



Programming Languages

OOP

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What is OOP? (part I)

The *object* idea:

- bundling of data (*data members*) and operations (*methods*) on that data
- restricting access to the data

An object contains:

- **data members** : arranged as a set of named fields
- **methods** : routines which take the object they are associated with as an argument
(known as *member functions* in C++)
- **constructors** : routines which create a new object

A class is a construct which defines the data, methods and constructors associated with all of its instances (objects).

The *inheritance* and *dynamic binding* ideas:

- classes can be extended (*inheritance*):
 - ◆ by adding new fields
 - ◆ by adding new methods
 - ◆ by *overriding* existing methods (changing behavior)

If class B extends class A, we say that B is a *subclass* or *derived* class of A, and A is a *superclass* or *base* class of B.

- dynamic binding : wherever an instance of a class is required, we can also use an instance of any of its subclasses; when we call one of its methods, the overridden versions are used.
- There should be an *is-a* relationship between a derived class and its base class.

Styles of OOLs

- in class-based OOLs, each object is an instance of a class (Java, C++, C#, Ada95, Smalltalk, OCaml, etc.)
- in prototype-based OOLS, each object is a clone of another object, possibly with modifications and/or additions (Self, Javascript)

Other common OOP features

- multiple inheritance
 - ◆ C++
 - ◆ Java (of interfaces only)
 - ◆ problem: how to handle diamond shaped inheritance hierarchy
- classes often provide package-like capabilities:
 - ◆ visibility control
 - ◆ ability to define types and classes in addition to data fields and methods

Java Features

- an imperative language (like C++, Ada, C, Pascal)
- is interpreted (like Scheme, APL)
- is garbage-collected (like Scheme, ML, Smalltalk, Eiffel, Modula-3)
- can be compiled
- is object-oriented (like Eiffel, more so than C++, Ada)
- a successful hybrid for a specific-application domain
- a reasonable general-purpose language for non-real-time applications

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- Work in progress: language continues to evolve
 - C# is latest, incompatible variant

Original design goals (white paper 1993)

- simple
- object-oriented (inheritance, polymorphism)
- distributed
- interpreted
- multi-threaded
- robust
- secure
- architecture-neutral

Obviously, “simple” was dropped.

Critical concern: write once – run everywhere

Consequences:

- portable interpreter
- definition through virtual machine: the JVM
- run-time representation has high-level semantics
- supports dynamic loading
- high-level representation can be queried at run-time to provide reflection
- dynamic features make it hard to fully compile, safety requires numerous run-time checks

Contrast with conventional systems languages

Conventional imperative languages are fully compiled:

- run-time structure is machine language
- minimal run-time type information
- language provides low-level tools for accessing storage
- safety requires fewer run-time checks because compiler (least for Ada and somewhat for C++) can verify correctness statically
- languages require static binding, run-time image cannot be easily modified
- different compilers may create portability problems

Notable omissions

- no operator overloading (syntactic annoyance)
- no separation of specification and body
- no enumerations until latest language release
- no generic facilities until latest language release

Statements

Most statements are like their C counterparts:

- `switch` (including C's falling through behavior)
- `for`
- `if`
- `while`
- `do ... while`
- `break` and `continue`
 - ◆ Java also has *labeled* versions of `break` and `continue`, like Ada.
- `return`

Java has no `goto`!

The simplest Java program

```
class HelloWorld {  
    public static void main (String[] args) {  
        System.out.println("Hello, □world");  
    }  
}
```

Encapsulation of type and related operations

```
class Point {
    private double x, y;    // private data members

    public Point (double x, double y) { // constructor
        this.x = x;    this.y = y;
    }

    public void move (double dx, double dy) {
        x += dx;    y += dy;
    }

    public double distance (Point p) {
        double xdist = x - p.x, ydist = y - p.y;
        return Math.sqrt(xdist * xdist + ydist * ydist);
    }

    public void display () { ... }
}
```

Extending a class

```
class ColoredPoint extends Point {
    private Color color;

    public ColoredPoint (double x, double y,
                        Color c) {
        super(x, y);
        color = c;
    }

    public ColoredPoint (Color c) {
        super(0.0, 0.0);
        color = c;
    }

    public Color getColor () { return color; }

    public void display () { ... } // now in color!
}
```

Dynamic dispatching

```
Point p1 = new Point(2.0, 3.0);
ColoredPoint cp1 = new ColoredPoint(2.0, 3.0, Blue);

Point p2 = p1;           // OK
Point p3 = cp1;         // OK

ColoredPoint cp2 = cp1; // OK
ColoredPoint cp3 = p1;  // Error

cp1.move(1.0, 1.0);    // cp1 and p3 affected

p1.display();          // Point's display
cp1.display();         // ColoredPoint's display
p3.display();          // ColoredPoint's display
```

Classes in C++

The same classes, translated into C++:

```
class Point {
    double m_x, m_y;    // private data members

public:

    Point (double x, double y)    // constructor
        : m_x(x), m_y(y) { }

    virtual ~Point () { }

    virtual void move (double dx, double dy) {
        m_x += dx;    m_y += dy;
    }

    virtual double distance (const Point& p) {
        double xdist = m_x - p.m_x, ydist = m_y - p.m_y;
        return sqrt(xdist * xdist + ydist * ydist);
    }

    virtual void display () { ... }
};
```


Extending a class

```
class ColoredPoint : public Point {
    Color color;

public:

    ColoredPoint (double x, double y,
                 Color c) : Point(x, y), color(c) {
        color = c;
    }

    ColoredPoint (Color c) : Point(0.0, 0.0), color(c) { }

    virtual Color getColor () { return color; }

    virtual void display () { ... } // now in color!
};
```

Dynamic dispatching

```
Point *p1 = new Point(2.0, 3.0);
ColoredPoint *cp1 = new ColoredPoint(2.0, 3.0, Blue);

Point *p2 = p1;           // OK
Point *p3 = cp1;         // OK

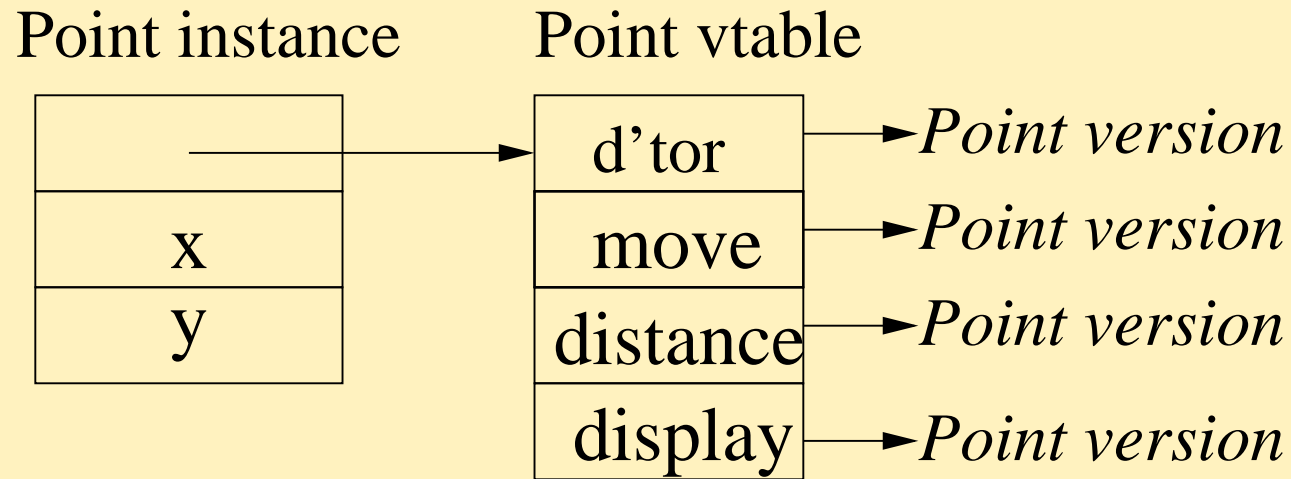
ColoredPoint *cp2 = cp1; // OK
ColoredPoint *cp3 = p1;   // Error

cp1->move(1.0, 1.0);     // cp1 and p3 affected

p1->display();           // Point's display
cp1->display();          // ColoredPoint's display
p3->display();           // ColoredPoint's display
```

Implementation: the vtable

A typical implementation of a class in C++; using `Point` as an example:



An extended vtable

For ColoredPoint, we have:

ColoredPoint instance

x
y
color

ColoredPoint vtable

d'tor	→ <i>ColoredPoint version</i>
move	→ <i>Point version</i>
distance	→ <i>Point version</i>
display	→ <i>ColoredPoint version</i>
getColor	→ <i>ColoredPoint version</i>

Non-virtual member functions are never put in the vtable

Method modifiers

- access modifiers:
 - ◆ `public`
 - ◆ `protected`
 - ◆ `package`
 - ◆ `private`
- `abstract`
- `static`
- `final`
- `synchronized`
- `native`
- `strictfp` (strict floating point)

A new construct: interfaces

A Java **interface** allows otherwise unrelated classes to satisfy a given requirement.

This is orthogonal to inheritance.

- **inheritance**: an **A** *is-a* **B** (has the attributes of a **B**, and possibly others)
- **interface**: an **A** *can-do* **X** (and possibly other unrelated actions)
- interfaces are a better model for multiple inheritance

See blackboard for implementation details (also in Scott, section 9.4.3)

Interface Comparable

```
public interface Comparable {  
    public int CompareTo (Object x) throws  
        ClassCastException;  
    // returns -1 if this < x,  
    //           0 if this = x,  
    //           +1 if this > x  
};
```

```
// Implementation needs to cast x to the proper class.
```

```
// Any class that may appear in a container should  
// implement Comparable, so the container can support  
// sorting.
```

Comparison with C++

Java	C++
methods	virtual member functions
public/protected/private members	similar
static members	same
abstract methods	pure virtual member functions
<code>final</code> methods	no analogous feature
<code>interface</code>	pure virtual class with no data members
implementation of an interface	virtual inheritance

Simulating a first-class function with an object

A simple first-class function:

```
fun mkAdder nonlocal = (fn arg => arg + nonlocal)
```

The corresponding C++ class:

```
class Adder {  
    int nonlocal;  
public:  
    Adder (int i) : nonlocal(i) { }  
    int operator() (int arg) { return arg + nonlocal; }  
};
```

`mkAdder 10` is roughly equivalent to `Adder(10)`.

First-class functions strike back

A simple unsuspecting object (in Java, for variety):

```
class Account {
    private float theBalance;
    private float theRate;

    Account (float b, float r) { theBalance = b;
                                theRate = r; }

    public void deposit (float x) {
        theBalance = theBalance + x;
    }
    public void compound () {
        theBalance = theBalance * (1.0 + rate);
    }
    public float balance () { return theBalance; }
}
```

First-class functions strike back, part 2

The corresponding first-class function:

```
(define (Account b r)
  (let ((theBalance b) (theRate r))
    (lambda (method)
      (case method
        ((deposit)
         (lambda (x) (set! theBalance
                           (+ theBalance x))))
        ((compound)
         (set! theBalance (* theBalance
                              (+ 1.0 theRate))))
        ((balance)
         theBalance))))))
```

`new Account(100.0, 0.05)` is roughly equivalent to
`(Account 100.0 0.05)`.

Comparing ML datatypes with inheritance

ML datatypes and OO inheritance organize data and routines in orthogonal ways:

	data variants	data operations
datatypes	all together/closed	scattered/open
classes	scattered/open	all together/closed

datatypes	easy to add new operations harder to add new variants
classes	easy to add new variants harder to add new operations

OOP Pitfalls: the circle and the ellipse

A couple of facts:

- In mathematics, an ellipse (from the Greek for absence) is a curve where the sum of the distances from any point on the curve to two fixed points is constant. The two fixed points are called foci (plural of focus).
from <http://en.wikipedia.org/wiki/Ellipse>
- A circle is a special kind of ellipse, where the two foci are the same point.

If we need to model circles and ellipses using OOP, what happens if we have class `Circle` inherit from class `Ellipse`?

Circles and ellipses

```
class Ellipse {  
    ...  
  
    public move (double dx, double dy) { ... }  
  
    public resize (double x, double y) { ... }  
}
```

```
class Circle extends Ellipse {  
    ...  
  
    public resize (double x, double y) { ??? }  
}
```

We can't implement a `resize` for `Circle` that lets us make it asymmetric!

Pitfalls: Array subclassing

In Java, if class B is a subclass of class A, then Java considers “array of B” to be a subclass of “array of A”:

```
class A { ... }  
class B extends A { ... }
```

```
B[] b = new B[5];  
A[] a = b;           // allowed (a and b are now aliases)
```

```
a[1] = new A();     // Bzzzt! (Type error)
```

The problem is that arrays are *mutable*; they allow us to replace an element with a different element.