Sample Solution for Homework 2

Problem 1 AMP, p.41, Exercise 11: Flaky Lock (12 Points)

• The protocol satisfies mutual exclusion. For a proof by contradiction, suppose it did not. By inspecting the protocol code, we conclude that the following happens-before relationships must hold:

write_A(turn = A)
$$\rightarrow$$
 write_A(turn = B) \rightarrow CS_A

and

write_B(turn = B)
$$\rightarrow$$
 write_A(turn = A) \rightarrow CS_B

Assume without loss of generality that thread A was the last thread to write to turn before entering the critical section. Then

write_B(turn = B) \rightarrow write_A(turn = A)

which contradicts the fact that A completed its outer waiting loop.

- As we show below, the algorithm is not deadlock-free. Hence, it is also not starvation-free.
- The following sequence of events leads to a deadlock:

$$\operatorname{write}_A(\operatorname{turn} = A) \to \operatorname{read}_A(\operatorname{busy} = \operatorname{false}) \to \operatorname{write}_B(\operatorname{turn} = B)$$

Problem 2 AMP, p.41, Exercise 14: *l*-Exclusion (13 Points)

We can turn the filter algorithm into an algorithm that solves the ℓ -exclusion problem by reducing the number of levels by $\ell - 1$ (assuming $n > \ell$).

```
class LFilter implements Lock {
  int[] level;
 int[] victim;
 public LFilter(int n, int l) {
    level = new int[max(n-l+1,0)];
    victim = new int[max(n-l+1,0)];
    for (int i = 0; i < n-l+1; i++) {</pre>
      level[i] = 0;
    }
  }
 public void lock() {
    int me = ThreadID.get();
    for (int i = 1; i < n-l+1; i++) { // attempt level i</pre>
      level[me] = i;
      victim[i] = me;
      // spin while conflicts exist
      int above = 1+1;
      while (above > 1 && victim[i] == me) {
        above = 0;
        for (int k = 0; k < n; k++) {
          if (level[k] >= i) above++;
        }
      }
    }
  }
 public void unlock() {
    int me = ThreadID.get();
    level[me] = 0;
  }
}
```