

# V22.0453-001 Honors Theory of Computation

## Problem Set 3

All problems are worth 10 points. Due on Tue Nov 9, 2010.

### Problem 1

A *Turing machine with stay put instead of left* is similar to an ordinary Turing machine, but the transition function is of the form:

$$Q \times \Gamma \mapsto Q \times \Gamma \times \{R, S\}$$

At each point the machine can move its head to right or let it stay in the same position. Show that this Turing machine variant is *not* equivalent to the usual version. What class of languages do these machines recognize ?

### Problem 2

Show that the collection of Turing-recognizable languages is closed under the operation of (a) union (b) concatenation (c) star and (d) intersection. What about complementation operation ?

### Problem 3

Show that the following language is decidable:

$$INFINITE_{CFG} = \{\langle G \rangle \mid G \text{ is a context-free grammar such that } L(G) \text{ is infinite}\}$$

### Problem 4

Show that the following language is decidable:

$$L = \{\langle R, S \rangle \mid R \text{ and } S \text{ are regular expressions such that } L(R) \subseteq L(S)\}$$

### Problem 5

Show that the following language is undecidable:

$$A = \{\langle M \rangle \mid M \text{ is a Turing machine that accepts } w \text{ whenever it accepts } w^R\}$$

Here  $w^R$  denotes the reverse of string  $w$ .

**Problem 6** In this problem, we explore the notion of *oracle reducibility*. If  $A$  is a language, then a *Turing machine with oracle  $A$*  is a Turing machine with a "magical" subroutine that decides membership in  $A$ . In other words, the subroutine, when given a string  $w$ , tells the machine whether or not  $w \in A$ . Let

$$HALT_{TM} = \{\langle M, x \rangle \mid M \text{ is a Turing machine that halts on } x\}$$

Show that there is a Turing machine with oracle  $HALT_{TM}$  that decides the following problem with only *two* questions to the oracle: Given three (machine, input) pairs  $\langle M_1, x_1 \rangle, \langle M_2, x_2 \rangle, \langle M_3, x_3 \rangle$ , decide for each pair whether the Turing machine halts on the corresponding input.

Note: This is trivial if one allows *three* questions. Just ask the oracle whether  $\langle M_i, x_i \rangle \in HALT_{TM}$  for  $i = 1, 2, 3$ .