

$$\begin{aligned}
\sum_{\rho^{1/5} \leq \tau < 1} c\tau^{-2}\sqrt{\rho}G(\rho) &= O(\rho^{\frac{1}{2}-\frac{2}{5}}\ln^2(\rho^{-1})) \\
\sum_{1 \leq \tau} c\tau^{-2}\sqrt{\rho}G(\tau^{-1}\rho) &\leq c' \sum_{1 \leq \tau} \tau^{-2}\sqrt{\rho}(\ln \tau + \ln(\rho^{-1})) \\
&= O(\sqrt{\rho} \ln(\rho^{-1}))
\end{aligned}$$

which are all  $o(\rho)$  so that (14) holds when  $\rho_0$  is picked sufficiently small.

We bound

$$2n^{-1/4} \sum_{\tau} b(\tau\sqrt{n}) \leq \sum_{\tau < \rho^{1/5}} \tau^{0.4} + \sum_{\rho^{1/5} \leq \tau < 1} \tau^{1/2}\rho + \sum_{1 \leq \tau} \tau^{-1/2}\rho$$

The first sum dominates and this is  $O(\rho^{4/25})$  as  $\rho \rightarrow 0$ . We have shown:

**Lemma 5.3** *There are absolute positive constants  $\rho_0, c$  so that if  $|X| = \rho n$ ,  $\rho < \rho_0$  then there exists a partial coloring  $\chi$  so that*

$$|\chi(A \cap X)| \leq cn^{1/4}\rho^{4/25}$$

for all  $A \in \mathcal{A}$  and with at least half the points of  $X$  colored.

The exponent  $\frac{4}{25}$  clearly could be improved by more careful calculation but it does not matter. We are done. Begin with  $X = [n]$ . Apply Lemma 3.1 and then Lemma 5.1 until  $|X| < \rho_0 n$  and then apply Lemma 5.3 until  $|X| < n^{1/4}$  and then color the remaining points arbitrarily. The final coloring  $\chi$  has

$$|\chi(A)| \leq cn^{1/4} + \sum_{i=0}^{\infty} c'n^{1/4}(\rho_0 2^{-i})^{4/25} + n^{1/4} \leq c^*n^{1/4}$$

for all  $A \in \mathcal{A}$  and has  $no$  points uncolored.

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Let  $f(m, s)$  denote the number of  $s$ -sets in  $\mathcal{C}_X$  so that  $f(m, s) \leq s^{-1}U$  with  $U$  as in §4. We first apply the elementary bound

$$f(m, s) \leq \frac{mn}{s^2}$$

**Lemma 5.2** *There is a partial coloring  $\chi$  of  $X$  with*

$$|\chi(A \cap X)| \leq cn^{1/4}$$

*for all  $A \in \mathcal{A}$  and with more than half the points of  $X$  colored.*

**Proof:** We follow the proof of Lemma 3.1 precisely. For  $s = \tau\sqrt{n}$  we set

$$b(\tau\sqrt{n}) = \sqrt{\tau\sqrt{n}} \sigma(\tau\sqrt{n}) \text{ with } \sigma(\tau\sqrt{n}) = \begin{cases} c'\tau^{-1} & \text{if } \tau \geq 1 \\ c'\tau^{-0.1} & \text{if } \tau < 1 \end{cases}$$

Again we need (8) and the bound (9) is the same.  $\square$

We iterate this result, beginning at  $X = [n]$ , resetting  $X$  to be the uncolored points at each iteration, stopping when  $|X| < \rho_0 n$ , with  $\rho_0$  a sufficiently small (as determined later) absolute constant. This is a constant number of iterations (recall the number of uncolored points is at least halved at each iteration) so together we have a partial coloring  $\chi$  with  $|\chi(A)| \leq cn^{1/4}$  for all  $A \in \mathcal{A}$  and a set  $X$  of fewer than  $\rho_0 n$  points uncolored.

**Remark:** Continuing this process until  $|X| < n^{1/4}$  and then coloring the remaining points arbitrarily would give a full coloring with all  $|\chi(A)| \leq cn^{1/4} \ln n$ . Our “slight” improvement of §4 will allow a slight improvement as  $X$  becomes smaller so that the sum converges to  $O(n^{1/4})$ .

Now fix  $X$  with  $|X| = \rho n$ ,  $n^{-3/4} \leq \rho \leq \rho_0$ . We set  $b(\tau\sqrt{n}) = \sqrt{\tau\sqrt{n}} \sigma(\tau\sqrt{n})$  with

$$\sigma(\tau\sqrt{n}) = \begin{cases} \tau^{-0.1} & \text{if } \tau < \rho^{1/5} \\ \rho & \text{if } \rho^{1/5} \leq \tau < 1 \\ \tau^{-1}\rho & \text{if } 1 \leq \tau \end{cases} \quad (13)$$

and set

$$f(\tau\sqrt{n}) = \begin{cases} m\tau^{-2} & \text{if } \tau < \rho^{1/5} \\ cm\tau^{-2}\sqrt{\rho} & \text{if } \tau \geq \rho^{1/5} \end{cases}$$

which, by a slight weakening of Proposition 4.1, is an upper bound on the number of  $\tau\sqrt{n}$ -sets in  $\mathcal{C}_X$ .

We first claim that for  $\rho$  appropriately small

$$m^{-1} \sum_{\tau} f(\tau\sqrt{n}) G(\sigma(\tau\sqrt{n})) \leq 0.2 \quad (14)$$

(we recall the convention from §3 —  $\tau$  in summation runs thru integral powers of 2). We split the sum by the ranges of (13). As functions of  $\rho$

$$\sum_{\tau < \rho^{1/5}} \tau^{-2} G(\tau^{-0.1}) = O(\rho^{-2/5} e^{-\rho^{-1/30}})$$

The upper bound (12) suggests to define the  $d_i$ 's by the initial condition  $d_1 = \varepsilon n/s$  and by the recurrence

$$d_{i+1} = d_i + \left\lfloor \frac{s^2 d_i^2}{9mn} \right\rfloor.$$

One may check that with our choice of parameters,  $s^2 d_1^2 / (9mn) \geq 2$ , and therefore  $d_{i+1} \geq d_i + s^2 d_i^2 / (18mn)$ . We need to check the validity of (11), but this follows by calculation from the assumption  $s \geq 5\sqrt{m}$ .

It remains to estimate the smallest  $k$  such that  $d_{k+1} \geq n/s$ . Set  $\alpha = s^2 / (18mn)$ . Then  $d_{i+1} \geq d_i(1 + \alpha d_i)$ . Given  $i$ , let us estimate the number  $j$  of steps needed so that  $d_{i+j} \geq 2d_i$ . We have  $d_{i+j} \geq d_i(1 + \alpha d_i)^j \geq d_i(j\alpha d_i)$ , so  $j \geq 1/(2\alpha d_i)$  suffices for the doubling. Therefore, the first doubling (from  $d_1$  to at least  $2d_1$ ) needs

$$\frac{1}{2d_1\alpha} = \frac{9m}{s\varepsilon} = O\left(\frac{n}{s}\sqrt{\rho}\right)$$

steps. Then the successive doubling times decrease geometrically, until the ratio of successive two members of the sequence of the  $d_i$ 's exceeds 2. The number of remaining steps needed for reaching  $n/s$  after this happens is at most  $\log_2((n/s)/d_1) = \log_2(1/\varepsilon)$ . Therefore  $k = O((n/s)\sqrt{\rho} + \log(1/\rho)) = O((n/s)\sqrt{\rho})$ , and  $\sum_{d \in J} |U(d)| = O((mn/s)\sqrt{\rho})$  as claimed.  $\square$

**Remark:** The set  $S = \{1, \dots, m\}$  gives a value  $U \sim \frac{m^2}{s} = \frac{nm}{s}\rho$ . Finding the maximal value of  $U = U(n, m, s)$  is an intriguing problem we do not pursue here but we conjecture that our Proposition 4.1 is *not* best possible.

## 5 The End of the Hunt

Let  $X \subseteq [n]$ ,  $|X| = m = \rho n$  with  $n^{-3/4} \leq \rho \leq 1$ . Our object is to find a partial coloring  $\chi$  of  $X$  so that  $|\chi(X \cap A)|$  is small for all  $A \in \mathcal{A}$  and at least half the points of  $X$  are colored. Once successful, we'll apply this process iteratively beginning with  $X = [n]$  (which we did in §3), resetting  $X$  to be the uncolored points, until  $|X| < n^{1/4}$  at which time the remaining points may be colored arbitrarily.

Following (6) set

$$\mathcal{C} = \mathcal{C}_X = \bigcup_{1 \leq d \leq n} \bigcup_{0 \leq i < d} \text{CINT}\{\{x \in X : x \equiv i \pmod{d}\}\}$$

For any  $A \in \mathcal{A}$  we may, as in §2.1, decompose  $A \cap X = B \setminus C$  with  $C \subset B$  and  $B, C$  both disjoint unions of sets of  $\mathcal{C}_X$  of different cardinalities. Lemma 3.3 now generalizes.

**Lemma 5.1** *If  $\chi$  is a partial coloring of  $X$  so that  $|\chi(Y)| \leq b(|Y|)$  for all  $Y \in \mathcal{C}_X$  then*

$$|\chi(A \cap X)| \leq 2 \sum_{s=2^i} b(s)$$

for all  $A \in \mathcal{A}$ .

**Proof.** A number  $x \in U(d) \cap U(d')$  can be specified by giving the number  $r = \lfloor x/\text{lcm}(dd') \rfloor$  plus the residue classes of  $x$  modulo  $d$  and modulo  $d'$ , by the Chinese Remainder Theorem. The number  $r$  can be chosen in at most  $\lfloor n/\text{lcm}(d, d') \rfloor$  ways, and we note that  $U(d)$  may intersect at most  $|U(d)|/s$  residue classes modulo  $d$ , and similarly for  $U(d')$ .  $\square$

**Lemma 4.3** *Let  $I \subseteq \{1, 2, \dots, n\}$  be a set such that  $d \geq d_0$  for all  $d \in I$ , and  $\gcd(d, d') \leq M$  for all distinct  $d, d' \in I$ . Suppose that  $d_0^2/n \leq M \leq s^2 d_0^2/(9mn)$ . Then*

$$\sum_{d \in I} |U(d)| \leq 2m.$$

**Proof.** If not add indices to  $I$  one by one until the sum first gets over  $2m$ . Stopping then would give a set  $I$  with the same assumptions where  $x = \sum_{d \in I} |U(d)|$  satisfies  $2m < x \leq 3m$ . We use Inclusion-Exclusion:

$$m \geq \left| \bigcup_{d \in I} U(d) \right| \geq \sum_{d \in I} |U(d)| - \sum_{d, d' \in I, d < d'} |U(d) \cap U(d')|. \quad (10)$$

By Lemma 4.2 and by the assumptions on  $I$ , we have

$$\sum_{d < d'} |U(d) \cap U(d')| \leq \sum_{d < d'} \frac{|U(d)| \cdot |U(d')|}{s^2} \left\lceil \frac{nM}{d_0^2} \right\rceil \leq \frac{1}{2s^2} \left( \sum_{d \in I} |U(d)| \right)^2 \left( \frac{nM}{d_0^2} + 1 \right)$$

The assumption on  $M$  implies  $nM/d_0^2 \geq 1$ . Thus, from (10), we further get

$$m \geq x - \frac{x^2}{2s^2} \frac{2nM}{d_0^2} > 2m - \frac{9m^2 nM}{s^2 d_0^2} \geq 2m - m = m$$

(using the upper bound on  $M$  in the assumption of the Lemma), a contradiction.  $\square$

**Proof of Proposition 4.1.** We may suppose that  $m, n, s, \rho^{-1}$  are all sufficiently large (otherwise the claim is satisfied trivially). We fix a parameter  $\varepsilon = 5\sqrt{\rho}$ . We let  $J$  be the interval

$$J = \left[ \varepsilon \frac{n}{s}, \frac{n}{s} \right]$$

(we may also suppose that  $\varepsilon mn/s$  is an integer). We note that the  $d$  lying outside the interval  $J$  only contribute at most  $\varepsilon mn/s = 5(nm/s)\sqrt{\rho}$  to  $U$ . Hence it suffices to bound  $\sum_{d \in J} |U(d)|$ .

We want to partition the interval  $J$  into consecutive intervals  $I_1, I_2, \dots, I_k$ , in such a way that Lemma 4.3 can be applied to each of them, giving the bound  $\sum_{d \in I_i} |U(d)| \leq 2m$ . It remains to calculate how small can  $k$  be made. If we denote  $I_i = [d_i, d_{i+1})$ , then we have  $\gcd(d, d') \leq d_{i+1} - d_i$  for any 2 distinct numbers  $d, d' \in I_i$ . Thus, in order to apply Lemma 4.3, it is enough to have

$$d_{i+1} - d_i \geq \frac{d_i^2}{n} \quad (11)$$

$$d_{i+1} - d_i \leq \frac{s^2 d_i^2}{9mn}. \quad (12)$$

$O(\tau^{-2} \exp(-\tau^{-0.2}/9))$  ( $\tau$  small) give convergent sums so we find an absolute constant  $T$  for which

$$\sum_{\tau \geq T} \tau^{-2} G(\tau^{-1}) + \sum_{\tau < T^{-1}} \tau^{-2} G(\tau^{-0.1}) < 0.1$$

As  $\lim_{x \rightarrow \infty} G(x) = 0$  we may now select  $c' \geq 1$  sufficiently large so that the finite sum

$$\sum_{T^{-1} < \tau < T} \tau^{-2} G(c\tau^{-0.1}) < 0.1$$

yielding (8). Hence by Lemma 3.3 and Corollary 2.4 there is a partial coloring of  $[n]$  with at least half of the points colored and with

$$|\chi(A)| \leq 2 \sum_{\tau} b(\tau \sqrt{n}) \leq 2c'n^{1/4} \left[ \sum_{\tau \geq 1} \tau^{-1/2} + \sum_{\tau < 1} \tau^{0.4} \right] \quad (9)$$

for all  $A \in \mathcal{A}$ . As the bracketed sums both converge this gives Lemma 3.1.

## 4 Number Theory

Let  $X \subseteq \{1, \dots, n\}$  be an  $m$ -element set. Let  $s$  be an integer,  $1 \leq s \leq n$ . For an integer  $d$ , let  $U(d)$  denote the set of all  $x \in X$  in residue classes modulo  $d$  for which at least  $s$  elements of  $X$  lie in that residue class. We are interested in the quantity

$$U = \sum_d |U(d)|$$

We can clearly restrict ourselves to the range  $1 \leq d \leq n/s$  (for larger  $d$ ,  $U(d) = \emptyset$ ). Also, for each  $d$ ,  $|U(d)| \leq m$ , and thus we get  $U \leq nm/s$ . This is tight for  $s = 1$  but, for large enough  $s$ , the following theorem gives an improvement. The intuition behind it is that while for some individual value of  $d$ , the members of  $X$  can be distributed among very few residue classes modulo  $d$  only, such a distribution cannot occur for too many values of  $d$  at once.

Set  $\rho = m/n$ . We have

**Proposition 4.1** *Suppose that  $5\sqrt{m} \leq s \leq m$ , then*

$$U \leq c \frac{nm}{s} \sqrt{\rho},$$

for an absolute constant  $c$ .

**Lemma 4.2** *For any pair  $d, d'$  of distinct natural numbers, we have*

$$|U(d) \cap U(d')| \leq \frac{|U(d)| \cdot |U(d')|}{s^2} \left\lceil \frac{n}{\text{lcm}(d, d')} \right\rceil.$$

**Proof:** With  $A = \{x_u : i \leq u \leq j\}$  set  $B = \{x_u : 1 \leq u \leq j\}$  and  $C = \{x_u : 1 \leq u \leq i-1\}$ . Take the binary expansion  $j = 2^{b_1} + 2^{b_2} + \dots$ ,  $b_1 > b_2 > \dots$  of  $j$ . Decompose  $B$  into the first  $2^{b_1}$  elements of  $X$  union the next  $2^{b_2}$  elements of  $X$  ..., and do likewise with  $C$ .  $\square$

We can think of any arithmetic progression as a subinterval of an entire residue class so that

$$\mathcal{A} = \bigcup_{1 \leq d \leq n} \bigcup_{0 \leq i < d} \text{INT}[\{x \in [n] : x \equiv i \pmod{d}\}]$$

We define the ‘‘canonical arithmetic progressions’’

$$\mathcal{C} = \mathcal{C}_n = \bigcup_{1 \leq d \leq n} \bigcup_{0 \leq i < d} \text{CINT}[\{x \in [n] : x \equiv i \pmod{d}\}] \quad (6)$$

**Lemma 3.3** *If  $\chi$  is a partial coloring of  $[n]$  so that*

$$\chi(X) \leq b(|X|)$$

*for all  $X \in \mathcal{C}$  then*

$$\chi(A) \leq 2 \sum_{s; s=2^i \leq n} b(s) \quad (7)$$

*for all  $X \in \mathcal{A}$ .  $\square$*

## 3.2 The Coloring

For  $s = 2^i \leq n$  how many  $s$ -sets are in  $\mathcal{C}_n$ ? We restrict  $1 \leq d \leq \frac{n-1}{s-1}$  (otherwise the residue classes have fewer than  $s$  elements) and for each  $d$  the  $s$ -sets are disjoint so there are at most  $\frac{n}{s}$  of them, giving an upper bound of  $\frac{n(n-1)}{s(s-1)}$  of them. For  $s = 1$  there are only  $n$  distinct singletons. Ignoring asymptotically insignificant terms we’ll say that  $\mathcal{C}_n$  has at most  $n^2 s^{-2}$  sets of size  $s$ .

**Remark:** For  $s \sim \sqrt{n}$  we have  $\sim n$  sets of size  $s$  and Corollary 2.4 gives a partial coloring with  $|\chi(A)| \leq cn^{1/4}$  for all such sets. We need simultaneously color the larger and smaller sets. To avoid a logarithmic term in applying (7) we’ll need a slightly better bound on  $|\chi(A)|$  when  $|A|$  is not near  $\sqrt{n}$ .

We parameterize  $s = \tau\sqrt{n}$  so that we have  $n\tau^{-2}$  sets of size  $s$ . We’ll assume for convenience that  $\sqrt{n}$  is a power of two so that  $\tau = 2^i$ ,  $i$  integral. We set

$$b(\tau\sqrt{n}) = \sqrt{\tau\sqrt{n}} \sigma(\tau\sqrt{n}) \text{ where } \sigma(\tau(\sqrt{n})) = \begin{cases} c'\tau^{-1} & \text{if } \tau \geq 1 \\ c'\tau^{-0.1} & \text{if } \tau < 1 \end{cases}$$

We claim that, for an appropriately large constant  $c'$ , (5) is now satisfied. We need show

$$\sum_{\tau \geq 1} \tau^{-2} G(c'\tau^{-1}) + \sum_{\tau < 1} \tau^{-2} G(c'\tau^{-0.1}) < \frac{1}{5} \quad (8)$$

where  $\tau$  in the sums runs over integral powers of 2 and  $G$  is given by (4). We will insist  $c' \geq 1$  so that  $G(c'y) \leq G(y)$ . Both  $\tau^{-2} G(\tau^{-1}) = O(\tau^{-2} \ln(\tau))$  ( $\tau$  large) and  $\tau^{-2} G(\tau^{-0.1}) =$

**Corollary 2.4** *Let  $\mathcal{A}$  be a family of subsets of an  $n$ -set  $\Omega$  consisting of at most  $f(s)$  sets of size  $s$ . If  $b(s) = \sigma(s)\sqrt{s}$  where  $\sigma(s)$  satisfies*

$$\sum_s f(s)G(\sigma(s)) \leq \frac{n}{5} \tag{5}$$

*then there is a partial coloring  $\chi$  of  $\Omega$  with  $|\chi(A)| \leq b(|A|)$  for all  $A \in \mathcal{A}$  and at least half the points of  $\Omega$  colored.  $\square$*

With these bounds we can already give a result which is interesting in its own right and may give significant insight into the somewhat technical computations to come.

**Theorem 2.5** *There is an absolute constant  $c$  so that the following holds for all  $n, s$ . If  $A_1, \dots, A_n \subseteq \Omega$  and  $|\Omega| = n$  and all  $|A_i| \leq s$  then there is a partial coloring  $\chi$  of  $\Omega$  with less than half the points of  $\Omega$  uncolored and with*

$$|\chi(A_i)| \leq c\sqrt{s}$$

*for all  $1 \leq i \leq n$ .*

**Proof.** From Lemma 2.3 we may pick  $c$  so that  $G(c) \leq 0.2$ . Now apply Corollary 2.4.  $\square$

The monotonicity of  $G$  allows a further generalization of Corollary 2.4. Suppose  $\mathcal{A}$  is a family of subsets of an  $n$ -set  $\Omega$  which breaks into subfamilies consisting of at most  $f(s)$  sets of size *at most*  $s$ . When (5) holds the conclusion of Corollary 2.4 then holds. In particular, given any  $A_1, \dots, A_n \subseteq \{1, \dots, n\}$  we have  $n$  sets of size at most  $n$  we may pick  $c$  so that  $G(c) < 0.2$  and then there exists a partial coloring  $\chi$  of  $\{1, \dots, n\}$  with all  $|\chi(A_i)| \leq c\sqrt{n}$  and at least half the points colored. The result was the core of [5].

### 3 The First Partial Coloring

Let  $\mathcal{A}$  denote the family of arithmetic progressions contained in  $\Omega = \{1, \dots, n\}$ . Here we show:

**Lemma 3.1** *There is a partial coloring of  $\Omega$  so that  $|\chi(A)| \leq cn^{1/4}$  for all  $A \in \mathcal{A}$  and at least half the points of  $\Omega$  are colored.*

#### 3.1 The Decomposition

Let  $X = \{x_1, \dots, x_l\}$  be any set of integers with  $x_1 < \dots < x_l$ . Define  $INT(X)$  to be the family of intervals — i.e., all sets  $\{x_u : i \leq u \leq j\}$  where  $1 \leq i \leq j \leq l$ . Now define  $CINT(X)$  (the *canonical intervals* on  $X$ ) by taking, for all powers of two  $s = 2^i \leq l$ , all sets  $\{x_{(j-1)s+1}, \dots, x_{js}\}$  with  $js \leq l$ . That is, we split  $X$  into consecutive intervals of length  $s = 2^i$ , ignoring the “extra”. The following observation is standard:

**Lemma 3.2** (Decomposition lemma) *Any  $A \in INT(X)$  can be written  $A = B \setminus C$  with  $C \subset B$  and with  $B$  and  $C$  both decomposable into disjoint unions of sets in  $CINT(X)$  of different sizes.*

**Corollary 2.2** *Let  $\mathcal{A}$  be a family of subsets of an  $n$ -set  $\Omega$  consisting of at most  $f(s)$  sets of size  $s$ . If  $b(s)$  satisfies*

$$\sum_s f(s)ENT(s, b(s)) \leq \frac{n}{5}$$

*then there is a partial coloring  $\chi$  with  $|\chi(S)| \leq b(|S|)$  for all  $S \in \mathcal{A}$  and fewer than half the points of  $\Omega$  uncolored.*

In applying Corollary 2.2 we need upper bounds on  $ENT(n, b)$ . The correct parameterization is  $b = \lambda\sqrt{n}$ . Roughly  $S_n$  is like  $\sqrt{n}N$  where  $N$  is standard normal so that  $ENT(n, \lambda\sqrt{n})$  should be like  $g(\lambda) = H(R_\lambda(N))$ . Analysis gives that for  $\lambda$  large  $g(\lambda) = \Theta(\lambda^2 e^{-\lambda^2/2})$  ( $i = \pm 1$  giving the dominant terms) while for as  $\lambda \rightarrow 0$   $g(\lambda) = \Theta(\ln(\lambda^{-1}))$ , the major contribution being  $p_i = \Theta(\lambda^{-1})$  for  $i = O(\lambda^{-1})$ . The following results are somewhat weaker and certainly not best possible but have the advantage of holding for *all*  $n, \lambda$ .

**Lemma 2.3** *There is an absolute constant  $c$  so that  $ENT(n, \lambda\sqrt{n}) \leq G(\lambda)$  where we define*

$$G(\lambda) = \begin{cases} ce^{-\lambda^2/9} & \text{if } \lambda \geq 10 \\ c & \text{if } 0.1 \leq \lambda \leq 10 \\ c \ln(\lambda^{-1}) & \text{if } \lambda < 0.1 \end{cases} \quad (4)$$

**Proof (Outline):** We employ the universal bound

$$\Pr[S_n \geq \tau\sqrt{n}] \leq e^{-\tau^2/2}$$

Set  $g_i = \exp(-\lambda^2(2i-1)^2/8)$ ,  $i \geq 1$  and  $g_0 = 1 - 2 \exp(-\lambda^2/8)$ . From (2)  $p_i, p_{-i} \leq g_i$  and  $p_0 \geq g_0$ . On  $[0, 1]$  the function  $-x \log_2 x$  increases to  $x = e^{-1}$  and then decreases. When  $\lambda \geq 10$ ,  $g_i \leq e^{-1}$  for all  $i \geq 1$  and  $g_0 \geq e^{-1}$  so

$$ENT(n, \lambda\sqrt{n}) \leq -g_0 \log_2 g_0 + 2 \sum_{i=1}^{\infty} -g_i \log_2 g_i$$

This is a continuous function of  $\lambda$  which is  $O(\lambda^2 e^{-\lambda^2/8})$  or, giving ground,  $O(e^{-\lambda^2/9})$ . When  $0.1 \leq \lambda \leq 10$  set  $I = \{-100, \dots, +100\}$ . The contribution to  $ENT(n, b)$  from  $i \in I$  is at most  $\log_2 |I| \leq 8$ . For  $i \notin I$  certainly  $g_i < e^{-1}$  so

$$ENT(n, \lambda\sqrt{n}) \leq 8 + 2 \sum_{i=101}^{\infty} -g_i \log_2 g_i \leq 9$$

For  $\lambda < 0.1$  set  $I = \{i : |i| < \lambda^{-20}\}$ . Again for  $i \notin I$  we have  $g_i \leq e^{-1}$  and

$$ENT(n, \lambda\sqrt{n}) \leq \log_2(2\lambda^{-20} + 1) + 2 \sum_{|i| > \lambda^{-20}} -g_i \log_2 g_i \leq 40 \ln(\lambda^{-1})$$

by computation.  $\square$

We may now further reexpress Corollary 2.2.



With this definition we give our general criterion.

**Lemma 2.1** *Let  $S_1, \dots, S_v \subseteq \Omega$  with  $|\Omega| = n$  and  $|S_i| = n_i$ . Suppose  $b_i, \varepsilon$  and  $\gamma \leq \frac{1}{2}$  are such that*

$$\sum_{i=1}^v ENT(n_i, b_i) \leq \varepsilon n$$

and

$$\sum_{j=0}^{\gamma n} \binom{n}{j} < 2^{n(1-\varepsilon)} \quad (3)$$

Then there is a partial coloring  $\chi$  of  $\Omega$  with

$$|\chi(S_i)| \leq b_i \text{ for all } i$$

and more than  $2\gamma n$  points  $x \in \Omega$  colored.

**Proof.** Consider the uniform probability space of all  $\chi : \Omega \rightarrow \{-1, +1\}$  and define the random variable

$$L(\chi) = (R_{b_1}(\chi(S_1)), \dots, R_{b_v}(\chi(S_v)))$$

By subadditivity of entropy

$$H(L) \leq \sum_{i=1}^v H(R_{b_i}(\chi(S_i))) = \sum_{i=1}^v ENT(n_i, b_i) \leq \varepsilon n$$

Hence some value of  $L$  has probability at least  $2^{-\varepsilon n}$  of being achieved. As all  $\chi$  have probability  $2^{-n}$  there is a set  $\mathcal{C}$  of at least  $2^{n(1-\varepsilon)}$  colorings  $\chi$  so that if  $\chi_1, \chi_2 \in \mathcal{C}$ , then  $L(\chi_1) = L(\chi_2)$ .

We naturally associate such colorings  $\chi$  with points on the Hamming Cube  $\{-1, +1\}^n$ . (With  $\Omega = \{1, \dots, n\}$  associate  $\chi$  with  $(\chi(1), \dots, \chi(n))$ .) A theorem of Kleitman [4] (basically an isoperimetric inequality) states that any  $\mathcal{C} \subseteq \{-1, +1\}^n$  of size bigger than  $\sum_{j=0}^l \binom{n}{j}$  with  $l \leq \frac{n}{2}$  contains two points at Hamming distance (i.e., the number of different coordinates) at least  $2l$ . (This is ‘‘best possible’’ as  $\mathcal{C}$  may be the set of all sequences with at most  $l$  coordinates  $+1$ .) Thus there are  $\chi_1, \chi_2 \in \mathcal{C}$  at Hamming distance at most  $2\gamma n$ . Set

$$\chi(x) = \frac{\chi_1(x) - \chi_2(x)}{2} \text{ for all } x \in \Omega.$$

Then  $\chi$  is a partial coloring. The number of colored points is precisely the Hamming distance which is at least  $2\gamma n$ . For each  $i$  the values  $\chi_1(S_i), \chi_2(S_i)$  have the same  $b_i$ -roundoff and therefore lie in a common open interval of length less than  $2b_i$ . Thus

$$|\chi(S_i)| = \left| \frac{\chi_1(S_i) - \chi_2(S_i)}{2} \right| < b_i$$

as desired.  $\square$

We note that (3) holds for, say,  $\gamma = \frac{1}{4}$  and  $\varepsilon = 0.2$ , this value will suffice for our purposes. Also, we shall always use a bound on  $|\chi(S)|$  dependent only on  $|S|$ . We'll use the Lemma in the following simpler form.

In words, we show the *existence* of a two-coloring  $\chi$  of the first  $n$  integers so that all arithmetic progressions  $A$  have imbalance  $|\chi(A)| \leq Cn^{1/4}$ . We remark that the proof does not give a construction of  $\chi$  in the usual sense and is indeed not satisfactory from an algorithmic point of view. The methods of §2 (see comments in [5]) are such that we have not been able to obtain an algorithm that would output this coloring  $\chi$  in time polynomial in  $n$ . Our proof involves variants of the probabilistic method, we give [1] as a general reference. The technique of our proof combines methods of [2], [5], [6].

Throughout the paper, we'll use the symbols  $c, c'$  etc. generically for denoting absolute constants, and in order to limit the number of symbols, we reuse them freely.

## 2 Entropy

Let  $A_1, \dots, A_v \subseteq \Omega$ . A *partial coloring* is a map  $\chi : \Omega \rightarrow \{-1, 0, +1\}$ . When  $\chi(x) = 0$  we call  $x$  uncolored, otherwise  $x$  is called colored. We define, for  $A \subseteq \Omega$ ,  $\chi(A) = \sum_{x \in A} \chi(x)$ . Our object will be to give a general condition under which there exists a partial coloring  $\chi$  with the  $|\chi(A_i)|$  “small” and “few”  $x \in \Omega$  uncolored.

For any positive integer  $b$  define the *b-roundoff function*  $R_b(x)$  as that  $i$  so that  $2bi$  is the nearest multiple of  $2b$  to  $x$ . In case of ties take the larger. Thus

$$R_b(x) = 0 \quad \text{if and only if} \quad -b \leq x < b \tag{2}$$

$$R_b(x) \geq i \quad \text{if and only if} \quad x \geq (2i - 1)b$$

$$R_b(x) \leq -i \quad \text{if and only if} \quad x < -(2i - 1)b$$

Let  $X$  be any discretely valued random variable. We use the standard definition of the *entropy function*  $H(X)$

$$H(X) = \sum_i -p_i \log_2(p_i)$$

where  $p_i = \Pr[X = i]$ , the summation is over the possible values of  $X$ , and  $0 \log_2 0$  is interpreted as 0. We shall use the following well known facts about entropy:

- Entropy is subadditive. That is, if  $X = (X_1, \dots, X_v)$  then  $H(X) \leq \sum_{i=1}^v H(X_i)$ .
- When  $X$  takes on at most  $K$  values it has entropy at most  $\log_2 K$ , the extreme case being a uniformly chosen value from a  $K$ -set. Moreover  $\sum_{i \in I} -p_i \log_2(p_i) \leq \log_2 |I|$  for any subset of values of  $X$ .
- When  $X$  has entropy less than  $K$  it takes on some value with probability at least  $2^{-K}$ .

Let  $S_n$ , as standard, denote the sum of  $n$  independent random variables, each uniform on  $\{-1, +1\}$ . When  $\chi : \Omega \rightarrow \{-1, +1\}$  is uniform and  $A \subseteq \Omega$ ,  $|A| = n$ , then  $\chi(A)$  has distribution  $S_n$ . Now we come to a key definition:

$$ENT(n, b) = H(R_b(S_n))$$

# Discrepancy in Arithmetic Progressions

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## Abstract

It is proven that there is a two-coloring of the first  $n$  integers for which all arithmetic progressions have discrepancy less than  $const.n^{1/4}$ . This shows that a 1964 result of K.F. Roth is, up to constants, best possible.

## 1 Results and History

Let  $\mathcal{A}$  be a family of subsets of a finite set  $\Omega$ . By a two-coloring of  $\Omega$  we shall mean a map  $\chi : \Omega \rightarrow \{-1, +1\}$ . For any  $X \subseteq \Omega$  we define  $\chi(X) = \sum_{x \in X} \chi(x)$ . The *discrepancy* of  $\mathcal{A}$  is defined by

$$\text{disc}(\mathcal{A}) = \min_{\chi} \max_{A \in \mathcal{A}} |\chi(A)| \quad (1)$$

Let  $\Omega = \{1, \dots, n\}$ , which we denote by  $[n]$ . Let  $\mathcal{A}$  denote the set of arithmetic progressions on  $[n]$ . The discrepancy of this set system was investigated in 1964 by K.F. Roth [7]. If we define the function  $ROTH(n) = \text{disc}(\mathcal{A})$ , his result can be written

$$ROTH(n) \geq cn^{1/4}$$

$c$  a positive absolute constant. That is, for any two coloring  $\chi$  of the first  $n$  integers there will be an arithmetic progression  $A$  on which the “imbalance”  $|\chi(A)|$  is at least  $cn^{1/4}$ .

It is interesting that Roth himself did not believe his result to be best possible and speculated that perhaps  $ROTH(n) = n^{1/2-o(1)}$ . Indeed a bound  $ROTH(n) = O(\sqrt{n \ln n})$  follows by elementary probabilistic considerations. In the early 1970's Sárközi (see [3]) showed  $ROTH(n) \leq n^{1/3+o(1)}$ . A breakthrough was given in 1981 by Beck [2] who showed  $ROTH(n) \leq cn^{1/4} \ln^{5/2} n$ . Here we show

### Theorem 1.1

$$ROTH(n) \leq Cn^{1/4}$$

with  $C$  an absolute constant.

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