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Robert Hooke





- Robert Hooke (1635-1703) was an experimental scientist, mathematician, architect, and astronomer. Secretary of the Royal Society from 1677 to 1682, ...
- Hooke was considered the "England's Da Vinci" because of his wide range of interests.
- His work Micrographia of 1665 contained his microscopical investigations, which included the first identification of biological cells.
- In his drafts of Book II, Newton had referred to him as the most illustrious Hooke—"Cl[arissimus] Hookius."
- Hooke became involved in a dispute with Isaac Newton over the priority of the discovery of the inverse square law of gravitation.



Hooke to Halley





 "[Huygen's Preface] is concerning those properties of gravity which I myself first discovered and showed to this Society and years since, which of late Mr. Newton has done me the favour to print and publish as his own inventions."











- "Now is this not very fine? Mathematicians that find out, settle & do all the business must content themselves with being nothing but dry calculators & drudges & another that does nothing but pretend & grasp at all things must carry away all the inventions...
- "I beleive you would think him a man of a strange unsociable temper."



Newton to Hooke













Image & Logic





The great distance between

- a glimpsed truth and
- a demonstrated truth
 - Christopher Wren/Alexis Claude Clairaut



Micrographia Principia





Micrographia



MICROGRAPHIA:



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Noses. 23. 1644.

BROUNCKER, P. R.S.



MICROGRAPHIA:

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"The Brain & the Fancy"



The truth is, the science of Nature has already been too long made only a work of the brain and the fancy. It is now high time that it should return to the plainness and soundness of observations on material and obvious things."

- Robert Hooke. (1635 - 1703), *Micrographia* 1665



Juffa Societatis Reg

Principia





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LAW L

That every body perfeveres in its flate of refling, or of moving uniformly in a right line, as far as it is not compelled to change that flate by external forces impreffed upon it.

LAW II.

That the change of motion is proportional to the moving force imprefied; and is produced in the direction of the right line, in which that force is imprefied.

LAW III.

That readion is always contrary and equal to action: or, that the mutual actions of two bodies upon each other are always equal, and directed to contrary parts.

LEGES MOTU

The set of the set of





"Induction & Hypothesis"



Hypotheses non fingo. I feign no hypotheses. Principia Mathematica.

- "Truth being uniform and always the same, it is admirable to observe how easily we are enabled to make out very abstruse and difficult matters, when once true and genuine Principles are obtained."
 - Halley, "The true Theory of the Tides, extracted from that admired Treatise of Mr. Issac Newton, Intituled, Philosophiae Naturalis Principia Mathematica," *Phil. Trans.* 226:445,447.
- This rule we must follow, that the argument of induction may not be evaded by hypotheses.



Morphogenesis



Alan Turing: 1952



 "The Chemical Basis of Morphogenesis," 1952, *Phil. Trans. Roy. Soc. of London*, Series B: Biological Sciences, **237**:37—72.

• A reaction-diffusion model for development.



"A mathematical model for the growing embryo."



- A very general program for modeling embryogenesis: The `model' is ``a simplification and an idealization and consequently a falsification."
- Morphogen: "is simply the kind of substance concerned in this theory..." in fact, anything that diffuses into the tissue and "somehow persuades it to develop along different lines from those which would have been followed in its absence" qualifies.



a: concentration *Da*: diffusion constant



Turing, A.M. (1952)."The chemical basis of morphogenesis." *Phil. Trans. Roy. Soc.* London B **237**: 37





Pearson, J. E.: Complex patterns in simple systems. *Science* **261**, 189–192 (1993). Made by A-PDF PPT2PDF



Reaction-diffusion: an example





Genes: 1952





 Since the role of genes is presumably catalytic, influencing only the rate of reactions, unless one is interested in comparison of organisms, they "may be eliminated from the discussion..."



Crick & Watson :1953

11.5

MOLECULAR STRUCTURE OF NUCLEIC ACIDS

A Strattare for Deoxyribose Nucleic Acid

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GENETICAL IMPLICATIONS OF THE STRUCTURE OF DEOXYRIBONUCLEIC ACID In J. D. WATSON and F. H. C. CRICK

Redicel Research Council Unit for the Souty of the Melecular Sources of Biological Systems, Controllat Laboratory, Committee

Tatls reportences of decorphonealesis acid (DNA) at dividing solis is undeputed. It is found in all dividing solis, heighly if not excitently in the reaches, where it is an exempted constituent of the diversesources. Many lines of evidence indivisits that is it the marrier of a part of ill rait solit the generative specificity of the diversemenanges and these the itself.



Fig. 1. This farere is proved thepresentils. The two theory organizations that they phototheormaging the theory phototheory and the theory pairs of takere toolating the charts of takere the method has pairs of taker the method has pairs the

Ttha that th postula a possibility

It has not escaped our notice that the specific pairing we have postulated immediately suggests a possible copying mechanism for the genetic material."

J.D.Watson F.H.C.Crick, Nature in agazine, 2 April 1953





Genome





• Genome:

- Hereditary information of an organism is encoded in its DNA and enclosed in a cell (unless it is a virus).
 All the information contained in the DNA of a single organism is its genome.
- DNA molecule can be thought of as a very long sequence of nucleotides or bases:

 $\Sigma = \{\mathsf{A},\mathsf{T},\mathsf{C},\mathsf{G}\}$



The Central Dogma





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The central dogma(due to Francis Crick in 1958) states that these information flows are all unidirectional: "The central dogma states that once `information^rhas passed into protein it cannot get out again. The transfer of information from nucleic acid to nucleic acid, or from nucleic acid to protein, may be possible, but transfer from protein to protein, or from protein to nucleic acid is impossible. Information means here the precise determination of sequence, either of bases in the nucleic acid or of amino acid residues in the protein."



RNA, Genes and Promoters





"The Brain & the Fancy"





"Work on the mathematics of growth as opposed to the statistical description and comparison of growth, seems to me to have developed along two equally unprofitable lines... It is futile to conjure up in the imagination a system of differential equations for the purpose of accounting for facts which are not only very complex, but largely unknown,...What we require at the present time is more measurement and less theory."

Eric Ponder, Director, CSHL (LIBA), 1936-



"Axioms of Platitudes"

-E.B. Wilson





- . Science need not be mathematical.
- . Simply because a subject is mathematical it need not therefore be scientific.
- 3. Empirical curve fitting may be without other than classificatory significance.
- Growth of an individual should not be confused with the growth of an aggregate (or average) of individuals.
- Different aspects of the individual, or of the average, may have different types of growth curves.



Genes for Segmentation



- Fertilization followed by cell division
- Pattern formation instructions for
 - Body plan (Axes: A-P, D-V)
 - Germ layers (ecto-, meso-, endoderm)
- Cell movement form gastrulation
- Cell differentiation



PI: Positional Information





- Positional value
 - Morphogen a substance
 - Threshold concentration
- Program for development
 - Generative rather than descriptive
- "French-Flag Model"









 The *bicoid* gene provides an A-P morphogen gradient









gap genes

- into broad regions by gap gene expression
- The first *zygotic* genes
- Respond to maternallyderived instructions
- Short-lived proteins, gives bell-shaped distribution from source



Transcription Factors in Cascade



- Hunchback (hb), a gap gene, responds to the dose of bicoid protein
- A concentration above threshold of bicoid activates the expression of *hb*
- The more *bicoid* transcripts, the further **back** *hb* expression goes



Transcription Factors in Cascade





- Krüppel (Kr), a gap gene, responds to the dose of hb protein
- A concentration above minimum threshold of hb activates the expression of Kr
- A concentration above maximum threshold of hb inactivates the expression of Kr

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Segmentation





- Parasegments are delimited by expression of pairrule genes in a periodic pattern
- Each is expressed in a series of 7 transverse stripes

Pattern Formation







- Edward Lewis, of the California Institute of Technology
- Christiane Nuesslein-Volhard, of Germany's Max-Planck Institute
- Eric Wieschaus, at Princeton
- Each of the three were involved in the early research to find the genes controlling development of the *Drosophila* fruit fly.




Completeness





Model Parameters





Complete Model





Complete Model





Is this your final answer?



- It is not uncommon to assume certain biological problems to have achieved a cognitive finality without rigorous justification.
- Rigorous mathematical models with automated tools for reasoning, simulation, and computation can be of enormous help to uncover
 - cognitive flaws,
 - qualitative simplification or
 - overly generalized assumptions.
- Some ideal candidates for such study would include:
 - prion hypothesis
 - cell cycle machinery
 - muscle contractility
 - processes involved in cancer (cell cycle regulation, angiogenesis, DNA repair, apoptosis, cellular senescence, tissue space modeling enzymes, etc.)
 - signal transduction pathways, and many others.



Computational Systems Biology







Why do we need a tool?



Simulate Biologists! Not Biology!!



Future Biology

int -

Functional genomic hypothesis generation and experimentation by a robot scientist

Ross D. King', Kenneth E. Whelan', Filon M. Jones', Philip G. K. Reiser', Christopher H. Bryant', Slaphen H. Haggieton', Douglas B. Kell' & Stephen G. Oliver'

¹Department of Computer Science, University of Hisley, Aberystueph ST23 3D4E EX

⁵School of Computing, The Robert Gordon University, Aberdone AB10 IER, UK ⁵Dipareneeus of Computing, Imperial Calley, London SWT 2022, UK ⁴Dipareneeus of Chemistery, UNIST, PO, Stor 08, Manchester M60 (QD, UK ⁵School of Scholgend Sciences, University of Manchester, 2200 Stopfend Bookling, Manchester M13 1972, UK



Biology of the future should only involve a biologist and his dog: the biologist to watch the biological experiments and understand the hypotheses that the data-analysis algorithms produce and the dog to bite him if he ever touches the experiments or the computers.



Simpathica is a modular system



Canonical Form:

$$\begin{cases} \dot{X}_{i} = \alpha_{i} \prod_{j=1}^{n+m} X_{j}^{g_{i}} - \beta_{i} \prod_{j=1}^{n+m} X_{j}^{h_{i}} \quad i = 1...n \\ C_{i}(X_{1}(t), \dots, X_{n+m}(t)) = \sum (\gamma_{i} \prod_{j=1}^{n+m} X_{j}^{f_{i}}) = 0 \end{cases}$$

Characteristics:

- Predefined Modular Structure
- Automated Translation from Graphical to Mathematical Model
- Scala bility

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Figure 1: Representation of an unmodified and of a reversible reaction.



Figure 2: Representation of a divergence and of a convergence branch point (the two processes in each reaction are independent of each other).



Figure 3: Representation of a single splitting reaction generating two products, X₂ and X₃, in stoichiometric proportions and of a single synthetic reaction involving two source components, X₁ and X₂ always in stoichiometric proportions.



Figure 4: The conversion of X_1 into X_2 is modulated (stimulation or inhibition is represented by the sign of the arrow) by X_3 . The reaction between X_1 and X_2 requires corraying X_3 , which in the process is converted into X_4 .





Formal Definition of S-system

Definition 1 (S-system). An S-system is a quadruple S = (DV, IV, DE, C) where:

- $-DV = \{X_1, \ldots, X_n\}$ is a finite non empty set of dependent variables ranging over the domains D_1, \ldots, D_n , respectively;
- $-IV = \{X_{n+1}, \ldots, X_{n+m}\}$ is a finite set of independent variables ranging over the domains D_{n+1}, \ldots, D_{n+m} , respectively;
- DE is a set of differential equations, one for each dependent variable, of the form

$$\dot{X}_{i} = \alpha_{i} \prod_{j=1}^{n+m} X_{j}^{g_{ij}} - \beta_{i} \prod_{j=1}^{n+m} X_{j}^{h_{ij}}$$

with $\alpha_i, \beta_i \geq 0$ called rate constants;

-C is a set of algebraic constraints of the form

$$C_j(X_1,\ldots,X_{n+m}) = \sum (\gamma_j \prod_{k=1}^{n+m} X_k^{f_{jk}}) = 0$$

with a called rate constraints



An Artificial Clock



- Three proteins:
 - LacI, tetR & \lambda cI
 - Arranged in a cyclic manner (logically, not necessarily physically) so that the protein product of one gene is rpressor for the next gene.

LacI $\rightarrow \neg$ *tetR; tetR* \rightarrow TetR

 $\mathsf{TetR} \to \neg \lambda cI; \lambda cI \to \lambda cI$

 $\lambda \mathbf{cI} \rightarrow \neg \mathbf{lacI}; \mathbf{lacI} \rightarrow \mathbf{LacI}$

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uet et al., Antoniotti et al., Wigler & Mishra



Cycles of Repression



- The first repressor protein, LacI from E. coli inhibits the transcription of the second repressor gene, tetR from the tetracyclineresistance transposon Tn10, whose protein product in turn inhibits the expression of a third gene, cI from I phage.
- Finally, CI inhibits lacI expression,
- completing the cycle.



Biological Model





- Standard molecular biology: Construct
 - A low-copy plasmid encoding the repressilator and
 - A compatible higher-copy reporter plasmid containing the tetrepressible promoter
 PLtet01 fused to an intermediate stability variant of gfp.



$$\begin{aligned} dx_2/dt &= \alpha_2 X_6^{g26} X_1^{g21} - \beta_2 X_2^{h22} \\ dx_4/dt &= \alpha_4 X_2^{g42} X_3^{g43} - \beta_4 X_4^{h44} \\ dx_6/dt &= \alpha_6 X_4^{g64} X_5^{g65} - \beta_6 X_6^{h66} \\ X_1, X_3, X_5 &= \text{const} \end{aligned}$$









Simpathica System

Sustem: Biology





 $\begin{array}{l} \text{IX} \left[X(s(t)) \land f(s(t), s(t + \land t) \land t) \Rightarrow X(s(t+ \land t)) \right] \\ \text{Made by A-PDF PPT2PDF} \end{array}$





- Ritt-Kolchin: Ideal Theoretic Approach
- Kolchin-Singer: Galois-Theoretic Approach
- Lie: Group-Theoretic Approach
- Understanding their interrelationship
- Effectiveness of various approaches



Differential Algebra

Assume that the system (SISO) is described as shown below:

$$\dot{x}_1 = p_1(X, u, \dot{u}, \dots, u^{(k)})$$

$$\vdots$$

$$\dot{x}_r = p_r(X, u, \dot{u}, \dots, u^{(k)})$$

$$0 = q_1(X, u)$$

$$\vdots$$

$$0 = q_s(X, u)$$

$$y = h(X, u)$$

Consider the following differential ideal I in the differential ring $\mathbb{R}\{X, u, y\}$:

$$I = [\dot{x}_1 - p_1, \dots, \dot{x}_r - p_r, q_1, \dots, q_s, y - h].$$

The input-output relation is then obtained by finding the contraction I^c of the ideal I to the ring $\mathbb{R}\{u, y\}$. The generators of $I^c = I \cap \mathbb{R}\{u, y\}$ give the differential polynomials involving u and y. However, the underlying algorithmic questions for

ain largely unsolved.

Systems Diviogy

Example System



and the output y is simply [B]:

$$y = [B].$$

We can simplify the above system to a polynomial system by following transformations:

$$x_1^2 = [A]$$
 and $x_2^2 = [B]$.





Thus,

$$I = [2x_1^5 \dot{x}_1 + x_1 - u, 2x_2 \dot{x}_2 + x_2 - x_1, x_2^2 - y].$$

After eliminating x_1 and x_2 , we obtain the following input-output relation:

$$\begin{array}{l} (20\dot{y}^8y^2 - 4\dot{y}^{10}y - 40\dot{y}^6y^3 + 40\dot{y}^4y^4 - 20\dot{y}^2y^5 + 4y^6)\ddot{y}^2 \\ + (4u\dot{y}^5y - 4\dot{y}^6y - 20\dot{y}^4y^2 + 40u\dot{y}^3y^2 + 20\dot{y}^2y^3 + 20u\dot{y}y^3 + 4y^4)\ddot{y} \\ - \dot{y}^2y^5 + 5\dot{y}^4y^4 - 10\dot{y}^6y^3 + 20u\dot{y}^3y^2 + 10\dot{y}^8y^2 + y^2 - 8\dot{y}^6y + 10u\dot{y}^5y \\ - u^2y + 2u\dot{y}y - \dot{y}^2y - 5\dot{y}^{10}y + \dot{y}^{12} + 8\dot{y}^2y^3 + 2u\dot{y}y^3 = 0.\Box \end{array}$$







- Various Approaches:
 - Ideas based on the H-bases (Gröbner Bases).
 - Ideas based on Ritt's Characteristic Sets.
 - Obstacles: Failure of a Hilbert-basis like theorem (only a weaker version, Ritt-Raudenbusch Basis Theorem, holds), existence of non-recursive differential ideals, etc.



Simpler Computational Models



- Kripke Models/Discrete Event Systems
- Hybrid Automata
- Their Connection to
 - Turing Machines
 - "Real" Turing Machines



Kripke Structure





- Formal Encoding of a Dynamical System:
- Simple and intuitive pictorial representation of the behavior of a complex system
 - A Graph with nodes representing system states labeled with information true at that state
 - The edges represent system transitions as the result of ome action



Computation Tree



asiens Sinion

- Finite set of states;
 Some are initial states
- Total transition relation: every state has at least one next state i.e. infinite paths
- There is a set of basic environmental variables or features ("atomic propositions")
- In each state, some atomic propositions are true



Hybrid Automata



- 1. $Z = \langle Z_1, ..., Z_k \rangle$ and $Z' = \langle Z'_1, ..., Z'_k \rangle$ are two vectors of variables ranging over the reals IR;
- 2. $(\mathcal{V}, \mathcal{E})$ is a graph; the objects, $v \in \mathcal{V}$, are called locations;
- 3. Each vertex $v \in V$ is labeled by the formula Inv(v)[Z];
- 4. \mathcal{F} is a function assigning to each vertex $v \in V$ a continuous vector field over \mathbb{R}^k ; we will use $f_v : \mathbb{R}^k \times \mathbb{R}^+ \longrightarrow \mathbb{R}^k$ to indicate the solution of the vector field $\mathcal{F}(v)$ and Dyn(v)[Z, Z', T] to identify the corresponding formula, i.e., $Dyn(v)[Z, Z', T] \equiv Z' = f_v(Z, T)$;
- 5. Each edge $e \in \mathcal{E}$ is labeled by the two formulæ Act(e)[Z] and Reset(e)[Z, Z']; $\overline{Reset}(e)[Z'] \stackrel{\text{def}}{=} \exists Z \operatorname{Reset}(e)[Z, Z'].$

Thermostat



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System: Dislogy

Intuition



Hybrid Automata - Intuitively

Intuitively, a hybrid automaton is a finite state automaton H with continuous variables X





Semantics

Hybrid Automata - Sematics



E

 $Act(\langle v, v' \rangle)[r],$

and

and

Hybrid Automata - Sematics

$\frac{r}{f(t')} = \frac{1}{s}$ $\frac{r}{r} = \frac{r}{r} =$



 $\sigma \equiv$ switching signal taking values in the set {1,2}





Chemotaxis



- Escherichia coli has evolved a strategy for responding to a chemical gradient in its environment
 - It detects the concentration of ligands through a number of receptors
 - It reacts by driving its flagella motors to alter its path of motion.
 - Either it "runs" moves in a straight line by moving its flagella counterclockwise (CCW), or it "tumbles" – randomly change its heading by moving its flagella clockwise (CW).
- The response is mediated through the molecular concentration of CheY in a phosphorylated form, which in turn is determined by the bound ligands at the receptors that appear in several forms.
- The more detailed pathway involves other
 - CheB (either with phosphorylation or without, B_p and B_0),
 - CheZ (Z),
 - bound receptors (LT) and
 - unbound receptors (T)
- Their continuous evolution is determined by a set of differential

rough kinetic mass action formulation.



 $t>100\wedge t'=0\wedge \theta'=0\wedge Y'_P=Y_P\wedge B'_P=B_P\wedge P'=P\wedge Z'=Z$

Fig. 2. An IDA capturing the run-tumble mechanism of E. coli.





Questions of Interest

- Controllability:
 - Assume that the system is at the "origin" initially. Can we find a control signal so that the state reaches a given position at a fixed time?
- Observability:
 - Can the state x be determined from observations of the output y over some time interval.
- Reachability: A computationally simpler problem:
 - Can we determine what states are reached as the system evolves autonomously or under a class of control signal.
- "HALTING" Problem:
 - Can the system reach a designated state at some time and then stay there?




Dynamics



Replacing differential equations by "equivalent" dynamics:

If f(X, T) is the solution of $X = \mathcal{F}(X, T)$, then

 $\dot{X} = \mathcal{F}(X, T)$ and $Dyn[X, X', T] \equiv X' = f(X, T)$,

are equivalent

Inclusion DynamicsWe are interested in inclusion dynamics defined by formulæ $f(r,2\delta)$ $f(r,2\delta)$ $f(r,2\delta)$ $f(r,3\delta)$ $f(r,3\delta)$ $f(r,4\delta)$ $F(r,3\delta)$ $F(r,4\delta)$ $F(r,4\delta)$



Michael's Form



– Let $F_X^v(T) \equiv \{X' \mid Dyn(v)[X, X', T] \land Inv(v)[X']\}$

- A Hybrid automaton is in Michael's form if
 - F_{x}^{v} is lower semi-continuous
 - For each $t\in I_X{}^{\vee}$ the set $F_{\!\!X}{}^{\vee}(t)$ is closed and convex
 - where I_X^{V} is the largest [0, t') such that $F_X^{V}(t) \neq \emptyset$, $\forall t \in [0, t')$.

Theorem

If H is in Michael's form, then $s \in F_r^v(t)$ iff $\langle v, r \rangle \xrightarrow{t}_C \langle v, s \rangle$.

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Reachability



Michael's Form and Reachability

For each automaton in Michael's form, we can write a formula Reach(H, ph)[X, X', T], where $ph = v_0, \ldots, v_n$ is a path on $\langle \mathcal{V}, \mathcal{E} \rangle$, such that

Reach(H, ph)[X, X', T] holds H reaches $\langle v_n, X' \rangle$ from $\langle v_0, X \rangle$ with a trajectory corresponding to *ph*

ist not be infinite!!





FOCoRe and IDA

FOCoRe (First Order Constant Reset hybrid automata) are first-order hybrid automata:

- in Michael's form
- with constant resets (i.e., *Reset(e)[X, X']* does not depend on X)

IDA (Independent Dynamics hybrid Automata) allows identity resets between locations whose dynamic does not change

We can reduce reachability problem for either FOCoRe or IDA T-automata to a satisfiability problem for formulæ of T



First Order Theory of Reals



- Tarski's theorem says that the first-order theory of reals with +, ×, =, and > allows quantifier elimination. Algorithmic quantifier elimination implies decidability.
- Every quantifier-free formula composed of polynomial equations and inequalities, and Boolean connectives defines a semialgebraic set. Thus a set S is semi-algebraic if:

$$S = \{ \langle \xi_1, \cdots, \xi_n \rangle \in \mathcal{R}^n | \psi(\xi_1, \cdots, \xi_n) = True \}, or$$

$$S = \bigcup_{i=1}^{I} \bigcap_{j=1}^{J_i} \{ \langle \xi_1, \cdots, \xi_n \rangle \in \mathcal{R}^n | sign(f_{i,j}(\xi_1, \cdots, \xi_n)) = s_{i,j} \}$$

where $\psi(\xi_1, \dots, \xi_n)$ is a quantifier-free formula involving n algebraic variables, $f_{i,j}s$ are multivariate polynomials over R and the $s_{i,j}s$ are -1, 0, +1.







- Hybrid Automata's inclusion dynamics, approximated by semi-algebraic formula. Dyn[X,X', T] ≡ Semialgebraic Set
- A more realistic approximation, for time invariant systems:
 - Dyn[X, X', h]

 $\approx \{X' \mid X' = X + \mathcal{F}(X,0) h + \delta, |\delta| < \epsilon\},\$

for a suitably chosen

 $\epsilon \equiv |\mathcal{F}(X,0) h^2/2! + \mathcal{F}(X,0) h^3/3! + \cdots |$



Another Example: Biological Pattern Formation



Figure 3: *Xenopus* embryo labeled by a marker for ciliated cell precursors seen as black dots.¹

- Embryonic Skin Of The South African Claw-Toed Frog
- "Salt-and-Pepper" pattern formed due to lateral inhibition in the Xenopus epidermal layer





Physically **adjacent** cells **laterally inhibit** each other's ciliation (Delta production)



Delta-Notch Pathway





Delta binds and activates its receptor Notch in neighboring cells (proteolytic release and nuclear translocation of the intracellular domain of Notch)

- Activated Notch suppresses ligand (Delta) production in the cell
- A cell producing more ligands forces its neighboring cells to produce less



Pattern formation by lateral inhibition with feedback: a mathematical model of Delta-Notch intercellular signalling Collier et al.(1996)

$$\frac{\mathrm{d}(N_P/N_0)}{\mathrm{d}\tau} = F(\bar{D}_P/D_0) - \mu N_P/N_0,$$

$$\frac{\mathrm{d}(D_P/D_0)}{\mathrm{d}\tau} = G(N_P/N_0) - \rho D_P/D_0.$$

Rewriting...

$$\dot{n}_P = f(\bar{d}_P) - n_P,$$

$$\dot{d}_P = v\{g(n_P) - d_P\}.$$

Where:

$$f(x) = \frac{x^k}{a + x^k}, \ g(x) = \frac{1}{1 + bx^k},$$



FIG. 1. Diagrammatic representation of the effective feedback loop between Notch and Delta in neighbouring cells. Details of the Notch signalling pathway are omitted for clarity. Key: -> Delta; -- Kotch.



One-Cell Delta-Notch Hybrid Automaton





(a) Transition diagram for a single cell Made by A-PDF PPT2PDF $H_{one_cell} = (Q, X, \Sigma, Init, f, Inv, R)$

 $Q = q_1, q_2, q_3, q_4$ $X = (x_1, x_2)^T \in \Re^2$ $\Sigma = \left\{ u_N = \sum_{i=1}^{6} x_{Delta,i} \right\}$ $Init = Q \times \{X \subset \Re^2 : x_1, x_2 > 0\}$ $f(q, x) = \begin{cases} [-\lambda_D x_1; -\lambda_N x_2]^T & \text{if } q = q_1 \\ [R_D - \lambda_D x_1; -\lambda_N x_2]^T & \text{if } q = q_2 \\ [-\lambda_D x_1; R_N - \lambda_N x_2]^T & \text{if } q = q_3 \\ [R_D - \lambda_D x_1; R_N - \lambda_N x_2]^T & \text{if } q = q_4 \end{cases}$ $Inv = \{q_1, \{-x_2 < h_D, u_N < h_N\}\} \cup$ $\{q_2, \{-x_2 \ge h_D, u_N < h_N\}\} \cup$

$$\{q_3, \{-x_2 < h_D, u_N \ge h_N\}\} \cup \{q_4, \{-x_2 \ge h_D, u_N \ge h_N\}\}$$

Ghosh et al.





Fig. 7. Phase plane projections for two cell system showing equilibria. Labels d_1 and d_2 are the Delta protein concentrations in cell 1 and 2 respectively.



State Reachability

Reaching State q_7 (2,3) When we ask $True \exists \mathcal{U} [-2n_1 > -1 \land 5d_2 < 1 \land -2n_2 < -1 \land 5d_1 > 1]$, we get: Iteration 1: $5d_1 - 1 \ge 0 \land 2n_1 - 1 \le 0 \land 5d_2 - 1 \le 0 \land 2n_2 - 1 \ge 0$ Iteration 2: $n_1 - 1 \le 0 \land [[2n_1 - 5d_1 \le 0 \land 5d_2 - 1 \le 0 \land 8n_2 - 5d_2 - 3 \ge 0 \land n_2 + n_1 - 1 = 0] \lor [8n_1 - 5d_1 - 3 \le 0 \land 4d_2 + d_1 - 1 = 0 \land 2n_2 - 1 \ge 0 \land 8n_2 + 5d_1 - 5 \ge 0] \lor [5d_1 - 1 \ge 0 \land 2n_1 - 5d_1 \le 0 \land 5d_2 + 2n_1 - 2 \le 0 \land 2n_2 - 1 \ge 0] \lor [5d_1 - 1 \ge 0 \land 2n_1 - 1 \le 0 \land 5d_2 - 1 \le 0 \land 8n_2 - 5d_2 - 3 \ge 0] \lor [2n_1 - 1 \le 0 \land 2n_1 - 1 \le 0 \land 5d_2 - 1 \le 0 \land 8n_2 - 5d_2 - 3 \ge 0] \lor [2n_1 - 1 \le 0 \land 5d_2 - 1 \le 0 \land 8n_2 + 5d_1 - 5 \ge 0] \lor [2n_1 - 1 \le 0 \land 5d_2 - 1 \le 0 \land 8n_2 - 5d_2 - 3 \ge 0 \land 8n_2 + 5d_1 - 5 \ge 0] \lor [2n_1 - 5d_1 \le 0 \land 5d_2 - 1 \le 0 \land 5d_2 - 1 \le 0 \land 8n_2 + 5d_1 - 5 \ge 0] \lor [2n_1 - 5d_1 \le 0 \land 5d_2 - 1 \le 0 \land 2n_2 - 1 \ge 0 \land 8n_2 + 5d_1 - 5 \ge 0]]$ $\equiv f_7$ (sav).



State Reachability

Notch C2 Notch C2 Delta

Reaching State q_{10} (3,2) When we ask *True* $\exists \mathcal{U} [-2n_1 < -1 \land 5d_2 > 1 \land -2n_2 > -1 \land 5d_1 < 1]$, we get: Iteration 1: $5d_1 - 1 \leq 0 \land 2n_1 - 1 \geq 0 \land 5d_2 - 1 \geq 0 \land 2n_2 - 1 \leq 0$ Iteration 2: $n_2 - 1 \leq 0 \land [[2n_1 - 1 \geq 0 \land 5d_2 + 8n_1 - 5 \geq 0 \land d_2 + 4d_1 - 1 = 0 \land 2n_2 + 5d_1 - 2 \leq 0] \lor [2n_1 - 1 < 0 \land 8n_1 - 5d_1 - 3 \geq 0 \land 5d_2 + 8n_1 - 5 \geq 0 \land n_2 + n_1 - 1 = 0] \lor [8n_1 - 5d_1 - 3 \geq 0 \land 5d_2 + 8n_1 - 5 < 0 \land 5d_2 + 2n_1 - 2 \geq 0 \land n_2 + n_1 - 1 = 0] \lor [2n_1 - 1 \geq 0 \land 5d_2 - 1 \geq 0 \land 2n_2 + 5d_1 - 2 \leq 0 \land n_2 + n_1 - 1 = 0] \lor [2n_1 - 1 \geq 0 \land 5d_2 - 1 \geq 0 \land 2n_2 + 5d_1 - 2 \leq 0 \land n_2 + n_1 - 1 = 0] \lor [5d_1 - 1 \leq 0 \land 2n_1 - 1 \geq 0 \land 5d_2 + 8n_1 - 5 \geq 0 \land 2n_2 - 5d_2 \leq 0] \lor [5d_1 - 1 \leq 0 \land 2n_1 - 1 \geq 0 \land 5d_2 + 8n_1 - 5 \geq 0 \land 2n_2 - 1 \leq 0] \lor [8n_1 - 5d_1 - 3 \geq 0 \land 5d_2 - 1 \geq 0 \land 2n_2 + 5d_1 - 2 \leq 0 \land 2n_2 - 1 \leq 0] \lor [8n_1 - 5d_1 - 3 \geq 0 \land 5d_2 - 1 \geq 0 \land 2n_2 + 5d_1 - 2 \leq 0 \land 2n_2 - 1 \leq 0] \lor [8n_1 - 5d_1 - 3 \geq 0 \land 5d_2 - 1 \geq 0 \land 2n_2 + 5d_1 - 2 \leq 0 \land 2n_2 - 1 \leq 0] \lor [8n_1 - 5d_1 - 3 \geq 0 \land 5d_2 - 1 \geq 0 \land 2n_2 + 5d_1 - 2 \leq 0 \land 2n_2 - 1 \leq 0] \lor [8n_1 - 5d_1 - 3 \geq 0 \land 5d_2 - 1 \geq 0 \land 2n_2 + 5d_1 - 2 \leq 0 \land 2n_2 - 1 \leq 0] \lor [8n_1 - 5d_1 - 3 \geq 0 \land 5d_2 - 1 \geq 0 \land 2n_2 + 5d_1 - 2 \leq 0 \land 2n_2 - 1 \leq 0] \lor [8n_1 - 5d_1 - 3 \geq 0 \land 5d_2 - 1 \geq 0 \land 2n_2 + 5d_1 - 2 \leq 0 \land 2n_2 - 1 \leq 0] \lor [8n_1 - 5d_1 - 3 \geq 0 \land 5d_2 - 1 \geq 0 \land 2n_2 + 5d_1 - 2 \leq 0 \land 2n_2 - 1 \leq 0] \lor [8n_1 - 5d_1 - 3 \geq 0 \land 5d_2 - 1 \geq 0 \land 2n_2 + 5d_1 - 2 \leq 0 \land 2n_2 - 1 \leq 0]$





 $\begin{array}{l} f_7 \wedge \neg f_{10} \,=\, n_1 - 1 \,\leq\, 0 \wedge \left[\left[2n_1 - 5d_1 \,\leq\, 0 \wedge 5d_2 - 1 \,< \\ 0 \wedge 8n_2 - 5d_2 - 3 \,\geq\, 0 \wedge n_2 + n_1 - 1 \,=\, 0 \right] \vee \left[2n_1 - 1 \,\leq\, \\ 0 \wedge 5d_2 - 1 \,\leq\, 0 \wedge 8n_2 - 5d_2 - 3 \,\geq\, 0 \wedge 2n_2 + 5d_1 - 2 \,>\, \\ 0 \right] \vee \left[2n_1 - 1 \,\leq\, 0 \wedge 5d_2 + 2n_1 - 2 \,\leq\, 0 \wedge 4d_2 + d_1 - 1 \,=\, \\ 0 \wedge n_2 + n_1 - 1 \,>\, 0 \right] \vee \left[2n_1 - 5d_1 \,\leq\, 0 \wedge 5d_2 - 1 \,\leq\, 0 \wedge n_2 + n_1 - 1 \,>\, \\ 0 \wedge 2n_2 - 1 \,\geq\, 0 \right] \vee \left[2n_1 - 1 \,\leq\, 0 \wedge 5d_2 - 1 \,\leq\, 0 \wedge 8n_2 - 5d_2 - 3 \,\geq\, \\ 0 \wedge 8n_2 + 5d_1 - 5 \,\geq\, 0 \right] \vee \left[8n_1 - 5d_1 - 3 \,<\, 0 \wedge 4d_2 + d_1 - 1 \,=\, \\ 0 \wedge 2n_2 - 1 \,\geq\, 0 \wedge 8n_2 + 5d_1 - 5 \,\geq\, 0 \right] \vee \left[5d_1 - 1 \,\geq\, 0 \wedge 2n_1 - 5d_1 \,<\, \\ 0 \wedge 5d_2 + 2n_1 - 2 \,\leq\, 0 \wedge 2n_2 - 1 \,\geq\, 0 \right] \vee \left[5d_1 - 1 \,\geq\, 0 \wedge 2n_1 - 5d_1 \,<\, \\ 0 \wedge 5d_2 - 1 \,<\, 0 \wedge 8n_2 - 5d_2 - 3 \,\geq\, 0 \right] \vee \left[2n_1 - 1 \,<\, 0 \wedge 5d_2 - 1 \,\leq\, \\ 0 \wedge 8n_2 - 5d_2 - 3 \,\geq\, 0 \wedge 8n_2 + 5d_1 - 5 \,\geq\, 0 \right] \right] \end{array}$

Since we have assumed no upper bound on the initial values and since we have been able to compute only two iterations, this formula does *not* evaluate to *True* given $n_1 < n_2 \land$ $d_1 > d_2$. However, when *Qepcad* simplifies the above formula assuming that $n_1 > n_2 \land d_1 < d_2$, it immediately evaluates





Logic & Model-Checking



Deciphering Design Principles in a Biological Systems

Step 1. Formally encode the behavior of the system as a hybrid automaton Step 2. Formally encode the properties of interest in a powerful logic Step 3. Automate the process of checking if the formal model of the system satisfies the formally encoded properties using Model Checking



Temporal Logic



- First Order Logic: Time is an explicitly quantified variable
- Propositional Modal logic: was invented to formalize modal notions and suppress the quantified variables – with operators "possibly P" and "necessarily P" (similar to "eventually" and "henceforth")





Branching versus Linear Time

- Temporal Logic:
 - Short hand for describing the way properties of the system change with time
 - Time is implicit
- Linear-time: Only one possible future in a moment
 - Look at individual computations
- Branching-time: It may be possible to split to different courses depending on possible futures
 - Look at the tree of computations



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Time is Branching



Computation Tree Logic (CTL)

- Branching Time temporal logic: interpreted over an *execution tree* where branching denotes non-deterministic actions
- Explicitly quantify over two modes the path and the time
- Each time we talk about a temporal property, we also specify whether it is true on all possible paths or whether it is true on at least one path - *Path quantifiers*
 - -A = "for all future paths"
 - E = "for some future path"

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 $\exists j \ge 0 [s_j \models g \text{ and } \forall i : 0 \le i < j [s_i \models f]]$ • $s \models EG f \Leftrightarrow \exists \pi = \langle s_0 s_1 \dots \rangle \text{ from } s \forall i \ge 0 : s_i \models f$

Semantics for CTL

- $s \models E(f \cup g) \Leftrightarrow \exists \pi = \langle s_0 s_1 \dots \rangle$ from s
- $s \models f \lor g \Leftrightarrow s \models f \text{ or } s \models g$ • $s \models EX f \Leftrightarrow \exists \pi = \langle S_0 S_1 \dots \rangle \text{ from } S s_1 \models f$
- $s \models p \Leftrightarrow p \in L(s)$ $s \models \neg p \Leftrightarrow p \notin L(s)$ • $s \models f \land g \Leftrightarrow s \models f \text{ and } s \models g$
- For p∈AP:





System Divings

Some CTL Operators



EFg AFg EGg AGg



CTL Model-Checking



Label the states with the terms in the formula:

 Proceed by marking each point with the set of valid subformulas

"Global" algorithm:

- Iterate on the structure of the property, traversing the whole of the model in each step
- Use fixed point unfolding to interpret Until:

 $\mathbf{E}(\psi_2 \mathbf{U}^+ \psi_1) \leftrightarrow \mathbf{E} \mathbf{X}(\psi_1 \lor \psi_2 \land \mathbf{E}(\psi_2 \mathbf{U}^+ \psi_1))$ $\mathbf{A}(\psi_1 \mathbf{U}^+ \psi_1) \leftrightarrow \mathbf{A} \mathbf{X}(\psi_1 \lor \psi_2 \land \mathbf{A}(\psi_2 \mathbf{U}^+ \psi_1))$



Naïve CTL Model-Checker





Other Model Checking Algorithms

- LTL Model Checking: Tableu-based...
- CTL* Model Checking: Combine CTL and LTL Model Checkers...
- Symbolic Model Checking
 - Binary Decision Diagram
 - OBDD-based model-checking for CTL
 - Fixed-point Representation
 - Automata-based LTL Model-Checking
- SAT-based Model Checking
- Algorithmic Algebraic Model Checking
- Hierarchical Model Checking





Purine Metabolism

- Provides the organism with building blocks for the synthesis of DNA and RNA.
- The consequences of a malfunctioning purine metabolism pathway are severe and can lead to death.
- The entire pathway is almost closed but also quite complex. It contains
 - several feedback loops,
 - cross-activations and
 - reversible reactions

 Thus is an ideal candidate for reasoning with computational tools.



Simple Model







Biochemistry of Purine Metabolism









- The main metabolite in purine biosynthesis is 5-phosphoribosyl-a-1pyrophosphate (PRPP).
 - A linear cascade of reactions converts PRPP into *inosine monophosphate* (*IMP*). IMP is the central branch point of the purine metabolism pathway.
 - IMP is transformed into AMP and GMP.
 - Guanosine, adenosine and their derivatives are recycled (unless used elsewhere) into *hypoxanthine* (*HX*) and *xanthine* (*XA*).
 - XA is finally oxidized into *uric acid* (*UA*).





Queries



- Variation of the initial concentration of PRPP does not change the steady state. (PRPP = 10 * PRPP1) implies steady_state()
- This query will be true when evaluated against the modified simulation run (i.e. the one where the initial concentration of PRPP is 10 times the initial concentration in the first run – PRPP1).
- Persistent increase in the initial concentration of PRPP does cause unwanted changes in the steady state values of some metabolites.
- If the increase in the level of PRPP is in the order of 70% then the system does reach a steady state, and we expect to see increases in the levels of IMP and of the hypoxanthine pool in a "comparable" order of magnitude.

Always (PRPP = 1.7*PRPP1) implies steady_state()



TRUE







- Consider the following statement:
- Eventually
- (Always (PRPP = 1.7 * PRPP1) implies steady_state() and Eventually
 - Always(IMP < 2* IMP1)) and Eventually (Always (hx_pool < 10*hx_pool1)))
- where IMP1 and hx_pool1 are the values observed in the unmodified trace. The above statement turns out to be false over the modified experiment trace..

- In fact, the increase in IMP is about 6.5 fold while the hypoxanthine pool increase is about 60 fold.
- Since the above queries turn out to be false over the modified trace, we conclude that the model "over-predicts" the increases in some of its products and that it should therefore be amended



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Spilens Sinlogy

ght solid arraws represent activation, while light dashed centwo retoring or loaving the pathway indicate parine ring and ribase is system.



Final Model



Purine Metabolism





Continuous-Time Logics

- Linear Time
 - Metric Temporal Logic (MTL)
 - Timed Propositional Temporal Logic (TPTL)
 - Real-Time Temporal Logic (RTTL)
 - Explicit-Clock Temporal Logic (ECTL)
 - Metric Interval Temporal Logic (MITL)
- Branching time
 - Real-Time Computation Tree Logic (RTCTL)
 - Timed Computation Tree Logic (TCTL)


TCTL: Syntax And Semantics



Basic Syntax And Semantics The basic syntax of TCTL is:

 $\phi ::= p \mid \neg \phi \mid \phi_1 \lor \phi_2 \mid \phi_1 \exists \mathcal{U} \phi_2 \mid \phi_1 \forall \mathcal{U} \phi_2 \mid z.\phi$

- z. The freeze quantification "z." binds the associated variable z to the current time. Thus the formula z.φ(z) holds at time t iff φ(t) does.
- ∀U and ∃U: universal and existential "until" operators. It is common notation to subscript the Until operator as p ∃U≤t q to indicate that q has to be satisfied within t time units. This is just a convenient notation for p ∃U (q ∧ (t ≤ 5)).

E.g. $(p)\forall \mathcal{U}(q)$ asks whether on any path leading off the state where the modal formula is being considered, p is true everywhere until the state where q is true. (q is required to be true somewhere, and p is required to be true until the previous instant, but not necessarily at the point where q becomes true)







 $\phi ::= X ~|~ p ~|~ \neg \phi ~|~ \phi_1 \lor \phi_2 ~|~ \phi_1 \triangleright \phi_2 ~|~ z.\phi ~|~ \mu X.\phi$

Note that though we mention only the least-fixpoint μ , the greatest-fixpoint ν can be expressed as $\neg \mu X.(\neg \phi[X := \neg X]).$





 $s_1 ~ \exists \mathcal{U} ~ s_2 ~=~ \mu X. (s_2 \lor (s_1 \triangleright X))$

- s2 is true now or
- s1 holds for one-step on some path after which s2 holds or
- **s1** holds for *one-step on some path* after which **s1** holds for *one more step on some path* after which **s2** holds **or**
- and so on..

Since the universal unti can be computed using $s_1 \forall U s_2 = \neg((\neg s_2) \exists U (\neg s_1 \land \neg s_2))$, it is sufficient to focus on $\exists U$.³



TCTL Model Checking



- Only "Until" requires "computation"
- Until: Iterative computation of "onestep" Until
- Least fixpoint computation:

1.
$$\psi := \text{false}$$

2. repeat
(a) $\phi := \psi$
(b) $\psi := \phi[X := \phi]$
(c) until $[\phi] = [\psi]$

3. return ϕ



Semi-Decidability Of TCTL



Global "time" variable

 Allows interpretation of the TCTL operators freeze (z.X) and subscripted until (Ua)

- While "one-step until" is decidable, the fixpoint is not guaranteed to converge
- So TCTL is "semi"-decidable



Mandelbrot Hybrid Automaton





 $S'(t) = \{x(t)^2 - y(t)^2 + C_r\} + i.\{2x(t)y(t) + C_i\} = \{x(t) + i.y(t)\}^2 + \{C_r + i.C_i\}$

Reachability Query: $(x_1^2 + x_2^2 \ge 4)$



Solution



- Bounded Model Checking
 - Fully O-minimal Systems for Dense CTL
- Constrained Systems
 - Linear Systems for Dense CTL
 - O-minimal for Dense CTL
 - SACoRe (Semi algebraic Constrained Reset) for TCTL
 - IDA (Independent Dynamics Automata) for TCTL





- "And though many things I have first Discovered could not find acceptance yet I finde there are not wanting some who pride themselves on arrogating of them for their own...
 - "-But I let that passe for the present."







- "So many are the links, upon which the true Philosophy depends, of which, if any can be loose, or weak, the whole chain is in danger of being dissolved;
- "it is to begin with the Hands and Eyes, and to proceed on through the Memory, to be continued by the Reason;
- "nor is it to stop there, but to come about to the Hands and Eyes again, and so, by a continuall passage round from one Faculty to another, it is to be maintained in life and strength."



The end...

