

LECTURE 11:

ADAPTIVE ENCODING MODELS, PPM

October 19, 2005

- We would like to get both the advantages of:
 - fast learning of a low-order model
 - ultimately better prediction of a high-order model
- We can do this by *varying* the order we use.
- PPM maintains frequencies for characters that have been seen before in *all* contexts that have occurred before, up to some maximum order.
Example: Suppose we have so far encoded the string

```

this_is_th

```
- If we are using contexts up to order two, then we will record frequencies for the following contexts:
 - Order 0: ()
 - Order 1: (t) (h) (i) (s) (_)
 - Order 2: (th) (hi) (is) (s_) (_i) (_t)

- So far, we've looked at models in which the symbols would be *independent*, if we knew what their probabilities were.
- If we don't know the probabilities, our predictions do depend on previous symbols, but the symbols are still "exchangeable" — their order doesn't matter.
- Very often, this isn't right: The probability of a symbol may depend on the *context* in which it occurs — eg, what symbol precedes it.
- **Example:** "U" is much more likely after "Q" (in English), than after another "U". Probabilities may also depend on position in the file, though modeling this is less common.
- We looked into using Markov models of various orders but there was a problem: with small amounts of data we can't estimate the required frequencies very well.

- The frequency tables maintained by PPM contain *only* the characters that have been seen before in that context.
Examples: if "x" has never occurred, none of the frequency tables will have an entry for "x".
If "x" *has* occurred before, but *not* after a "t", the frequency table for order 1 context (t) will not contain "x".
- **The main idea:** If we need to encode a character that doesn't appear in the context we're using, we transmit an "escape" flag, and switch to a lower-order context.
- What if we escape from every context? We end up in a special "order -1" context, in which every character has a frequency of 1.

- Two details about frequencies need to be resolved.
- First, what characters do we count in a context?
 - We might count *every* character that appears following the characters making up the context.
 - We might count a character in a context *only* when it does not appear in a higher-order context.
- One could argue for either way, but we'll go for the second option.
- Second, what do we use as the frequency of the “escape” symbol? There are many possibilities. We'll just always give it a frequency of one, no matter how many times we escape a given context.

Order -1: `_:1 a:1 b:1 ... z:1`
 Order 0: `() Escape:1 t:2 h:1 i:2 s:1 _:1`
 Order 1:
 `(t) Escape:1 h:2`
 `(h) Escape:1 i:1`
 `(i) Escape:1 s:2`
 `(s) Escape:1 _:1`
 `(_) Escape:1 i:1 t:1`
 Order 2:
 `(th) Escape:1 i:1`
 `(hi) Escape:1 s:1`
 `(is) Escape:1 _:2`
 `(s_) Escape:1 i:1 t:1`
 `(_i) Escape:1 s:1`
 `(_t) Escape:1 h:1`

Loop until end of file:

 Read the next character, c .

 Let d_K, d_{K-1}, \dots, d_1 be the preceding K characters.

 Set the context size, k , to the maximum, K .

 While (d_k, \dots, d_1) hasn't been seen previously:

 Set k to $k - 1$.

 While $k \geq 0$ and c hasn't been seen in context (d_k, \dots, d_1) :

 Transmit an escape flag using context (d_k, \dots, d_1) .

 Set k to $k - 1$.

 If $k = -1$: {Transmit c using the special “order -1” context. Set k to 0. }

 Else {Transmit c using context (d_k, \dots, d_1) .}

 While $k \leq K$:

 Create context (d_k, \dots, d_1) if it doesn't exist.

 Increment the count for c in context (d_k, \dots, d_1) .

 Set k to $k + 1$.

- One reason PPM works well for files like English text is that it can implicitly learn the vocabulary — the dictionary of words in the language. This is because early letters of a word like “Ontario” almost completely determine the remaining letters.
- A more direct approach is to store a dictionary explicitly. When a word is encountered, a short code for it is sent, rather than the letters.
- The “LZ” (for Lempel-Ziv) family of data compression algorithms build a dictionary adaptively, based on the text seen previously. The “gzip” program is an example.
- Dictionary Methods, like `gzip` or `compress` are not quite as good as the best adaptive models but they can be *much* faster during encoding and decoding because they don't have to do operations for every single symbol.

- A version of PPM (written by Bill Teahan) and gzip applied to the three English text files from before:

PPM			
Uncompressed file size	Compressed file size	Compression factor	Bits per character
2344	1042	2.25	3.56
20192	5903	3.42	2.34
235215	51323	4.58	1.75

GZIP			
Uncompressed file size	Compressed file size	Compression factor	Bits per character
2344	1160	2.02	3.96
20192	7019	2.88	2.78
235215	70030	3.36	2.38

- Speed: On the long file, PPM took 2.2 to encode, 2.3s to decode; gzip needed only 60ms to encode, <1ms to decode.

- Compression using adaptive dictionaries may be less elegant, but has it's own advantages:
 - Dictionary methods can be quite fast (especially at decoding), since whole sequences of symbols are specified at once.
 - The idea that the data contain many repeated strings fits many sources quite well — eg, English text, machine-language programs, files of names and addresses.
- The main disadvantage is that compression may not be as good as a model based method:
 - Dictionaries are inappropriate for some sources — eg, noisy images.
 - Even when dictionaries work well, a good model-based method may do better — and can't do worse, if it uses the same modeling ideas as the dictionary method.

- N -th order Markov models and PPM models cleanly separate the *model* for symbol probabilities from the *coding* based on those probabilities.
- Such models have several advantages:
 - Coding can be nearly optimal (eg, using arithmetic coding).
 - It's easy to try out various modeling ideas.
 - You can get very good compression, if you use a good model.
- The big disadvantage:
 - The coding and decoding involves operations for every symbol and every bit, plus possibly expensive model updates, which limits how fast these methods can be.

- This scheme was devised by Ziv and Lempel in 1977. There are many variants, including the method used by gzip.
- The idea of LZ77 is to use the past text as the dictionary — avoiding the need to transmit a dictionary separately. We need a buffer of size W that contains the previous S characters plus the following $W - S$ characters.
- We encode up to $W - S$ characters at once by sending the following:
 - A pointer to a past character in the buffer (an integer from 1 to S).
 - The number of characters to take from the buffer (an integer from 0 to $W - S - 1$, or maybe more).
 - The single character that follows the string taken from the buffer.

- Suppose we look at the past 16 characters, and look ahead at the next 8 characters.
- After encoding the first 16 characters of the following string, we would proceed as follows:

```
W a y _ o v e r _ t h e r e _ i s _ w h e r e _ i t _ i s
```

No match with string in window.

Transmit (-,0,s)

```
W a y _ o v e r _ t h e r e _ i s _ w h e r e _ i t _ i s
```

Match 3 back with _

Transmit (3,1,w)

```
W a y _ o v e r _ t h e r e _ i s _ w h e r e _ i t _ i s
```

Match with 9 back with here_i

Transmit (9,6,t)