Final Exam Computational Geometry and Modeling G22.3033.007, Spring 2005, Professor Yap

May 4, 2005

Instructions

- Out: May 4, 2005
- Due: May 6, noon.
- You may refer to any text books and web resources. But you may not consult with anyone. Unless the information is generic, you should give references.
- You are welcome to use any software (Core Library, Maple, etc) but please include source programs you use.
- 1. Question 1. Let $A(X,Y) \in K[X,Y]$ where $K \subseteq \mathbb{C}$ is a field. Define A = A(X,Y) to be square-free if A is non-zero and there does not exist $B \subseteq K[X,Y]$ such that $\deg(B) > 0$ and B^2 divides A. This definition seems to depend on the choice of K, not just on A(X,Y). E.g., if A(X,Y) is square-free when viewed as an element of K[X,Y], we may be able to choose larger field K' containing K and show that A(X,Y) is no longer square-free when considered as an element of K'[X,Y]. The following shows that this cannot happen. This result is important for us because the curve A = 0 to said to be reduced if A is square-free.

(a) Let A_x and A_y denote the derivative of A with respect to X and Y, respectively. Prove that if A is not square-free iff $\deg(\operatorname{GCD}(A, A_x)) > 0$, iff $\deg(\operatorname{GCD}(A, A_y)) > 0$.

- (b) Conclude that the notion of square-freeness of A does not depend on the choice of K.
- 2. Question 2. When is a curve irreducible? Ruppert's criterion for (absolute) irreducibility of a bivariate polynomials is given in the appendix below. It is known that the curve $Y^n + X^{n+1} = 0$ is irreducible for all $n \ge 1$. Using Ruppert's criterion, show that this is the case when n = 2. HINT: If you work by hand, you should organize your work carefully, as the matrix can be as large as 24×13 . EXTRA CREDIT if you can prove that this curve is irreducible for all n.
- 3. Question 3. We have emphasized that we mainly want to study curves in real affine space. But consider the following two circles C, C' of radii r > 0 and R > 0:

$$C: X^2 + Y^2 = r^2, \qquad C': X^2 + Y^2 = R^2$$

viewed as curves in *projective complex space*, $\mathbb{P}^3(\mathbb{C})$. Determine as much as you can about the nature (topology, dimension, etc) of their intersection. HINT: first homogenize the above equations.

4. Question 4. In the Delaunay Tetrahedralization of a set $S \subseteq \mathbb{R}^3$ of n points, we usually assume that the set S is **non-degenerate** in the sense that no 4 points are co-planar and no 5 points are co-spherical. This is certainly true of Edelsbrunner's book. Suppose we now allow degenerate point sets.

(a) Define the **generalized Delaunay face** of S to be a maximal set of 3 or co-planar points that lie on the boundary of an empty circle. (Empty means no points of S lie in the interior of this circle.)

Similarly, a generalized Delaunay tetrahedron of S to be a maximal set of 4 or more co-spherical points that lie on the boundary of an empty sphere. Describe in detail a data structure to represent this kind of generalized Delaunay tetrahedralization.

(b) Generalize the flipping algorithm (Edelsbrunner, Chapter 5.4) to deal with this situation. Prove the correctness of your algorithm.

- 5. Question 5. Suppose α is an algebraic number, say, represented by the isolated interval representation. Let B(X) and C(X) be integer polynomials.
 - (i) Describe an algorithm to test if α is a zero of B(X).
 - (ii) Describe an algorithm to test if $C(\alpha)$ is a zero of B(X).

Notes on Irreducibility Test

Recall that a polynomial $A(X,Y) \in K[X,Y]$ $(K \subseteq \mathbb{C})$ is said to be **absolutely irreducibile** if it has no factors in C[X,Y]. We simply say "irreducible" for "absolute irreducible". We have the following irreducibility criteria from Ruppert. First, write $\deg(A) \leq (m,n)$ if $\deg_X(A) \leq m$ and $\deg_Y(A) \leq n$.

THEOREM 1 Let $\deg(A) \leq (m, n)$. A is irreducible iff

$$\frac{\partial}{\partial Y}\frac{g}{f} - \frac{\partial}{\partial X}\frac{h}{f} \tag{1}$$

does not have a solution in $g, h \in \mathbb{C}[X, Y]$ with $\deg(g) \leq (m - 1, n)$ and $\deg(h) \leq (m, n - 2)$.

First note that (1) is equivalent to

$$f\frac{\partial g}{\partial Y} - g\frac{\partial f}{\partial Y} = f\frac{\partial h}{\partial X} - h\frac{\partial f}{\partial X}.$$
(2)

Note that the choice

$$g = \partial f / \partial X, \qquad h = \partial f / \partial Y$$

would satisfy (2) but this deg(h) does not satisfy the bounds on degrees of g and h in this theorem. Writing

$$F := f \frac{\partial g}{\partial Y} - g \frac{\partial f}{\partial Y} - f \frac{\partial h}{\partial X} + h \frac{\partial f}{\partial X}$$

we see that $\deg(F) \leq (2m-1, 2n-1)$. Let $\mathbf{x} = (X^{2m-1}Y^{2n-1}, X^{2m-2}Y^{2n-1}, \dots, X, 1)$ denote a (4mn)-vector containing all power products that could appear F. Write $g = \sum_{i=0}^{m-1} \sum_{j=0}^{n} g_{ij}X^iY^j$ and $h = \sum_{i=0}^{m} \sum_{j=0}^{n-2} h_{ij}X^iY^j$ where g_{ij}, h_{ij} are indeterminates. There are 2mn + n - 1 indeterminates, and we denote by \mathbf{v} the (2mn + n - 1)-vector containing these indeterminates. Note that the F is linear in this set of determinates. Similarly, F is linear in the coefficients of g and h. Hence, we may write the equation F = 0 in the matrix form:

$$\mathbf{x}^T R(f) \mathbf{v} = 0$$

where R(f) is a $(4mn) \times (2mn + n - 1)$ matrix whose entries are determined by the coefficients of f. Thus F = 0 is equivalent to R(f) being singular. Then Ruppert's criterion says f is reducible iff R(f) has rank less than 2mn + n - 1.