

# Basic Algorithms, Assignment 11

## Solutions

1. Here is a variant implementation **EasyPrim** that doesn't use the parent  $\pi$ . With a given set  $S$  the data structure will be a minheap  $Q$  consisting of *all* crossing edges  $\{x, y\}$  with  $key[x, y] = w(x, y)$ . The initial step, the minheap of all edges  $\{s, x\}$  is the same. Now  $V - 1$  times we EXTRACT-MAX getting a crossing edge  $\{x, y\}$ ,  $x \in S, y \notin S$  of minimal weight. We add edge  $\{x, y\}$  to  $T$  and  $y$  to  $S$ . The UPDATE is different. For each  $z \in Adj[y]$  we examine  $\{y, z\}$ . If  $z \in S$  then  $\{y, z\}$  is no longer a crossing edge so we delete it (caution: this takes time!) from the minheap. If  $z \notin S$  then  $\{y, z\}$  has become a crossing edge so we add it to the minheap, using its weight as its key.

Analyze the time for **EasyPrim**. Show that the total time is  $O(E \lg E)$ .

**Solution:** The initial step takes  $O(V)$  as before. The minheap has at most size  $E$ . The EXTRACT-MAX takes  $O(\lg E)$  each time, a total of  $O(V \lg E)$ . Adding or deleting from the minheap takes time  $O(\lg E)$ . For every edge it can at most be added to the minheap and later deleted from the minheap so the cost for the edge is  $O(\lg E)$  which in total gives  $O(E \lg E)$ . This is the dominant time. [Note that  $O(E \lg E)$  and  $O(E \lg V)$  are the same as  $E \leq V^2$  so that in  $\Theta$ -land **Prim** and **EasyPrim** have the same time – but in reality, with the constants, **Prim** is faster.]

2. Consider Prim's Algorithm for MST on the complete graph with vertex set  $\{1, \dots, n\}$ . Assume that edge  $\{i, j\}$  has weight  $|j - i|^3$ . Let the root vertex  $r = 1$ . Show the pattern as Prim's Algorithm is applied.

**Solution:** The set  $S$ , initially  $\{1\}$ , will grow to  $\{1, 2\}, \dots, \{1, 2, \dots, i\}, \dots, \{1, \dots, n\}$ . When  $S = \{1, \dots, i\}$  the closest point to  $S$  will be  $i + 1$  with  $\pi[i + 1] = i$  and  $key[i + 1] = 1$ . In particular, Let  $n = 500$  and consider the situation when the tree created has 211 vertices and  $\pi$  and  $key$  have been updated.

- (a) What are these 211 vertices and what are the edges.

**Solution:**  $1, \dots, 211$ . The edges are  $\{1, 2\}, \{2, 3\}, \dots, \{210, 211\}$ .

- (b) What are  $\pi[309]$  and  $key[309]$ .

**Solution:**  $\pi[309] = 211$  (all other of  $1, \dots, 202$  are further) and  $key[309] = (309 - 211)^3$ .

3. Find  $d = \gcd(144, 89)$  and  $x, y$  with  $144x + 89y = 1$ . [Remark: This is part of a pattern with two consecutive numbers from the Fibonacci

sequence 0, 1, 1, 2, 3, 5, 8, 13, 21, ...]

**Solution:**

$$\begin{aligned}
 \text{EUCLID}(144, 89) &= \\
 \text{EUCLID}(89, 55) &= \text{EUCLID}(55, 34) = \text{EUCLID}(34, 21) = \\
 &= \text{EUCLID}(21, 13) = \text{EUCLID}(13, 8) = \text{EUCLID}(8, 5) = \\
 &= \text{EUCLID}(5, 3) = \text{EUCLID}(3, 2) = \text{EUCLID}(2, 1) = \\
 &= \text{EUCLID}(1, 0) = 1
 \end{aligned}$$

with all quotients 1 except the last. For EXTENDED – EUCLID we get a chart like Figure 31.1:

$a$	$b$	$\lfloor a/b \rfloor$	$d$	$x$	$y$
144	89	1	1	34	-55
89	55	1	1	-21	34
55	34	1	1	13	-21
34	21	1	1	-8	13
21	13	1	1	5	-8
13	8	1	1	-3	5
8	5	1	1	2	-3
5	3	1	1	-1	2
3	2	1	1	1	-1
2	1	2	1	0	1
1	0	-	1	1	0

so  $x = 34$  and  $y = -55$ . (Note that the  $x$ 's and  $y$ 's form a Fibonacci like pattern as well!)

4. Find  $\frac{311}{507}$  in  $Z_{1000}$ .

**Solution:** Here we first find  $\text{EUCLID}(1000, 507)$ :

$$\begin{aligned}
 \text{EUCLID}(1000, 507) &= \text{EUCLID}(507, 493) = \text{EUCLID}(493, 14) = \\
 &= \text{EUCLID}(14, 3) = \text{EUCLID}(3, 2) = \text{EUCLID}(2, 1) = \\
 &= \text{EUCLID}(1, 0) = 1
 \end{aligned}$$

For EXTENDED – EUCLID we get a chart like Figure 31.1:

$a$	$b$	$\lfloor a/b \rfloor$	$d$	$x$	$y$
1000	507	1	1	181	-357
507	493	1	1	-176	181
493	14	1	35	5	-176
14	3	1	4	-1	5
3	2	1	1	1	-1
2	1	1	2	0	1
1	0	-	1	1	0

so that  $1000(181) - 357(507) = 1$  so in  $Z_{1000}$  we have  $(-357)(507) = 1$  so  $\frac{1}{507} = -357 = 643$ . Finally  $\frac{311}{507} = 311 \cdot 643 = 199973 = 973$ . So the answer is 973. (You might prefer to write it as  $-27$  which is the same.) To check:  $973 \cdot 507 = 493311 = 311$ .

5. Solve the system

$$x \equiv 34 \pmod{101}$$

$$x \equiv 59 \pmod{103}.$$

**Solution:** We write  $x = 103y + 59$  (we could start with either and this one is a bit easier) so that in  $Z_{101}$  we want  $103y + 59 = 34$  or  $2y = -25 = 76$  and  $y = 38$ . (Usually division is complicated but here it worked out like normal division.) Then  $x = 103(38) + 59 = 3973$ . The general answer is given as  $x \equiv 3973 \pmod{10403}$  as  $10403 = 103 \cdot 101$ .