

SOCIAL NETWORKS

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LECTURE # 19

A General Description of Evolutionarily Stable Strategies (ESS)

		Organism 2.	
		S	T
Organism 1	S	(a, a)	(b, c)
	T	(c, b)	(d, d)

Two-player Two-strategy Game
(Symmetric).

Condition for (S, S) to be evolutionarily stable

$S = ESS$

Let $\epsilon > 0$

$\left. \begin{array}{l} 1-\epsilon \text{ fraction uses } S \\ \epsilon \text{ fraction uses } T. \end{array} \right\} \begin{array}{l} \text{Expected payoff for} \\ \text{organism playing, } S \\ = a(1-\epsilon) + b\epsilon \end{array}$

$\left. \begin{array}{l} \text{Exp. payoff for org playing, } T \\ = c(1-\epsilon) + d\epsilon \end{array} \right\}$

We need the following inequality

$$a(1-\epsilon) + b\epsilon > c(1-\epsilon) + d\epsilon$$

$$\lim_{\epsilon \rightarrow 0} a > c$$

$$\therefore S = \text{ESS} \quad \text{if } a > c.$$

$$S \neq \text{ESS} \quad \text{if } a < c.$$

$$\text{if } a = c, \quad S = \text{ESS}, \quad \text{if } b > d.$$

Thm: In a two player, two-strategy, symmetric ^{game}, S is evolutionarily stable

when either (i) $a > c$

or (ii) $a = c$ and $b > d$.



Note: $(S, S) = \text{Nash Equilibria}$

$$u_1(S, S) \geq u_1(T, S) \quad \boxed{\equiv a > c}$$

$$\& \quad u_2(S, S) \geq u_2(S, T)$$

\Rightarrow

$$S = \text{ESS} \Rightarrow (S, S) = \text{N.E.}$$

Corollary:

If strategy S (in a two-player, two-strategy symmetric game) is ESS (evolutionarily stable), then $(S, S) =$ Nash equilibrium (P.S.N.E.)

Not all Nash Equilibria are E.S.S.

Stag-Hunt Game.

		Hunter 2.	
		Stag	Hare
Hunter 1	Stag	4,4	0,3
	Hare	3,0	3,3

$$a = 4$$

$$b = 0$$

$$c = 3$$

$$d = 3$$

$$a > c$$

$$d > b$$

Two ESS.

\swarrow Stag \searrow Hare.

Modify $c = 4$

	S	H	
S	4,4	0,4	$a = 4$
H	4,0	3,3	$c = 4$ $b = 0$ $d = 3$

$a = c, b < d$

\Rightarrow Stag \neq ESS.

But $d > b \Rightarrow$ Hare = ESS

$(H, H) = N.E.$

\approx

Note: ESS and N.E. are similar but underlying mechanisms differ

N.E \rightarrow Rationality (Fictitious Play)

E.S.S. \rightarrow Replicator Dynamics

Reasoning vs. Learning.

Evolutionarily Stable Mixed Strategies (ESMS)

"How to handle games in which no strategy is ESS."

Hawk-Dove Game

		Animal 2	
		D	H
Animal 1	D	(3, 3)	(1, 5)
	H	(5, 1)	(0, 0)

Symmetric
2-player
2-strategy

$$\left. \begin{array}{l} a=3 \\ b=1 \\ c=5 \\ d=0 \end{array} \right\} \begin{array}{l} a < c \\ d < b \end{array}$$

Hawk = Aggressive
Dove = Passive

- ◊ No ESS
- ◊ But it has two P.S.N.E.
= $\{(D, H), (H, D)\}$
- ◊ Neither D nor H is a best response to itself.
- ◊ Evolutionary Game in which strategy can be "mixed."
(by any individual)

Probability, $p \in [0,1]$

$q = 1-p$

$P(S) \sim B$

Play S with probability p
ET with " q } $S \sim \text{Ber}(p)$.

Organism 1 chooses S with prob p
X

Organism 2 chooses S with prob p'
Y

	X	Y
X	pp' (a, a)	pq' (b, c)
Y	qp' (c, b)	qq' (d, d)

Expected Pay-off for the player 1 \Rightarrow

$V(p, p') = pp'a + pq'b + qp'c + qq'd$

$$\left\{ \begin{aligned} V(p, p) &= p^2 a + 2pq(b+c) + q^2 d \\ V(p', p) &= pp'a + pq'c + qp'b + qq'd \\ V(p', p') &= p'^2 a + p'q'(b+c) + q'^2 d \end{aligned} \right.$$

$$p \neq p'$$

$\forall p \neq p' \exists \epsilon > 0 \forall x < \epsilon$

$$(1-x)V(p,p) + xV(p,p')$$

$$> (1-x)V(p',p) + xV(p',p')$$

In the general symmetric game,
 p is an evolutionarily stable mixed
 strategy,

if there is a (small) positive number ϵ
 such that, when any other mixed strategy
 p' invades p ($p \neq p'$) at any level
 $x < \epsilon$, the fitness of the organism
 playing p is strictly greater than
 the fitness of an organism playing
 p' .

$p = 1/3 \equiv$ Probability of playing Dove

To be shown: $p = 1/3 \equiv$ E.S.M.S.