Virtual Machines: Concepts & Applications

Lecture 1: So ... What Is A Virtual Machine?

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Who Am I?

• Mohamed Zahran (aka Z)
• Computer architecture/OS/Compilers Interaction
• http://www.mzahran.com
• Office hours: Tue 2:00-4:00 pm
• Room: WWH 320
Formal Goals of The Course

• Understand VM architectures and applications
• Study key implementation technologies
• Focus on architecture and microarchitecture aspects; as well as software aspects
• Cover significant case studies
My Wish List for This Course

• Be an expert in how programming languages, compilers, OS, and computer architecture interact together!
• Be able to use the technology learned in this course in many different situations
• Build a vision about technology and its future
• Enjoy the course!
The Textbook
Grades

• You have 3 sources of study information:
  – Slides
  – Notes you take in class
  – Reading material from the textbook

• Exam is open book/notes

• Grade distribution:
  – Homework assignments 30%
  – Project 30%
  – Final exam 40%
We Have to Admit that ... 

• Computers are very complicated structures!
• In order to manage/design them, we need to be able to manage extreme complexity!
• The best way to do that is through: levels of abstraction with well defined interfaces
Abstraction

- Computer systems are built on levels of abstraction
  - Higher level of abstraction hide details at lower levels
  - Example: files are an abstraction of a disk
Abstraction provides:
Simplified interface to underlying resources.
Problem → Algorithm Development → Programmer

High Level Language

Assembly Language

Machine Language

Control Unit (Interpreter)

Microarchitecture

Microsequencer (Interpreter)

Logic Level

Device Level → Semiconductors → Quantum
Advantages of Well-defined Interfaces

• Major design tasks are decoupled
• Different hardware and software development schedules
• Example of interfaces:
  – Instruction set architecture (ISA)
  – OS interface (system calls)
• Software can run on any machine supporting a compatible interface
Major Program Interfaces

- ISA Interface -- supports all conventional software

Application Binary Interface (ABI) -- supports application software only
There are also disadvantages...

- Software compiled for one ISA will not run on hardware with a different ISA
  - Apple Mac (PowerPC) binaries on an x86?
- Even if ISAs are the same, OSes may differ
  - Windows 10 applications on a Linux x86?
- Binary may not be optimized for the specific hardware platform it runs on
  - Intel Pentium 4 binaries on an AMD Athlon?
Disadvantages (contd.)

- Innovation may be inhibited by fixed ISA
  - Hard to add new instructions
    - OR remove obsolete ones
  - What was the most recent (successful) new ISA?
    Or new OS?

- Difficult for software to interact directly with implementation
  - Performance features
  - Power management
  - Fault tolerance
  - Software is *supposed* to be implementation independent
Diversity in instruction sets, OSes, and programming languages

- Encourages innovations
- Discourages stagnation

BUT

In practice, diversity leads to reduced interoperability.

How to deal with this in our world of networked computers, where it is advantageous to move software as freely as data?
A Look at Hardware Resources

- Conventional system software manages hardware resources directly
  - An OS manages the physical memory of a specific size
  - I/O devices are managed as physical entities
- Difficult to share resources except through OS
  - All users of hardware must use the same OS
  - All users are vulnerable to attack from other users sharing the resource (via security holes in OS)

Can we do better?
Virtualization is the answer!

- Real system is transformed so that it appears to be different!
- An isomorphism from guest to host
  - Map guest state to host state
  - Implement “equivalent” functions

\[
\text{Guest} \quad \begin{array}{c}
\text{Guest} \\
v(S_i) \\
S_i \quad e(S_i) \\
S_i \quad v(S_i) \\
\text{Host} \\
e'(S_i') \\
S_i' \quad e'(S_i') \\
S_j' \quad v(S_j) \\
S_j \quad \end{array} \quad \text{Host}
\]
Virtualization

• Similar to abstraction
 Except
  – Details not necessarily hidden

• Construct Virtual Disks
  – As files on a larger disk
  – Map state
  – Implement functions
The “Machine”

- Different perspectives on what the Machine is:
  - OS developer
  - Compiler developer
  - Application programmer
Virtual Machines

add *Virtualizing Software* to a *Host* platform and support *Guest* process or system on a *Virtual Machine* (VM)

Example: System Virtual Machine

![Diagram of virtual machine setup]

**Guest**

**VMM**

**Host**

VMM: Virtual Machine Manager
Example of VM Usages

- A virtualizing software installed on an Apple Macintosh can provide a Windows/x86_64 VM capable of running PC application programs.
- Multiple, replicated VMs can be implemented on a single hardware platform to provide groups/individuals with their own OS environments.
Example of VM Usages

• A large multiprocessor server can be divided into smaller virtual servers.
• VM can provide dynamic, on-the-fly optimization of program binaries.
• ...

The Family of Virtual Machines

• Lots of things are called “virtual machines”
  IBM VM/370
  Java
  VMWare

  Some things *not* called “virtual machines”, *are* virtual machines
  IA-32 EL
  Dynamo
  Transmeta Crusoe
The process of virtualization involves two steps:
1. Mapping of virtual resources or state to real resources of the underlying machine.
2. Using real machine instructions and/or system calls to carry out the actions specified by the VM instructions.
Process VMs

- Execute application binaries with an ISA different from hardware platform
- Provide user application with a virtual ABI environment
- Can provide: replication, emulation, and optimization
- Examples: IA-32 EL, FX!32
Process Virtual Machines

- Constructed at ABI level
- Runtime manages guest process
- Not persistent
- Guest processes may intermingle with host processes
- As a practical matter, guest and host OSes are often the same
- Dynamic optimizers are a special case
- Examples: IA-32 EL, FX!32, Dynamo
Example of Process VM: FX!32

Application compiled for source ISA (in this example IA-32)

Executed on a machine with target ISA (in this example DEC Alpha)
Process VM: Replication

• OS providing the illusion of multiprogramming
• Each user process thinks it has the complete machine to itself
Process VM: Emulation

• Supports program binaries compiled for a different instruction set than the host hardware → emulates one instruction set on hardware designed for another instruction set

• Interpretation Vs Translation
Process VM: Optimization
Same-ISA Dynamic Binary Optimizers

• Optimize Binary at Runtime
• Instruction sets for host and guest are the same
• Example HP Dynamo
  – Can optimize for dynamic properties of program
  – Can optimize for a specific processor implementation
Process VM:
High Level Language Virtual Machines

- VM environment does not directly correspond to any real platform
- VM environment designed for:
  - ease of portability
  - to match features of HLL used for program development.

\[ \text{HLL Program} \rightarrow \text{Intermediate Code} \rightarrow \text{Object Code (ISA)} \rightarrow \text{Memory Image} \]

\[ \text{HLL Program} \rightarrow \text{Portable Code (Virtual ISA)} \rightarrow \text{Virt. Mem. Image} \rightarrow \text{Host Instructions} \]

- Traditional
- HLL VM
HLL VM: Example JVM

Java Binary Classes

VM implementation
Sparc Workstation

VM implementation
x86 PC

VM implementation
Apple Mac
System Virtual Machines

- Provide a complete system environment
- Constructed at ISA level
- Persistent
- Examples: IBM VM/360, VMware, Transmeta Crusoe

A single host hardware platform can support multiple guest OS environments simultaneously.
System Virtual Machines

• Native VM System
  – VMM privileged mode
  – Guest OS user mode
  – Example: classic IBM VMs

• User-mode Hosted VM
  – VMM runs as user application

• Dual-mode Hosted VM
  – Parts of VMM privileged; parts non-privileged
  – Example VMware
Examples
Examples

Windows apps.

Windows

Mac apps.

Mac OS

PowerPC
Co-Designed VMs

Different objective:

Enable innovative ISA and/or hardware implementation for improved performance and/or power.
Co-Designed VMs

- Perform both translation and optimization
- VM provides interface between standard ISA software and implementation ISA
A computer user has a Java application running on laptop PC

Java application

JVM

Linux x86

The user has Linux installed on Windows PC via VMware.

Linux x86

VMware

Windows x86

The IA-32 hardware is in fact a Transmeta Crusoe implementing VLIW ISA.

Windows x86

Binary Translation

Crusoe VLIW
Taxonomy

Process VMs

- **same ISA**
  - Multiprogrammed Systems
  - HP Dynamo

- **different ISA**
  - IA-32 EL FX132
    - Java VM
    - MS CLI

System VMs

- **same ISA**
  - IBM VM/370
    - VMware

- **different ISA**
  - Virtual PC for Mac
    - Transmeta Crusoe
Conclusions

• Virtualization is a key technology for this interconnected world where the cloud is the main player.
• VMs have been investigated and built by OS developers, language designers, compiler developers, and hardware designers!
• In this course, we will study the underlying concepts and technologies that are common across the whole spectrum.