15
16
       Machine-dependent thread initialization function:
17
            void thread_init(tcb **t, void (*fn) (void *), void *arg);
18
19
20
       Machine-dependent thread-switch function:
21
22
            void swtch(tcb *current, tcb *next);
23
        Implementation of swtch(current, next):
24
25
26
            # gcc x86-64 calling convention:
            # on entering swtch():
27
            # register %rdi holds first argument to the function ("current")
28
            # register %rsi holds second argument to the function ("next")
29
30
31
            # Save call-preserved (aka "callee-saved") regs of 'current'
            pushq %rbp
32
33
            pushq %rbx
            pushq %r12
34
            pushq %r13
35
36
            pushq %r14
            pushq %r15
37
38
            # store old stack pointer, for when we swtch() back to "current" later
39
            movq %rsp, (%rdi)
                                                      # %rdi->saved_rsp = %rsp
41
            movq (%rsi), %rsp
                                                     # %rsp = %rsi->saved_rsp
42
            # Restore call-preserved (aka "callee-saved") regs of 'next'
43
44
            popq %r15
45
            popq %r14
46
            popq %r13
47
            popq %r12
48
            popq %rbx
49
            popq %rbp
50
            # Resume execution, from where "next" was when it last entered swtch()
52
53
54
```

```
56 2. Example use of swtch(): the yield() call.
57
       A thread is going about its business and decides that it's executed for
58
       long enough. So it calls yield(). Conceptually, the overall system needs
59
60
       to now choose another thread, and run it:
62
       void yield() {
63
           tcb* next = pick_next_thread(); /* get a runnable thread */
64
65
           tcb* current = get_current_thread();
66
67
            swtch(current, next);
68
            /* when 'current' is later rescheduled, it starts from here */
69
70
71
72
   3. How do context switches interact with I/O calls?
73
       This assumes a user-level threading package.
74
75
       The thread calls something like "fake_blocking_read()". This looks
76
77
       to the _thread_ as though the call blocks, but in reality, the call
78
       is not blocking:
79
       int fake_blocking_read(int fd, char* buf, int num) {
80
81
           int nread = -1;
82
83
           while (nread == -1) {
84
85
                /* this is a non-blocking read() syscall */
86
87
                nread = read(fd, buf, num);
88
                if (nread == -1 && errno == EAGAIN) {
89
90
                     * read would block. so let another thread run
91
                     * and try again later (next time through the
92
                     * loop).
93
                    yield();
95
96
97
           }
99
           return nread;
100
101
102
103
104
```

swtch.txt

Sep 04, 2023 10:32