

CSCI-UA.0201

Computer Systems Organization

Data Representation – Bits and Bytes

Thomas Wies

wies@cs.nyu.edu

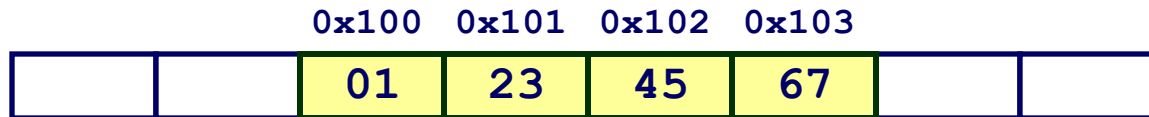
<https://cs.nyu.edu/wies>

Byte Ordering Example

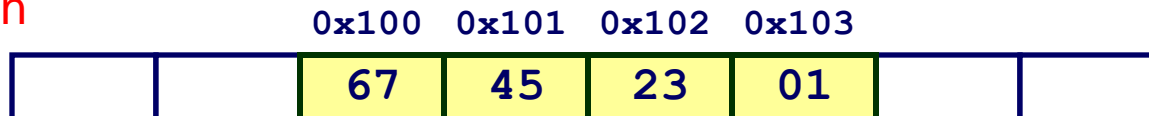
- Big Endian
 - Most significant byte has lowest address
- Little Endian
 - Most significant byte has highest address
- Example
 - Variable x has 4-byte representation `0x01234567`
 - Address given by `&x` is `0x100`

Most
Significant
Byte

Big Endian



Little Endian



Examining Data Representations

- Code to print Byte Representation of data

```
void show_bytes(unsigned char * start, int len){
    int i;
    for (i = 0; i < len; i++)
        printf("%p\t%2x\n", start+i, start[i]);
    printf("\n");
}
```

printf directives:

%p: Print pointer

%x: Print integer in hexadecimal

show_bytes Execution Example

```
int a = 0x12345678;  
printf("int a = 0x12345678;\n");  
show_bytes((unsigned char *) &a, sizeof(int));
```

Result (Linux):

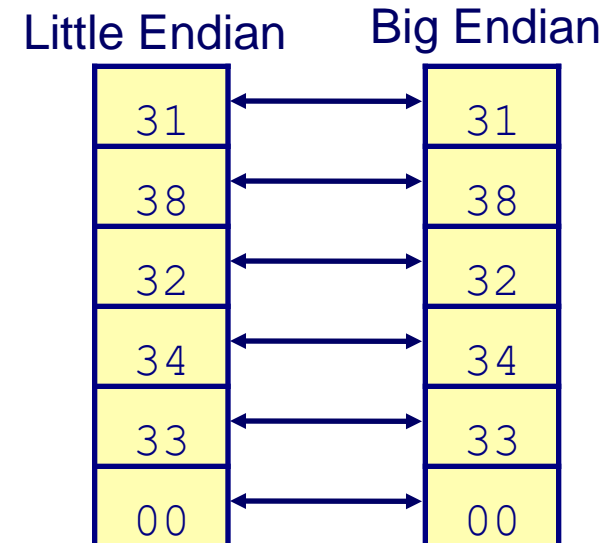
```
int a = 0x12345678;  
0x11ffffcb8  0x78  
0x11ffffcb9  0x56  
0x11ffffcba  0x34  
0x11ffffcbb  0x12
```

Representing Strings

```
char S[6] = "18243";
```

- Strings in C
 - Represented by array of characters
 - Each character encoded in ASCII format
 - Standard 7-bit encoding of character set
 - Character '0' has code 0x30
 - Digit i has code $0x30+i$
 - String should be null-terminated
- **Byte ordering not an issue**

Byte ordering is an issue for a single data item.
An array is a group of data items.



How to Manipulate Bits?

Boolean Algebra

- Applying Boolean operations, such as XOR, NAND, AND, ..., to bits to generate new bit values.

And

- $A \& B = 1$ when both $A=1$ and $B=1$

A	B	A & B
0	0	0
0	1	0
1	0	0
1	1	1

Or

- $A | B = 1$ when either $A=1$ or $B=1$

A	B	A B
0	0	0
0	1	1
1	0	1
1	1	1

Not

- $\sim A = 1$ when $A=0$

A	$\sim A$
0	1
1	0

Exclusive-Or (Xor)

- $A \wedge B = 1$ when either $A=1$ or $B=1$, but not both

A	B	$A \wedge B$
0	0	0
0	1	1
1	0	1
1	1	0

Boolean Algebra

- Applying Boolean operations, such as XOR, NAND, AND, ..., to bits to generate new bit values.

NAND

■ The reverse of AND

A	B	$\sim(A \& B)$
0	0	1
0	1	1
1	0	1
1	1	0

NOR

■ The reverse of OR

A	B	$\sim(A B)$
0	0	1
0	1	0
1	0	0
1	1	0

Exclusive-NOR (Xor)

■ The reverse of XOR

A	B	$\sim(A \wedge B)$
0	0	1
0	1	0
1	0	0
1	1	1

Application of Boolean Algebra

- Applied to **Digital Systems** by Claude Shannon
 - 1937 MIT Master's Thesis
 - Reason about networks of relay switches
 - Encode closed switch as 1, open **switch** as 0



Transistor

Lifting Operations to Bit Vectors

- Operate on Bit Vectors (e.g. an integer is a bit vector of 4 bytes = 32 bits)

– Operations applied bitwise

01101001	01101001	01101001	
<u>& 01010101</u>	<u> 01010101</u>	<u>^ 01010101</u>	<u>~ 01010101</u>
01000001	01111101	00111100	10101010

Bit-Level Operations in C

- Operations $\&$, $|$, \sim , \wedge Available in C
 - Apply to any “integral” data type
 - `long`, `int`, `short`, `char`, `unsigned`
- Examples (Char data type)
 - $\sim 0x41 = 0xBE$
 - $\sim 01000001_2 = 10111110_2$
 - $\sim 0x00 = 0xFF$
 - $\sim 00000000_2 = 11111111_2$
 - $0x69 \& 0x55 = 0x41$
 - $01101001_2 \& 01010101_2 = 01000001_2$
 - $0x69 | 0x55 = 0x7D$
 - $01101001_2 | 01010101_2 = 01111101_2$

Contrast: Logic Operations in C

- Contrast to Logical Operators

`&&`, `||`, `!`

- View 0 as "false"
- Anything nonzero as "true"
- Always return 0 or 1

- Examples

– `!0x41` = `0x00`

– `!0x00` = `0x01`

– `!!0x41` = `0x01`

– `0x69 && 0x55` = `0x01`

– `0x69 || 0x55` = `0x01`

– `p && *p` (avoids null pointer access – short circuiting)

Type `bool` in C

- Did not exist in standard C89/90
- It was introduced in C99 standard
- You may need to use the following switch with gcc:

```
gcc -std=c99 ...
```

```
#include <stdbool.h>
```

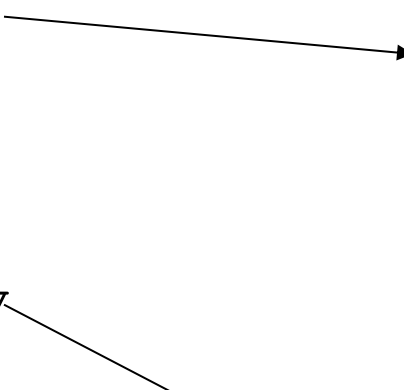
```
bool x;
```

```
x = false; ← lower case
```

```
x = true;
```

Shift Operations

- Left Shift: $x \ll y$
 - Shift x left by y positions
 - Throw away extra bits on left
 - Fill with 0's on right
- Right Shift: $x \gg y$
 - Shift x right y positions
 - Throw away extra bits on right
 - type1: **Logical shift**
 - Fill with 0's on left
 - type 2: **Arithmetic shift** (covered later)
 - Replicate most significant bit on right
- Undefined Behavior
 - Shift amount < 0 or \geq size of x



Argument x	01100010
$\ll 3$	00010 000
Log. $\gg 2$	00 011000
Arith. $\gg 2$	00 011000

Argument x	10100010
$\ll 3$	00010 000
Log. $\gg 2$	00 101000
Arith. $\gg 2$	11 101000

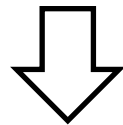
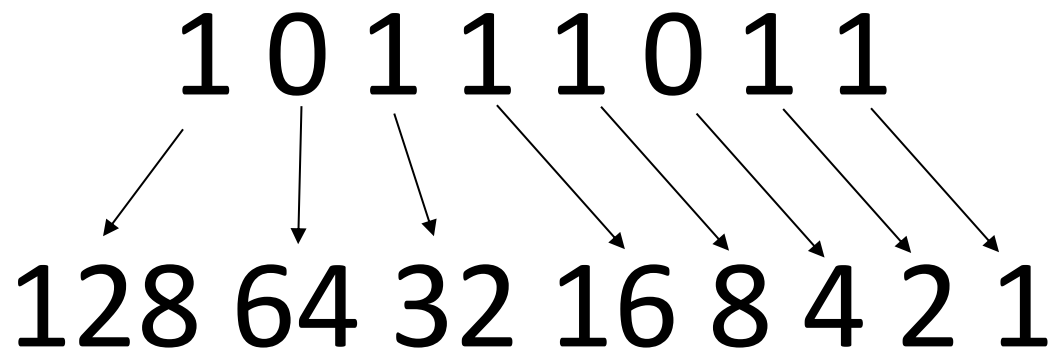
How to present Integers
in binary?

Two Types of Integers

- Unsigned
 - positive numbers and 0
- Signed numbers
 - negative numbers as well as positive numbers and 0

Unsigned Integers

$$B2U(X) = \sum_{i=0}^{w-1} x_i \cdot 2^i$$



187

Unsigned Integers

- An n -bit unsigned integer represents 2^n values:
from 0 to $2^n - 1$.

2^2	2^1	2^0	
0	0	0	0
0	0	1	1
0	1	0	2
0	1	1	3
1	0	0	4
1	0	1	5
1	1	0	6
1	1	1	7

Unsigned Binary Arithmetic

- Base-2 addition – just like base-10
 - add from right to left, propagating carry

$\begin{array}{r} 10010 \\ + \underline{1001} \\ \hline 11011 \end{array}$	$\begin{array}{r} \text{carry} \\ \curvearrowright \\ 10010 \\ + \underline{1011} \\ \hline 11101 \end{array}$	$\begin{array}{r} \curvearrowleft \curvearrowleft \curvearrowleft \curvearrowleft \\ 1111 \\ + \underline{\quad 1} \\ \hline 10000 \end{array}$
	$\begin{array}{r} 10111 \\ + \underline{111} \\ \hline \end{array}$	

What About Negative Numbers?

People have tried several options:

Sign Magnitude:

$$000 = +0$$

$$001 = +1$$

$$010 = +2$$

$$011 = +3$$

$$100 = -0$$

$$101 = -1$$

$$110 = -2$$

$$111 = -3$$

One's Complement

$$000 = +0$$

$$001 = +1$$

$$010 = +2$$

$$011 = +3$$

$$100 = -3$$

$$101 = -2$$

$$110 = -1$$

$$111 = -0$$

Two's Complement

$$000 = +0$$

$$001 = +1$$

$$010 = +2$$

$$011 = +3$$

$$100 = -4$$

$$101 = -3$$

$$110 = -2$$

$$111 = -1$$

- Issues: balance, number of zeros, ease of operations
- Which one is best? Why?

Signed Integers

- With **n bits**, we have **2^n distinct values**.
 - assign about half to positive integers and about half to negative
- **Positive integers**
 - just like unsigned: zero in *most significant* (MS) bit
00101 = 5
- **Negative integers**
 - In two's complement form

In general: **a 0 at the MS bit indicates positive**
and a 1 indicates negative.

Two's Complement

- *Two's complement* representation developed to make circuits easy for arithmetic.
 - for each positive number (X), assign value to its negative ($-X$), such that $X + (-X) = 0$ with “normal” addition, ignoring carry out

$$\begin{array}{r} 00101 \quad (5) \\ + \underline{11011} \quad (-5) \\ \hline 00000 \quad (0) \end{array}$$

$$\begin{array}{r} 01001 \quad (9) \\ + \underline{10111} \quad (-9) \\ \hline 00000 \quad (0) \end{array}$$

Two's Complement Signed Integers

- MS bit is sign bit.
- Range of an **n-bit number**: -2^{n-1} through $2^{n-1} - 1$.
 - The most negative number (-2^{n-1}) has no positive counterpart.

-2^3	2^2	2^1	2^0		-2^3	2^2	2^1	2^0	
0	0	0	0	0	1	0	0	0	-8
0	0	0	1	1	1	0	0	1	-7
0	0	1	0	2	1	0	1	0	-6
0	0	1	1	3	1	0	1	1	-5
0	1	0	0	4	1	1	0	0	-4
0	1	0	1	5	1	1	0	1	-3
0	1	1	0	6	1	1	1	0	-2
0	1	1	1	7	1	1	1	1	-1

Converting Binary (2's C) to Decimal

1. If MS bit is one (i.e. number is negative), take two's complement to get a positive number.
2. Get the decimal as if the number is unsigned (using power of 2s).
3. If original number was negative, add a minus sign.

n	2^n
0	1
1	2
2	4
3	8
4	16
5	32
6	64
7	128
8	256
9	512
10	1024

Examples

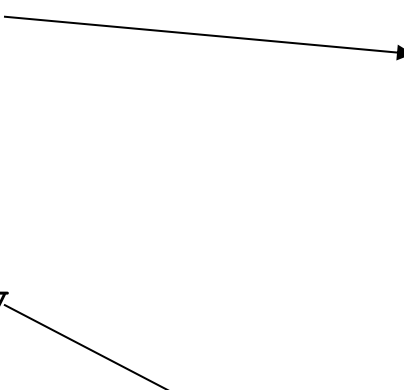
$$\begin{aligned} X &= 00100111_{\text{two}} \\ &= 2^5 + 2^2 + 2^1 + 2^0 = 32 + 4 + 2 + 1 \\ &= 39_{\text{ten}} \end{aligned}$$

$$\begin{aligned} X &= 11100110_{\text{two}} \\ -X &= 00011010 \\ &= 2^4 + 2^3 + 2^1 = 16 + 8 + 2 \\ &= 26_{\text{ten}} \\ X &= -26_{\text{ten}} \end{aligned}$$

n	2^n
0	1
1	2
2	4
3	8
4	16
5	32
6	64
7	128
8	256
9	512
10	1024

Shift Operations

- Left Shift: $x \ll y$
 - Shift x left by y positions
 - Throw away extra bits on left
 - Fill with 0's on right
- Right Shift: $x \gg y$
 - Shift x right y positions
 - Throw away extra bits on right
 - type1: **Logical shift**
 - Fill with 0's on left
 - type 2: **Arithmetic shift** (covered later)
 - Replicate most significant bit on right
- Undefined Behavior
 - Shift amount < 0 or \geq size of x



Argument x	01100010
$\ll 3$	00010000
Log. $\gg 2$	00011000
Arith. $\gg 2$	00011000

Argument x	10100010
$\ll 3$	00010000
Log. $\gg 2$	00101000
Arith. $\gg 2$	11101000

Numeric Ranges

Example: Assume 16-bit numbers

	Decimal	Hex	Binary
Unsigned Max	65535	FF FF	11111111 11111111
Signed Max	32767	7F FF	01111111 11111111
Signed Min	-32768	80 00	10000000 00000000
-1	-1	FF FF	11111111 11111111
0	0	00 00	00000000 00000000

Values for Different Sizes

	W			
	8	16	32	64
Unsig. Max	255	65,535	4,294,967,295	18,446,744,073,709,551,615
Signed Max	127	32,767	2,147,483,647	9,223,372,036,854,775,807
Signed Min	-128	-32,768	-2,147,483,648	-9,223,372,036,854,775,808

■ C Programming

- `#include <limits.h>`
- Declares constants, e.g.,
 - `INT_MAX`
 - `LONG_MAX`
 - `INT_MIN`
 - `UINT_MIN`
 - ...