1) Consider a computer system that has a cache with 4096 blocks. Each block can store 16 bytes. What will be the value stored in the TAG field of the cache block that holds the memory block containing the address 0xABCDEF (This hexadecimal number gives you a hint on the address length):

(i) [2 points] if it is a direct-mapped cache

- block size = 16 = $2^4$ so offset is 4 bits -> 1 hexadecimal digit
- #sets = #blocks/associativity = 4096/1 = $2^{12}$ so sets field is 12 bits -> 3 hex digits
- Tag is then the leftmost two digits from the address = 0xAB

(ii) [2 points] if it is a 16-way set-associative cache

- block size is same as above
- #sets = 4096/16 = $2^8$ so sets field is 8 bits -> 2 hex digits
- Tag is then the leftmost 3 digits -> 0xABC

(iii) [1 point] if it is fully associative

- block size same as above
- No sets field
- TAG is then all the address except the rightmost hex digit -> 0xABCDE
2. Consider a small 2-way set associative cache with a total of 32 blocks and a block size of 256 bytes. The cache uses LRU replacement policy. Assume that the cache is initially empty. The CPU accesses the following memory locations, in that order: 0x55c88, 0x55774, 0x5479c, 0x54c00, 0x55784, 0x56c80, 0x56718, 0x54738.

(a) [3 points] How the address is split to do cache lookup?

<table>
<thead>
<tr>
<th>Tag</th>
<th>Set</th>
<th>Offset</th>
<th>H/M</th>
<th>Comments</th>
</tr>
</thead>
<tbody>
<tr>
<td>55</td>
<td>c</td>
<td>88</td>
<td>M</td>
<td>Compulsory miss</td>
</tr>
<tr>
<td>55</td>
<td>7</td>
<td>74</td>
<td>M</td>
<td>Compulsory miss</td>
</tr>
<tr>
<td>54</td>
<td>7</td>
<td>9c</td>
<td>M</td>
<td>Compulsory miss</td>
</tr>
<tr>
<td>54</td>
<td>c</td>
<td>00</td>
<td>M</td>
<td>Compulsory miss</td>
</tr>
<tr>
<td>55</td>
<td>7</td>
<td>84</td>
<td>H</td>
<td></td>
</tr>
<tr>
<td>56</td>
<td>c</td>
<td>80</td>
<td>M</td>
<td>Replaces block with Tag 55</td>
</tr>
<tr>
<td>56</td>
<td>7</td>
<td>18</td>
<td>M</td>
<td>Replaces block with Tag 54</td>
</tr>
<tr>
<td>54</td>
<td>7</td>
<td>38</td>
<td>M</td>
<td>Replaces block with Tag 55</td>
</tr>
</tbody>
</table>

(b) [8 points] For each memory reference, indicate whether it will result in hit or miss and, if miss, indicate the type of miss (compulsory, capacity, or conflict)

(c) [1 point] From your solution in part (b) above, what is the hit rate of the cache?

The memory is accessed 8 times and had only 1 hit, as indicated by the table above. This gives a hit rate of $1/8 = 0.125$

(d) [3 points] Assume the memory access latency is 10 cycles and the cache, mentioned in this problem, has an access latency of 3 cycles. Also assume the hit rate
you calculated in part (c) above. Did this system benefit from having a cache?

Justify

(partial credits for this problem are noted next to the arrows)

If we did not have cache, then each access will take 10 cycles. $\leftarrow 0.5 \text{ point}$

With a cache, the average access time $= m + (1-p)M = 3 + (1-0.125)10 = 11.75 \leftarrow 1.5 \text{ point}$

This is larger than accessing the memory directly. So the system did not benefit from the cache. $\leftarrow 1 \text{ point}$
3. Almost all programs need external libraries.
   (a) In how many places can needed libraries be added to your code?
   
   In 3 places

   (b) What are these places?
   
   - statically during linking
   - dynamically during loading
   - dynamically during execution

   (c) For each one of those places, state one advantage and one disadvantage.
   
   - statically during linking
     - +: The executable is self-content. So, if you run it on a different machine, you do not have to worry about the version of the installed libraries or whether they are installed.
     - -: The executable file becomes much larger in size.
     - -: Is not very efficient in terms of memory usage.

   - dynamically during loading
     - +: The executable file is smaller.
     - +: Makes better use of the memory because libraries are shared among programs.
     - -: Loading becomes slower.

   - dynamically during execution
     - + Better memory usage because libraries are loaded only when needed.
     - +: loading time is not slower.
     - +: The executable file is not large.
     - -: Puts a burden on the programmer.
     - -: Slowdown during execution to load needed libraries
4. int array1[M][N];
   int array2[N][M];

   int copy(int i, int j)
   {
      array1[i][j] = array2[j][i];
   }

Suppose the above code generates the following assembly code (assume array2 and array1 are the base addresses of the corresponding arrays):

   copy:
   movl %rdi, %ecx  \(\rightarrow ecx = i\)
   movl %rsi, %ebx  \(\rightarrow ebx = j\)
   leal (%ecx,%ecx,8), %edx  \(\rightarrow edx = 9*ecx = 9i\)
   sall $2, %edx   \(\rightarrow edx = edx*4 = 36i\)
   movl %ebx, %eax  \(\rightarrow eax = ebx = j\)
   sall $4, %eax   \(\rightarrow eax = eax*16 = 16j\)
   subl %ebx, %eax  \(\rightarrow eax = eax – ebx = 16j – j = 15j\)
   sall $2, %eax    \(\rightarrow eax = eax *4 = 60j\)
   movl array2(%eax,%ecx,4), %eax \(\rightarrow eax = arra2+60j+4i\)
   movl %eax, array1(%edx,%ebx,4) \(\rightarrow array1+36i +4j = eax\)
   ret

What are the values of M and N: \(M = 15\) and \(N = 9\)

Show how did you reach your answer:
- An int is 4 bytes
- See the explanation next to the assembly code above
- Accessing \(array1[i][j] = array1 + (\text{#elements in a row} * 4*i) + (j*4) = array1+4Ni+4j\) \(\rightarrow array2+36i + 4j \rightarrow N = 9\)
- Same \(array2[j][i] = array2 +4Mj+4j \rightarrow array2+60j + 4i \rightarrow M = 15\)
5. From the assembly code and since both fun1 and fun2 have the same arguments in the same order:
   - edx \rightarrow a
   - eax \rightarrow b
   - By convention, the result of the function is returned in eax
   - cmpl %eax, %edx \rightarrow edx – eax \rightarrow a – b
   - Combining the above bullet with jg2 .L9 we find that it does:
     \text{if}(a \geq b) \text{ return } b \rightarrow \text{if}(a < b) \text{ return } b \rightarrow \text{so fun1}

6.

<table>
<thead>
<tr>
<th>Hexadecimal</th>
<th>binary</th>
<th>Decimal (assuming unsigned)</th>
<th>Decimal (assuming signed)</th>
</tr>
</thead>
<tbody>
<tr>
<td>0x8A</td>
<td>1000 1010</td>
<td>2+8+128 = 138</td>
<td>-(2s compl of num)</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>= -(01110110)</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>= -(64+32+16+4+2)</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td>= - 118</td>
</tr>
<tr>
<td>0x21</td>
<td>0010 0001</td>
<td>32+1 = 33</td>
<td>33</td>
</tr>
</tbody>
</table>