Computer Systems Organization

Lecture 1: Overview

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Who Am I?

• Mohamed Zahran (aka Z)
• Computer architecture/OS/Compilers Interaction
• http://www.mzahran.com
• Office hours: Tuesdays 4:30-6:30 pm
  – or by appointment
• Room: WWH 320
• Course webpage:
  http://cs.nyu.edu/courses/fall13/CSCI-UA.0201-003/
Formal Goals of This Course

• What happens under the hood when you boot your computer and start running applications?
• How do software and hardware interact?
• How to write a program and execute it on a computer?
• This course is programmer-centric
  – Understanding of underlying system makes a more effective programmer
  – Bring out the hidden hacker in everyone
Informal Goals of This Course

• To get more than an A
• To learn computer systems and enjoy it
• To use what you have learned in MANY different contexts
• To start your research project if you want
• To enjoy the course!
Textbook

COMPUTER SYSTEMS
A Programmer's Perspective

Bryant • O’Hallaron

BO
Course Components

• Lectures
  – Higher level concepts

• Programming labs (5 of them → 60%)
  – The heart of the course
  – 1-3 weeks each
  – Provide in-depth understanding of some aspect of systems

• One final exam (40%)
Course Syllabus

• Basic C
  – L1 (Rabin-Karp Lab)

• Assembly: Representation of program and data
  – L2 (Bomblab) and L3 (Buflab)

• Memory hierarchy and systems hardware
  – L4 (cache lab)

• Exceptional control flow: exceptions, process control, signals

• Virtual Memory: address translation, allocation
  – L5 (Malloclab)

• Concurrency
Policies: Assignments (Labs)

• You must work alone on all assignments
  – Post all questions on the mailing list,
  – You are encouraged to answer others’ questions, but refrain from explicitly giving away solutions.

• Hand-ins
  – Assignments due at 11:59pm on the due date
  – Everybody has 5 grace days for the whole semester
  – Zero score if a lab is handed in more than 3 days late
Facilities

• Lab environments:
  – Use official class VM image
    • Download (free) virtualbox for Windows/Linux/Mac
    • Download VM appliance from course web page
Course Theme: Abstraction Is Good But Don’t Forget Reality

• Most CS courses emphasize abstraction
• This class:
  – Helps you peek ```under-the-hood```
• Useful outcomes
  – Become more effective programmers
    • Debug problems
    • Tune performance
  – Prepare for later “systems” courses in CS
    • Compilers, Operating Systems, Computer Architecture, Distributed Systems
Reality #1: Ints are not Integers, Floats are not Reals

- $x^2 \geq 0$? Overflow!!
- $(x + y) + z = x + (y + z)?$ $1e20+(-1e20+3.14)! = 3.14$

Source: xkcd.com/571
Reality #2: You've Got to Know Assembly

- No need to program in assembly
- Knowledge of assembly helps one understand machine-level execution
  - Debugging
  - Performance tuning
  - Writing system software (e.g. compilers, OS)
  - Creating / fighting malware
    - x86 assembly is the language of choice!
Reality #3: Memory Matters

- Memory is not unbounded
  - It must be allocated and managed
- Memory referencing bugs especially wicked
- Memory performance is not uniform
  - Cache and virtual memory effects can greatly affect performance
Memory Referencing Errors

- C/C++ let programmers make memory errors
  - Out of bounds array references
  - Invalid pointer values
  - Double free, use after free
- Errors can lead to nasty bugs
  - Corrupt program objects
  - Effect of bug observed long after the corruption
double fun(int i)
{
    double d[1] = {3.14};
    int a[2];
    a[i] = 1073741824; /* Possibly out of bounds */
    return d[0];
}

fun(0)   =  3.14
fun(1)   =  3.14
fun(2)   =  
fun(3)   =  
fun(4)   =  

Carnegie Mellon

Code Security Example

/* Kernel memory region holding user-accessible data */
#define KSIZE 1024
char kbuf[KSIZE];

/* Copy at most maxlen bytes from kernel region to user buffer */
int copy_from_kernel(void *user_dest, int maxlen) {
    /* Byte count len is minimum of buffer size and maxlen */
    int len = KSIZE < maxlen ? KSIZE : maxlen;
    memcpy(user_dest, kbuf, len);
    return len;
}

• Similar to code found in FreeBSD
• There are legions of smart people trying to find vulnerabilities in programs
/* Kernel memory region holding user-accessible data */
#define KSIZE 1024
char kbuf[KSIZE];

/* Copy at most maxlen bytes from kernel region to user buffer */
int copy_from_kernel(void *user_dest, int maxlen) {
    /* Byte count len is minimum of buffer size and maxlen */
    int len = KSIZE < maxlen ? KSIZE : maxlen;
    memcpy(user_dest, kbuf, len);
    return len;
}

#define MSIZE 528

void getstuff() {
    char mybuf[MSIZE];
    copy_from_kernel(mybuf, MSIZE);
    printf("%s\n", mybuf);
}
Malicious Usage

/* Kernel memory region holding user-accessible data */
#define KSIZE 1024
char kbuf[KSIZE];

/* Copy at most maxlen bytes from kernel region to user buffer */
int copy_from_kernel(void *user_dest, int maxlen) {
    /* Byte count len is minimum of buffer size and maxlen */
    int len = KSIZE < maxlen ? KSIZE : maxlen;
    memcpy(user_dest, kbuf, len);
    return len;
}

#define MSIZE 528

void getstuff() {
    char mybuf[MSIZE];
    copy_from_kernel(mybuf, -MSIZE);
    . . .
}
Reality #4: Asymptotic performance is not always sufficient

- Constant factors matter
- Even operation count might not predict performance
- Must understand system to optimize performance
  - How programs compiled and executed
  - How to measure performance and identify bottlenecks
  - How to improve performance without destroying code modularity and generality
Memory System Performance Example

- Performance depends on access patterns

```c
void copyji(int src[2048][2048],
            int dst[2048][2048])
{
    int i,j;
    for (j = 0; j < 2048; j++)
        for (i = 0; i < 2048; i++)
            dst[i][j] = src[i][j];
}

void copyij(int src[2048][2048],
            int dst[2048][2048])
{
    int i,j;
    for (i = 0; i < 2048; i++)
        for (j = 0; j < 2048; j++)
            dst[i][j] = src[i][j];
}
```

21 times slower
(Pentium 4)
Example Matrix Multiplication

Matrix-Matrix Multiplication (MMM) on 2 x Core 2 Duo 3 GHz (double precision)

Gflop/s

- Standard desktop computer and compiler
- Both implementations have **exactly** the same operations count ($2n^3$)

Best code (K. Goto)

Triple loop

160x
Matrix-Matrix Multiplication (MMM) on 2 x Core 2 Duo 3 GHz

Gflop/s

Reason for 20x: Blocking or tiling, loop unrolling, array scalarization

Effect: fewer register spills, L1/L2 cache misses, and TLB misses
Reality #5: Computer is more than the CPU

- They need to do I/O (get data in and out)
- They communicate with each other over networks
  - Concurrent operations by autonomous processes
  - Coping with unreliable media
  - Cross platform compatibility
  - Complex performance issues
A Little Bit of History

Eckert and Mauchly

- 1\textsuperscript{st} working electronic computer (1946)
- 18,000 Vacuum tubes
- 1,800 instructions/sec
- 3,000 ft\textsuperscript{3}
A Little Bit of History

• Maurice Wilkes

EDSAC 1 (1949)

http://www.cl.cam.ac.uk/UoCCL/misc/EDSAC99/

1st stored program computer
650 instructions/sec
1,400 ft³
A Little Bit of History

- 1954 IBM developed 704
- All programming done in assembly
- Software costs exceed hardware costs!
A Little Bit of History

• Fortran I (project 1954-57)
• The main idea is to translate high level language to assembly
• Many thought this was impossible!
• In 1958 more than 50% of software in assembly!
• Development time halved!

John Backus
(December 3, 1924 – March 17, 2007)
Source Code to Execution

- C source
  - Compiler
    - Assembly
    - Assembler
    - Object File
    - Linker
      - Library
      - Executable
      - Loader
        - DLL
Conclusions

• This first lecture was just an overview. More fun is yet to come!

• Computer system can be viewed as layers of abstractions → knowing these layers helps us see the big picture