Fundamental Algorithms, Assignment 11 Due April 27/8 in Recitation.

My work has always tried to unite the true with the beautiful and when I had to choose one or the other I usually chose the beautiful. – Hermann Weyl

- 1. Consider Dumb Prim for MST. The high level idea is the same but to find the minimal weight of an edge $\{v, w\}, v \in S, w \notin S$, one looks at all the weights and finds the minimum in the usual way. (There is no updating in Dumb Prim.) Assume that all pairs $\{v, w\}$ have a weight. Let n be the number of vertices.
 - (a) When |S| = i what is the time to add a vertex to S as a function of n and i.
 - (b) What is the total time for Dumb Prim as a sum over *i*.
 - (c) Evaluate the above sum as $\Theta(g(n))$ for some nice function g(n). (Caution: The time is *not* an increasing function of *i*. For example, when i = n - 1 the time is quite quick.)
 - (d) Compare the time for Dumb Prim with Prim as discussed in class
- 2. Consider Prim's Algorithm for MST on the complete graph with vertex set $\{1, \ldots, n\}$. Assume that edge $\{i, j\}$ has weight $(j - i)^2$. Let the root vertex r = 1. Show the pattern as Prim's Algorithm is applied. In particular, Let n = 100 and consider the situation when the tree created has 73 elements and π and key have been updated.
 - (a) What are these 73 elements.
 - (b) What are $\pi[84]$ and key[84].
- 3. Find $d = \gcd(89, 55)$ and x, y with 89x + 55y = 1. [Remark: This is part of a pattern with two consecutive numbers from the Fibonacci sequence $0, 1, 1, 2, 3, 5, 8, 13, 21, 34, 55, 89, \ldots$]
- 4. Find $\frac{211}{507}$ in Z_{1000} .
- 5. Solve the system $x \equiv 34 \mod 101$ $x \equiv 59 \mod 103.$
- 6. Using the Island-Hopping Method to find 2^{1000} modulo 1001 using a Calculator but NOT using multiple precision arithmetic. (You should never have an intermediate value more than a million.)

7. (extra from last week!) Suppose that during Kruskal's Algorithm (for MST) and some point we have SIZE[v] = 37. What is the interpretation of that in the case when $\pi[v] = v$? What is the interpretation of that in the case when $\pi[v] = u \neq v$? Let w be a vertex. How many different values can $\pi[w]$ have during the course of Kruskal's algorithm? How many different values (as a function of V, the number of vertices) can SIZE[w] have during the course of Kruskal's algorithm? (That is, the maximal number possible.)

The universe is not only queerer than we suppose but queerer than we *can* suppose.

– J.B.S. Haldane